

ON-LINE SYSTEM IDENTIFICATION IN REAL TIME
USING A MINICOMPUTER

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By

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SCOPE AND CONTENTS: By making use of measurements of the input and output of an unknown system, the characterizing parameters are found by using matrix pseudoinverse method. For noise-corrupted measurements, least squares estimation of the parameters are obtained.

Stochastic approximation is employed to improve the estimation.

Both methods are tested on-line in real time using the PDP-11/45 minicomputer while the system is simulated on the TR-20 analog computer.

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TABLE OF CONTENTS

	<u>Page</u>
CHAPTER ONE: INTRODUCTION	1
CHAPTER TWO: SYSTEM IDENTIFICATION	4
2.1 Definition of System Identification	4
2.2 Classification of Identification Methods	5
2.3 Formulation of an Identification Problem	6
CHAPTER THREE: PSEUDOINVERSE APPLIED TO SYSTEM IDENTIFICATION	13
3.1 Definition of Pseudoinverse ...	13
3.2 Properties of the Pseudoinverse	16
3.3. Application to the Solution of a System of Linear Equations	17
3.4 Recursive Identification Algorithm ...	19
3.5 Determination of the Order of the System	22
CHAPTER FOUR: STOCHASTIC APPROXIMATION FOR THE REMOVAL OF BIAS	27
4.1 Effect of Measurement Noise on the Estimation	27
4.2 Removal of Bias by Filtering ...	29
4.3 Stochastic Approximation	30
4.4 Formulation of Noise Model ...	33
4.5 Application of Stochastic Approximation	37
CHAPTER FIVE: ANALOG SIMULATION AND HARDWARE DESCRIPTION	39
5.1 Introduction	39

	<u>Page</u>
5.2 Interface Hardware 	42
5.3 Operating Sequence 	47
5.4 Pseudorandom Input Signal ...	50
5.5 Analog Simulation 	53
5.6 Error Analysis 	55
CHAPTER SIX: SOFTWARE AND RESULTS ...	57
6.1 Introduction	57
6.2 Organization of the Program ...	58
6.3 Results 	64
CHAPTER SEVEN: CONCLUSIONS 	73
APPENDIX: Program Listing 	75
REFERENCES 	107

<u>Figure</u>	<u>LIST OF FIGURES</u>	<u>Page</u>
2.1	Error in Estimation	12
3.1	Error Function V versus Estimated Order \hat{N}	26
5.1	General Schematic of Experiment ...	41
5.2a	Circuit Diagram of Interface Panel ...	44
5.2b	Circuit Diagram of Control Logic ...	45
5.2c	Circuit Diagram of Clock Pulse Generator	46
5.3	Special Registers	49
5.4	Timing Diagram	49
5.5	Spectral Density of Y	52
5.6	Equivalent Impulse Functions	52
5.7	Analog Simulation	54
6.1	Flowchart of Program	59
6.2	Sample Print-out of the Results ...	61
6.3	Table of Subroutines	62
6.4	Table of Estimates of Parameters and Normalised Errors	68
6.5	Plots for Pseudoinverse Algorithm (5% and 100% noise) ...	69
6.6a	Plot for Combined Pseudoinverse and Stochastic Approximation (5% noise)	70
6.6b	Plot for Combined Pseudoinverse and Stochastic Approximation (25% noise)	70

	<u>Page</u>
6.7a Plot for Step Responses (25% noise)	71
6.7b Plot for Impulse Responses (25% noise)	71
6.8 Table of Mean Square Errors in Responses	72

CHAPTER ONE

INTRODUCTION

The terms system identification and process parameter estimation have been used interchangeably by many people. It is the process of obtaining the characteristic parameters of a system by artificially exciting the system with some specially selected signals. Depending on the situation and the need, a great variety of techniques have been developed.

The rapid development in the area of system identification can be attributed to many roots. Prominent among them is that there is a real need to plan and devise better process control. In order to do so, there is the inevitable prerequisite of a good knowledge of the process at hand. Due to the different conditions under which the process will operate, an accurate calculation beforehand is very difficult or sometimes impossible. To complicate the problem further, in the real world environment, we always have random disturbances generally called noise in the system. This naturally makes the characterization of the system even more difficult.

The techniques of system identification lend themselves to a host of applications. The study of the dynamics of many high performance systems such as space vehicles is another example. However, aside from the fact that there exists a need to utilize these techniques in various situations, the availability of the necessary tools is also an important factor. On

the theoretical side, we see tremendous advances in control theory. Also of vital importance is the phenomenal development of computer hardware technology and software. Nowadays, computers are fast enough to accommodate complex calculations in relatively complicated algorithms in much less time than a decade ago. Yet, they are cheap enough to be considered for implementation in many industrial processes. Parallel developments are also found in peripheral devices for data acquisition, interfacing and data transmission. The dramatic reduction in size and weight further facilitate their usage in many more situations.

The present study represents an attempt tending to the goal of utilising the modern computer technology in system identification in a noisy environment.

In the remaining parts of this thesis, the various aspects of implementing two on-line system identification algorithms based on matrix pseudoinverse will be presented. Chapter two is primarily concerned with the definition and classification of different methods of system identification. Ways of constructing a convenient model are also discussed. Chapter three reviews the basics of matrix pseudoinverse together with its application to develop an on-line identification algorithm. Before leaving for the discussion of a second algorithm, we will study methods for the determination of the order of the model of the system. Chapter four begins with an investigation of the effects of measurement noise on estimation and then proposes a method using stochastic approximation to improve the estimation. All the practical aspects of the experiments are presented in chapter five. Results of the experiments with tables and plots are pre-

presented in the next chapter. Chapter seven is the final conclusion of the whole work with an appendix giving the program listing.

CHAPTER TWO

SYSTEM IDENTIFICATION

2.1 Definition of System Identification

While there is not yet a unique definition for system identification agreed upon by everybody, we shall adopt the one suggested by Zadeh [1]:

"Identification is the determination, on the basis of input and output, of a system within a specified class of systems, to which the system under test is equivalent"

Before we proceed, clarification of some of the terms used in the above definition is in order. Systems here refer to the mathematical models of physical systems. It is the first step in the formulation of an identification problem and indeed of any analysis. Modelling of a system generally refers to its representation by a pulse transfer function or by the state variables. Identification, therefore, is the determination of the coefficients in either representation. The meaning of equivalence can be better understood by referring to another concept called identifiability. Astrom and Bohlin [2] suggested that a system is identifiable if the estimates of its parameters are consistent. A necessary condition is that the information matrix be positive definite. When two models are identifiable and result in the same consistent estimates, they are said to be equivalent.

2.2 Classification of Identification Methods

The problem of parameter identification can be viewed as an optimization problem. The solution to this problem consists of finding the extremum of a certain specified loss function as a function of the parameters to be identified.

Computationally, identification methods can be divided into two categories, namely on-line methods and off-line methods.

Off-line method is a one-shot technique where the loss function is an explicit mathematical relation between the measurements and the estimates of interest. It is also known as the open loop method. The solution is available after a fixed finite number of elementary operations. But processing of the data can only start after all measurements are completed. In general, algorithms belonging to this category require considerable computer memory for the processing and storage of the data. Many methods making use of auto- and cross-correlation [3] belong to this category.

On-line methods, on the other hand, are closed loop methods. An iterative scheme whereby the estimation of parameters are being continuously updated as new measurements are made. These intermediate estimates are approximate solutions only. The final solution is approached asymptotically and is therefore, in principle, available only after an infinite number of elementary operations. The objective function in this case is an implicit function of the parameters and some form of model-adjustment strategy is employed to search for the extremum of the objective function. For example, the matrix pseudoinverse method being investigated in the present study

adopts the squares of the residual errors of the system dynamic equation as the objective function. This function is being minimized through an iterative scheme.

The stochastic approximation method can also be formulated in an iterative form for on-line application. This method is being used to identify the noise model in the present study. Both methods will be described in greater detail later.

The on-line methods have the obvious advantage over the off-line methods in that intermediate results are available for use during the identification process if desired. This type of estimation will also be able to respond to a change in the system dynamics. When applied in real time, however, they pose more problems in their implementation. A useful on-line method must have the important property that it be computationally efficient. The reason is obvious because we have to complete all necessary calculations in updating the estimates before the next set of new measurements can be made. An inefficient algorithm will impose severe constraints on the choice of sampling frequencies. Moreover, in an industrial setting, identification is only part of the control loop. The computer is likely to be used for other purposes as well. Thus, the ease of implementation, and hence the feasibility for realistic application, depends heavily on the computational simplicity.

2.3 Formulation of an Identification Problem

The formulation of an identification problem begins with the model-

ling of the system. Preferably, the model chosen should have the following properties:-

- (a) It is based on the input-output measurements and does not depend on other measurements that might be difficult or impossible to be made directly;
- (b) It must assume a simple form with the smallest possible number of parameters to be identified and can readily be used for controller design;
- (c) It should be able to accommodate the stochastic behaviour of the system due to random disturbances.

Several basic assumptions are also made in the modelling which would greatly reduce the amount of work without severely limiting the usefulness. We shall assume that the system is linear or can be adequately approximated by a linearised model. Further the system is assumed to be time-invariant. If the system is only slowly time-varying, the changes can be reflected on the estimates obtained by on-line algorithms. Finally, the physical system is assumed to be of finite order.

There are two convenient types of models available, namely the state-variable model and the pulse transfer function model.

The state-variable model of a linear, finite dimensional and time-invariant discrete time system is given by

$$x(k+1) = Ax(k) + Bu(k) \quad (2.1)$$

$$y(k) = Cx(k) \quad (2.2)$$

where $k \geq 0$ is an integer,

x is the state vector of dimension m ,

A is a $m \times m$ matrix,

B is a $m \times 1$ vector,

C is a $1 \times m$ vector,

u is the input,

and y is the output.

The parameters are the elements in matrices A , B and C . Gupta [4] has investigated this type of model in detail applied to system identification.

The transfer function counterpart of equations (2.1) and (2.2) are as follows:-

$$\begin{aligned} H(z^{-1}) &= \frac{C(z^{-1})}{R(z^{-1})} \\ &= \frac{a_0 + a_1 z^{-1} + \dots + a_m z^{-m}}{1 + b_1 z^{-1} + \dots + b_n z^{-n}} \end{aligned} \quad (2.3)$$

where $C(z^{-1})$ is a polynomial for the output of order m

$R(z^{-1})$ is a polynomial for the input of order n

The parameter vector for identification can be defined as

$$\phi^T = [a_0 a_1 \dots a_m b_1 b_2 \dots b_n] \quad (2.4)$$

where superscript T denotes matrix transpose.

The transfer function model is more suitable for the equation-error approach since it is a direct relation between the input and output. For this reason, it is chosen for the present study.

In discrete time, equation (2.3) can be written as

$$c_i = \sum_{j=0}^m a_j r_{i-j} - \sum_{j=1}^n b_j c_{i-j} \quad (2.5)$$

where $c_i = c(iT)$, output of system at $t = iT$
 $r_i = r(iT)$, input of system at $t = iT$
 T = sampling period
 i = integer

In matrix notation, letting i range from 1 to same integer k , we can again rewrite (2.5) as

$$A'_k \phi = c_k \quad (2.6)$$

where

$$A'_k = \begin{bmatrix} r_0 & r_{-1} & \dots & r_{1-m} & -c_{-1} & -c_{-2} & \dots & -c_{1-n} \\ r_1 & r_0 & \dots & r_{2-m} & -c_0 & -c_{-1} & \dots & -c_{2-n} \\ \vdots & \vdots & & \vdots & \vdots & \vdots & & \vdots \\ \vdots & \vdots & & \vdots & \vdots & \vdots & & \vdots \\ r_{k-1} & r_{k-2} & \dots & r_{k-m} & -c_{k-2} & -c_{k-3} & \dots & -c_{k-n} \end{bmatrix} \quad (2.7)$$

and

$$c_k^T = [\hat{c}_1 \ c_2 \ \dots \ c_k] \quad (2.8)$$

Equation (2.6) can be solved analytically if $m+n+1$ pairs of r_i and c_i are known exactly, assuming known initial conditions, or twice as many pairs of measurements if initial conditions are unknown.

However, the measurements are usually contaminated with noise and we have the input-output measurements as follows:

$$u_i = r_i + w_i \quad (2.9)$$

$$y_i = c_i + v_i \quad (2.10)$$

where $\{w_i\}$ and $\{v_i\}$ are noise sequences for input and output respectively.

Therefore, equations (2.6), (2.7) and (2.8) will be replaced by

$$A_k \hat{\phi}_{\sim} = y_k \quad (2.11)$$

where

$$A_k = \begin{bmatrix} u_0 & u_{-1} & \dots & u_{1-m} & -y_{-1} & \dots & y_{1-n} \\ u_1 & u_0 & \dots & u_{2-m} & -y_0 & \dots & y_{2-n} \\ \vdots & \vdots & & \vdots & \vdots & & \vdots \\ \vdots & \vdots & & \vdots & \vdots & & \vdots \\ u_{k-1} & u_{k-2} & \dots & u_{k-m} & -y_{k-2} & \dots & y_{k-n} \end{bmatrix} \quad (2.12)$$

$$\hat{y}_k^T = [y_1 \ y_2 \ \dots \ y_k] \quad (2.13)$$

and the parameter vector

$$\hat{\phi}_{\sim}^T = [\hat{a}_0 \hat{a}_1 \ \dots \ \hat{a}_m \hat{b}_1 \hat{b}_2 \ \dots \ \hat{b}_n] \quad (2.14)$$

In subsequent chapters, an algorithm will be developed to solve $\hat{\phi}$ in an iterative manner suitable for on-line application based on the noise-corrupted input-output measurements. The error of estimation E_R is defined as the difference between the estimated vector $\hat{\phi}$ and the true parameter vector ϕ ,

$$E_R = \phi - \hat{\phi} \quad (2.15)$$

Very often, it is convenient to express the error as a scalar quantity which can be normalised for easy comparison. Thus, we define the normalised error of the estimated parameter vector as

$$e_R = \frac{||\phi - \hat{\phi}||^2}{||\phi||^2} \quad (2.16)$$

where $|| \quad ||$ denotes the norm.

Pictorially, the error due to an incorrect estimation of the parameter vector can be depicted in figure 2.1.

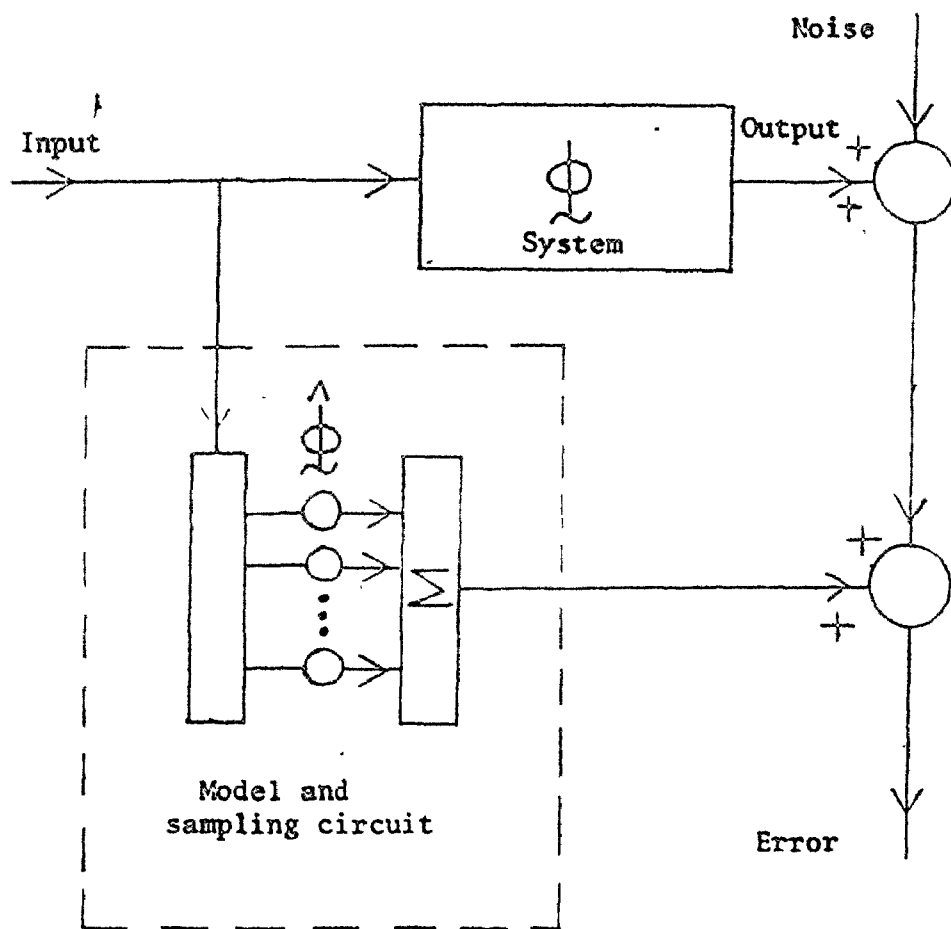


Figure 2.1 Error in Estimation

CHAPTER THREE

PSEUDOINVERSE APPLIED TO SYSTEM IDENTIFICATION

3.1 Definition of Pseudoinverse

In solving the system of linear equations

$$Ax = y$$

we have

$$x = A^{-1}y$$

where A^{-1} is called the inverse matrix of matrix A such that A must be a square non-singular matrix with the property that

$$AA^{-1} = A^{-1}A = I$$

Penrose [5] generalized the idea of matrix inverse to include cases where A is rectangular and gave it the name generalized inverse or pseudoinverse of a matrix. Greville [6] and others also have investigated its properties and applications. One way of defining [6] the pseudoinverse is as follows:

Let A be a matrix of dimension $m \times n$ with rank equal to r . It can be factorized into two matrices B and C such that

$$A = BC \tag{3.1}$$

where B is a $m \times r$ matrix of rank r

C is a $r \times n$ matrix of rank r

$r \leq m, n$ integers

The factorization can be found by first selecting B such that its columns are the linearly independent columns of A . Since A is of rank r , the dimension of B must be $m \times r$. Then C is chosen such that it will satisfy equation (3.1).

The pseudoinverse of A is then given by

$$\begin{aligned} A^+ &= C^T [CC^T]^{-1} [B^T B]^{-1} B^T \quad \text{if } A \neq 0 \\ &= A^T \quad \text{if } A = 0 \end{aligned} \quad (3.2)$$

It can be proved [7] that there always exists a unique pseudoinverse A^+ as defined above for any matrix A .

In the special case when $n=r$, (3.1) reduces to

$$A = BI$$

and (3.2) is simplified to

$$\begin{aligned} A_L^+ &= [B^T B]^{-1} B^T \\ &= [A^T A]^{-1} A^T \end{aligned} \quad (3.3)$$

where A_L^T is called the left pseudoinverse of A with the rows in the row space of A^T such that $A_L^+ A = I$

Example:

$$\text{Let } A = \begin{bmatrix} 1 & 0 \\ 0 & 2 \\ 2 & 1 \end{bmatrix}$$

Here $m=3$, $n=r=2$

$$\begin{aligned} \text{Using (3.3)} \quad A_L^+ &= [A^T A]^{-1} A^T \\ &= \frac{1}{21} \begin{bmatrix} 5 & -4 & 8 \\ -2 & 10 & 1 \end{bmatrix} \end{aligned}$$

$$A_L^+ A = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} = I$$

Similarly, when $m=r$, (3.1) reduces to

$$A = IC$$

and (3.2) becomes

$$\begin{aligned} A_R^+ &= C^T [CC^T]^{-1} \\ &= A^T [AA^T]^{-1} \end{aligned} \quad (3.4)$$

where A_R^+ is called the right pseudoinverse of A with the columns in the column space of A^T such that $AA_R^T = I$

Example:

$$\text{Let} \quad A = \begin{bmatrix} 1 & 2 \end{bmatrix}$$

$$\text{Here} \quad m=r=1, \quad n=2$$

$$\text{Using (3.4)} \quad A_R^T = [A^T A]^{-1} A^T$$

$$= \frac{1}{5} \begin{bmatrix} 1 & 2 \end{bmatrix}$$

$$AA_R^+ = 1 = I$$

Comparing equations (3.2), (3.3) and (3.4), we can deduce that

$$A^+ = A_L^+ A_R^+ \quad (3.5)$$

3.2 Properties of the Pseudoinverse

Before going on to consider the use of pseudoinverse for system identification, we have to look more closely into its properties. Many of these properties will be used in the derivation of the identification algorithm

For every real $m \times n$ matrix A , there exists a unique real pseudo inverse A^+ as defined earlier which will satisfy the following identities

$$\begin{aligned} A^+ A A^+ &= A^+ \\ A A^+ A &= A \\ [A A^+]^T &= A A^+ \\ [A^+ A]^T &= A^+ A \end{aligned} \quad (3.6)$$

where the superscript T denotes transpose of a matrix.

In fact, the above identities are used by Penrose [4] to define pseudoinverse in his original paper. They form a set of necessary and sufficient conditions to prove the existence and uniqueness of the pseudo-inverse.

Other properties are summarized below [8]:-

1. $A^{++} = A$
2. $A^{T+} = A^{+T}$
3. $A^+ = A^{-1}$ if A is square and nonsingular
4. $[\lambda A]^+ = \lambda^+ A^+$, λ scalar

$$\text{and } \lambda^+ = \frac{1}{\lambda} \text{ for } \lambda \neq 0$$

$$\lambda^+ = 0 \text{ for } \lambda = 0$$

$$5. [A^T A]^+ = A^+ A^{T+}$$

$$[AA^T]^+ = A^{+T} A^+$$

$$\text{in general } [AB]^+ \neq B^+ A^+$$

$$6. \text{ The ranks of } A, A^T A, A, A^+ A \text{ are all equal to the trace of } A^+ A$$

3.3 Application to the Solution of A System of Linear Equations

In this section, we shall prove an important theorem which serves to illustrate the crucial value of matrix pseudoinverse in the solution of a system of linear equations. This in turn plays an important role in the development of the algorithm for system identification.

Recall the definition of the pseudoinverse given in section 3.1 that

$$\text{if } A = BC \tag{3.1}$$

$$\text{then } A^+ = C^T [CC^T]^{-1} [B^T B]^{-1} B^T \tag{3.2}$$

Theorem [9]:

With the pseudoinverse A^+ of matrix A defined as in equations (3.1) and (3.2), the solution of the system of linear equations

$$\chi = A\chi \tag{3.7}$$

which will minimize

$$(a) \text{ the sum of squares of the residuals } \mathcal{E}^T \mathcal{E}$$

$$\text{where } \mathcal{E} = \chi - A\chi \tag{3.8}$$

(b) the sum of squares of the unknown $x^T x$ is given by

$$x = A^+ y \quad (3.9)$$

Proof: Let $S = (y - Ax)^T (y - Ax)$

$$= x^T x - x^T A^T y - y^T A x + x^T A^T A x$$

minimizing S

$$\frac{\partial S}{\partial x_j} = 0, \quad j = 1, 2, 3, \dots, n$$

$$\text{we have } A^T A x = A^T y \quad (3.10)$$

Now, substitute equation (3.1) into (3.10), we have

$$C^T [B^T B] C x = C^T B^T y \quad (3.11)$$

Multiply both sides by $[B^T B]^{-1} [C^T]^{-1}$ and rearranging, we get

$$\begin{aligned} x &= C^T [C^T]^{-1} [B^T B]^{-1} B^T y \\ &= A^+ y \end{aligned} \quad (3.12)$$

by applying equation (3.2). Hence the assertions (a) and (b) follow.

In other words, the solution given by $x = A^+ y$ has the important consequence that it is also the optimal solution in the sense that the square of the residual error is minimized. The identification algorithm developed according to this same principle will likewise yield optimal estimates in the same sense.

3.4 Recursive Identification Algorithm

Recall that we have formulated the parameter estimation problem in section 2.3 as

$$A_k \hat{\phi}_k = y_k \quad (2.11)$$

Applying what we have just developed in the preceding section, we have at the k th iteration

$$\hat{\phi}_k = A_k^+ y_k \quad (3.13)$$

with $\hat{\phi}_k$ as the optimal estimation of ϕ_k .

If $\hat{\phi}_k$ has $p=m+n+1$ elements as defined in section 2.3, we may form the following special cases:

(a) For $k \leq p$,

$$\hat{\phi}_k = A_k^T [A_k A_k^T]^{-1} y_k \quad (3.14)$$

where $A_k^T [A_k A_k^T]^{-1}$ is the right pseudoinverse of A_k . It gives the so called minimum norm solution of $\hat{\phi}_k$.

(b) For $k > p$

$$\hat{\phi}_k = [A_k^T A_k]^{-1} A_k^T y_k \quad (3.15)$$

where $[A_k^T A_k]^{-1} A_k^T$ is the left pseudoinverse of A_k . It gives the least squares solution of $\hat{\phi}_k$.

In both cases, A_k is assumed to have full rank. This condition can be guaranteed if one of the following conditions [10] is imposed on the

input sequence $r(iT)$:-

1. $r(iT)$ is a sequence composed of n discrete Fourier components and all natural modes are present in the output sequence C_i .
2. $r_i = 0$ for $k < n$
 $= 1$ for $k \geq n$
3. $r(iT)$ is a random signal.

More will be said about the input signal for the experiments in chapter five.

The transformation of equation (3.13) into an iterative formula has been considered by Wells [11] and Sinha and Pille [12]. The information matrix A_{k+1} is considered to be formed in the following manner:

$$A_{k+1} \triangleq \begin{bmatrix} A_k \\ \mathcal{R}_{k+1}^T \end{bmatrix} \quad (3.16)$$

where $\mathcal{R}_{k+1}^T = [u_{k+1} \ u_k \ \dots \ u_{k-m} \ y_k \ y_{k-1} \ \dots \ y_{k-n}]$ (3.17)

a row vector containing the latest set of measurements.

Similarly, the output vector χ_{k+1} is of the form

$$\chi_{k+1} = \begin{bmatrix} \chi_k \\ y_{k+1} \end{bmatrix} \quad (3.18)$$

where y_{k+1} is the latest measurement of the output of the system at the $(k+1)$ th instant corresponding to the input u_{k+1} .

The result is that, when a new pair of input-output measurements is made, a new row is added to the information matrix A_{k+1} and a new element is added to the output vector χ_{k+1} .

With these arrangements, we are now in a position to derive the iterative algorithm [12]. The major results for the k th iteration are summarized below:

Let p be the dimension of vector $\hat{\phi}_k$

For $k \leq p$ (minimum norm solution)

$$\hat{\phi}_{k+1} = \hat{\phi}_k + \frac{Q_{k+1} \hat{a}_{k+1}}{\hat{a}_{k+1}^T Q_k \hat{a}_{k+1}} (y_{k+1} - \hat{a}_{k+1}^T \hat{\phi}_k) \quad (3.19)$$

where

$$Q_{k+1} = Q_k - \frac{[Q_k \hat{a}_{k+1}]^T [Q_k \hat{a}_{k+1}]^T}{\hat{a}_{k+1}^T Q_k \hat{a}_{k+1}} \quad (3.20)$$

and

$$P_{k+1} = P_k + \frac{[Q_k \hat{a}_{k+1}] [Q_k \hat{a}_{k+1}]^T [1 + \hat{a}_{k+1}^T P_k \hat{a}_{k+1}]}{[\hat{a}_{k+1}^T Q_k \hat{a}_{k+1}]^2} - \frac{[P_k \hat{a}_{k+1}] [Q_k \hat{a}_{k+1}]^T + [Q_k \hat{a}_{k+1}] [P_k \hat{a}_{k+1}]^T}{\hat{a}_{k+1}^T Q_k \hat{a}_{k+1}} \quad (3.21)$$

with the initial conditions

$$Q_0 = I, \quad P_0 = 0 \text{ and } \hat{\phi}_0 = 0 \quad (3.22)$$

For $k > p$ (least squares estimation)

$$\hat{\phi}_{k+1} = \hat{\phi}_k + \frac{[P \hat{a}_{k+1}] [y_{k+1} - \hat{a}_{k+1}^T \hat{\phi}_k]}{1 + \hat{a}_{k+1}^T P_k \hat{a}_{k+1}} \quad (3.23)$$

$$\text{where } P_{k+1} = P_k - \frac{[P_k \hat{a}_{k+1}] [P_k \hat{a}_{k+1}]^T}{1 + \hat{a}_{k+1}^T P_k \hat{a}_{k+1}} \quad (3.24)$$

The matrices P_k and Q_k are defined as

$$P_k = A_k + A_k^{*T} \quad (3.25)$$

$$Q_k = I - A_k + A_k \quad (3.26)$$

Thus they are both symmetric. The dimensions of matrices P_k and Q_k are both $p \times p$ while that of the vectors $\hat{\phi}_k$ and \hat{a}_k are both p . The storage requirement is minimal. Further, only a total of

$$N = 3p^2 + 4p$$

multiplications are required to calculate equations (3.23) and (3.24) so that this algorithm is regarded as computationally efficient.

The results of the experiments applying this algorithm to the identification of a second order system will be described in greater details in chapter six. Meanwhile, making use of some of the definitions just presented, we shall discuss methods to determine the order of the model.

3.5 Determination of the order of the System

The algorithm discussed in the previous section assumes that we know the value of

$$p = m+n+1 \quad (3.27)$$

where m and n are the number of coefficients in the numerator and denominator of equation (2.3) respectively. Naturally, the values of m and n depend on the order of the model used for identification.

One method to determine the order of the model is to assume certain low values for m and n in equation (3.27) and proceed with the estimation of $\hat{\phi}$. Then, increase m and n by one and repeat the estimation procedures. The step responses of the last two trials are compared. If they are the same or sufficiently close, the latter model is of unnecessary high order. Otherwise, the trial process is carried on until we meet such a requirement. This is admittedly a very crude trial and error approach. We might have problems if the noise level is high.

The following method is a much more systematic approach proposed by Sinha and Pille [12]. It is based on the following theorem.

Theorem:

Consider a system transfer function

$$H(z) = \frac{P(z^{-1})}{Q(z^{-1})} \quad (3.28)$$

where $P(z^{-1})$ and $Q(z^{-1})$ are polynomials of order m and n respectively.

Assume that, in the model, the order of both the numerator and denominator is N which may be chosen arbitrary large. Let

$$q = \text{tr } Q_k \quad (3.29)$$

where

$$Q_k = I - A_k^* A_k \quad (3.30)$$

which can be obtained iteratively from equation (3.20). If k is incremented from 1 to M ($M \leq 2N$) until q becomes a constant, the true order of the system is given by

$$n = N - q \quad (3.31)$$

Proof:

$$\begin{aligned} \text{Rank } A_k &= \text{tr}(A_k^* A_k) \\ &= 2N - \text{tr } Q_k \\ &= 2N - q \end{aligned} \quad (3.32)$$

The maximum rank of A_k is $N+n$, since A_k has $2n + (N-n)$ degrees of freedom. Thus, when A_k has attained a maximum rank, q becomes a constant, and

$$\begin{aligned} N+n &= 2N - q \\ \text{or } n &= N - q \end{aligned}$$

In the event when the order $m < n$, the algorithm should yield zero for the estimates of $\hat{a}_{m+1}, \hat{a}_{m+2}, \dots, \hat{a}_n$ etc. in equation (2.14) and the order of the system is apparent.

However, the above derivation assumed the effect of noise to be insignificant so that q would approach a constant value. When the noise level is high, we might have difficulty in applying this method.

A third method to the same end takes the presence of noise into account. It is worthwhile to note that, in a noisy environment, several independent methods are sometimes necessary since each method has its own limitations due to the assumptions made in each of them. If they all give the same result, we can be confident that it is the correct value.

The system dynamic equation defined in section 2.3 is

$$x_k = A_k \hat{\phi}_k + e \quad (2.11)$$

where

x_k = output measurement vector with noise

A_k = information matrix with noisy measurements

$\hat{\phi}_k$ = estimation vector at the kth instant

$$= [\hat{a}_0 \ \hat{a}_1 \ \dots \ \hat{a}_m \ \hat{b}_1 \ \hat{b}_2 \ \dots \ \hat{b}_n]^T \quad (2.14)$$

Let the error function be defined as

$$v = \frac{1}{k} e^T e \quad (3.35)$$

According to Van den Boom and Van den Enden [13], assume the noise is white and taking the probability limit of v , we have

$$\text{plim}_{k \rightarrow \infty} [v] = \text{plim}_{k \rightarrow \infty} \left[\frac{1}{k} e^T e \right] \quad (3.34)$$

Let \hat{N} be an estimation of the true order N of the system, this leads to the following consideration by taking into account the asymptotic properties of the estimates $\hat{\phi}_k$.

$$\text{plim}_{k \rightarrow \infty} [\hat{a}_i] = \begin{cases} \neq a_i & \text{if } \hat{N} < N \quad \text{due to truncation effect} \\ = a_i & \text{if } \hat{N} = N \\ = a_i & i \leq N \\ = 0 & i > N \end{cases} \quad \text{if } \hat{N} > N \quad (3.35)$$

$$\text{plim}_{k \rightarrow \infty} [\hat{b}_i] = \begin{cases} \neq b_i & \text{if } \hat{N} < N \quad \text{due to truncation effect} \\ = b_i & \text{if } \hat{N} = N \\ = b_i & i \leq N \\ = 0 & i > N \end{cases} \quad \text{if } \hat{N} > N \quad (3.36)$$

These conditions result in

$$\text{plim} \left[\frac{1}{k} \hat{\epsilon}^T \hat{\epsilon} \right] = \begin{cases} > 0 & \text{for } \hat{N} < N \\ = 0 & \text{for } \hat{N} \geq N \end{cases} \quad (3.37)$$

An important consequence is apparent in that there is a marked change in the behaviour of the error function v when $\hat{N} = N$. Pictorially, its behaviour is shown in figure 3.1.

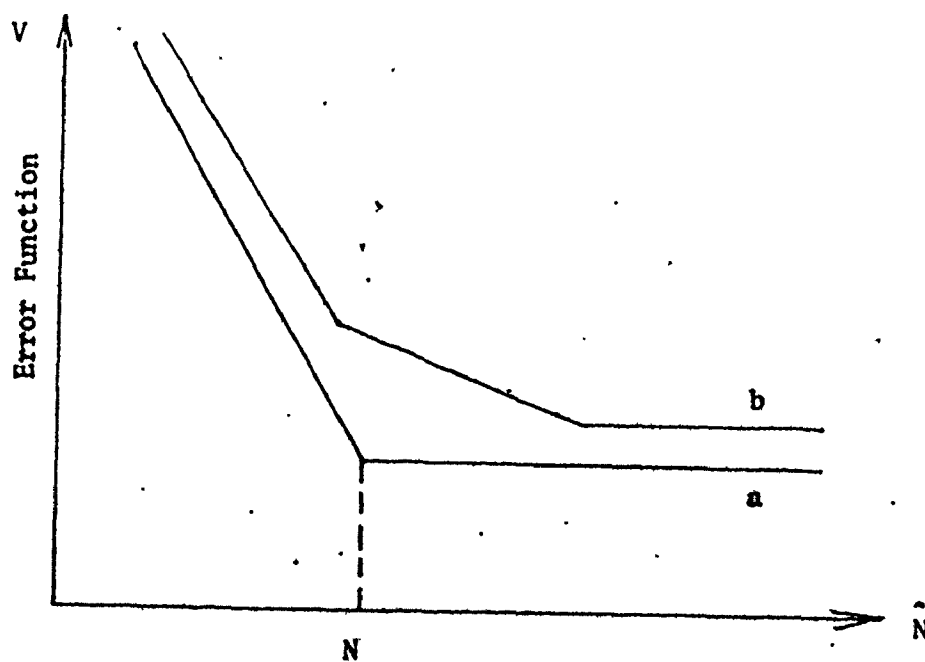


Figure 3.1 Error Function V versus Estimated Order \hat{N}

CHAPTER FOUR

STOCHASTIC APPROXIMATION FOR THE REMOVAL OF BIAS

4.1 Effect of Measurement Noise on the Estimation

Again refer to the basic dynamic equation

$$A_k \hat{\phi}_k = \chi_k \quad (2.11)$$

where the solution is given by

$$\hat{\phi}_k = A_k^{-1} \chi_k \quad (3.13)$$

Matrix A_k is the information matrix containing the input-output measurements. As noted earlier, these measurements are contaminated with noise. We shall analyze the effect of noise on the final estimates.

Let matrix A_k be decomposed into two component matrices as follows (dropping the subscript k):

$$A = B + N \quad (4.1)$$

where B is the noise-free component of A and N is the noise matrix.

Similarly, the output can be decomposed to

$$\chi = \xi + \gamma \quad (4.2)$$

where ξ is the true output and γ is the output noise.

As shown in the previous chapter, for least squares estimation, the pseudoinverse A^+ is given by

$$A^+ = [A^T A]^{-1} A^T$$

Substituting equation (4.1)

$$\begin{aligned} A^+ &= [(B+N)^T (B+N)]^{-1} [B+N]^T \\ &= [B^T B + B^T N + N^T B + N^T N]^{-1} [B+N]^T \end{aligned} \quad (4.3)$$

If the noise is white and hence uncorrelated with the input and output, we have

$$B^T N \rightarrow 0 \quad N^T B \rightarrow 0$$

Then, equation (4.3) becomes

$$A^+ = [B^T B + N^T N]^{-1} [B+N]^T \quad (4.4)$$

Using the identity [14]

$$[B^T B + N^T N]^{-1} = [B^T B]^{-1} [I + (N^T N)^{-1} (B^T B)]^{-1} [B+N]^T \quad (4.5)$$

equation (4.4) becomes

$$A^+ = [I-Z] B^+ + [I-Z] [B^T B]^{-1} N^T \quad (4.6)$$

where

$$\begin{aligned} Z &\triangleq [I + (N^T N)^{-1} (B^T B)]^{-1} \\ &= [N^T N + B^T B]^{-1} [N^T N] \end{aligned} \quad (4.7)$$

Substituting (4.6) and (4.7) into (4.1), we get .

$$\begin{aligned}
 \hat{\phi} &= [I-Z] B^+ \xi + [I-Z] [B^T B]^{-1} N^T \xi \\
 &= [I-Z] B^+ \chi + [I-Z] [B^T B]^{-1} N^T \chi
 \end{aligned} \tag{4.8}$$

For uncorrelated white noise

$$N^T \xi \rightarrow 0$$

$$B^+ \chi \rightarrow 0$$

$$N^+ \chi \rightarrow 0$$

Hence, equation (4.8) becomes

$$\begin{aligned}
 \hat{\phi} &= [I-Z] B^+ \xi \\
 &= [I-Z] \phi
 \end{aligned} \tag{4.9}$$

$$\text{Since } \phi = B^+ \xi$$

$$\text{Thus } \phi - \hat{\phi} = [N^T N + B^T B]^{-1} [B^T N] \tag{4.11}$$

$$\neq 0 \quad \text{if } N \neq 0$$

The difference given by equation (4.11) is called the bias of the estimation due to the presence of measurement noise.

4.2 Removal of Bias by Filtering

The result in the previous section holds true only for white uncorrelated additive noise. In practice, the noise is coloured. Coloured noise can be modelled as the output resulting from passing white noise through a finite order transfer function. Recognizing this fact, we can devise a method to find out this noise model. From this model, we can find its inverse

with which a filter is constructed. With this filter, we can remove the noise component from the measurements by passing them through this filter.

This approach is not unlike the Wiener-Hopf filtering method. In order to achieve this, we have to first identify the noise model and secondly develop a filtering mechanism. All these have to be integrated with the pseudoinverse algorithm we have just discussed. Further it must be an iterative algorithm so that it can be applied on-line. Finally, the amount of extra computation involved should be as little as possible in order to end up with a still efficient algorithm.

Sen and Sinha [14] have proposed a scheme by applying stochastic approximation to find the noise filter working in parallel with the pseudo-inverse algorithm. The derivation of this algorithm will be given in a later section while we pause to introduce the basic principles of stochastic approximation.

4.3 Stochastic Approximation

Stochastic approximation may be regarded as a scheme for successive approximation of a sought quantity when the observations involve random errors due to the stochastic nature of the problem. It has the following advantages:

- (a) Only a small interval of data needs processing.
- (b) Only simple computations are required.
- (c) A priori knowledge of the process statistics is not required, nor is the functional relationship between the desired para-

eters and the observed data. The only requirements are that it satisfies certain regularity conditions and that a unique solution exists.

Many major contributions are made to the area by various people. A comprehensive survey paper by Sakrison [15] gives a good general picture of various aspects of the subject.

First, let us look at the Robbins-Monro approach [16] which is the statistical analogue of the simple gradient method for finding the root of the equation

$$h(x) = 0 \quad (4.12)$$

which is
$$x_{i+1} = x_i - K_i h(x_i) \quad (4.13)$$

where $\{K_i\}$ is a sequence of real numbers which must satisfy certain conditions to ensure that the algorithm will converge.

When there is additive random noise, $h(x)$ becomes

$$Z(x_i) = h(x_i) + v_i \quad (4.14)$$

where $\{v_i\}$ is a zero mean noise sequence.

Making use of the fact that the expectation of $Z(x_i)$ is $h(x_i)$, equation (4.13) may be modified to

$$x_{i+1} = x_i - K_i Z(x_i) \quad (4.15)$$

Robbins and Monro showed that equation (4.15) will converge if the following conditions are met:

$$\begin{aligned} \lim_{i \rightarrow \infty} K_i &= 0, \\ \sum_{i=1}^{\infty} K_i &= \infty, \end{aligned} \quad (4.16)$$

$$\text{and} \quad \sum_{i=1}^{\infty} K_i^2 < \infty$$

A simple sequence which meets these requirements is

$$K_i = \frac{\alpha}{\beta + i} \quad (4.17)$$

where α and β are positive constants. Also, it is required that $h(x)$ be bounded on either side of a true solution by straight lines, such that it is not possible to overshoot the solution x which cannot be corrected by a K_i satisfying equation (4.17).

Kiefer and Wolfowitz [17] extended the method to find the extremum of an unknown unimodal regression function $\theta(u)$. This approach is the exact analogue of the gradient approach in the deterministic optimization procedure which yields

$$U_{i+1} = U_i - K_i \frac{d \theta(U_i)}{d U_i} \quad (4.18)$$

The stochastic counterpart is

$$U_{i+1} = U_i - K_i \frac{d \eta(U_i)}{d U_i} \quad (4.19)$$

where $\eta(u) = \theta(u) + \xi$

and ξ is the random noise component.

Since the differentiation in equation (4.19) does not exist in general, one may use the following approximation

$$\frac{d \eta(U_i)}{d U_i} \approx \frac{\eta(U_i + \Delta U_i) - \eta(U_i - \Delta U_i)}{2 \Delta U_i} \quad (4.20)$$

Convergence is guaranteed if the following conditions are satisfied:

$$\begin{aligned} \lim_{i \rightarrow \infty} K_i &= 0, \\ \lim_{i \rightarrow \infty} \Delta U_i &= 0, \\ \sum_{i=1}^{\infty} K_i &= \infty, \\ \sum_{i=1}^{\infty} K_i^2 &< \infty, \\ \text{and } \sum_{i=1}^{\infty} \left[\frac{k_i}{\Delta U_i} \right]^2 &< \infty \end{aligned} \quad (4.21)$$

A basic idea [15] of stochastic approximation is that a stochastic counterpart exists for any deterministic algorithm. Fu et al [20], Sinha and Griscik [18] and Kwany [19] have proposed specific formulae to implement the above ideas.

4.4. Formulation of the Noise Model

Let the system dynamic equation be represented in the following form:

$$[1 + A(z^{-1})]y_i = [b_0 + B(z^{-1})]U_i + e_i \quad (4.22)$$

where $e_i = [1 + A(z^{-1})]\eta_i$, the residual error (4.23)

$\{\eta_i\}$ is a zero mean random noise sequence

$$A(z^{-1}) = a_1 z^{-1} + a_2 z^{-2} + \dots + a_n z^{-n}$$

$$B(z^{-1}) = b_1 z^{-1} + b_2 z^{-2} + \dots + b_m z^{-m}$$

To guarantee stability of the process, the roots of $[1 + A(z^{-1})]$ are assumed to lie inside the unit circle.

We assume that the noise sequence $\{\eta_i\}$ can be described as a filtering of a well behaved, zero mean white noise signal ξ_i i.e.

$$\eta_i + \sum_{j=1}^p d_j \eta_{i-j} = \xi_i + \sum_{i=1}^q g_i \xi_{i-j}$$

or equivalently,

$$\eta_i = \frac{1 + G(z^{-1})}{1 + D(z^{-1})} \xi_i \quad (4.24)$$

where $D(z^{-1}) = d_1 z^{-1} + d_2 z^{-2} + \dots + d_p z^{-p}$

$$G(z^{-1}) = g_1 z^{-1} + g_2 z^{-2} + \dots + g_q z^{-q}$$

Again, the roots of $[1 + D(z^{-1})]$ are assumed to lie within the unit circle.

Combining (4.23) and (4.24), we have

$$e_i = \frac{[1 + A(z^{-1})][1 + G(z^{-1})]}{[1 + D(z^{-1})]} \xi_i \quad (4.25)$$

The above residual error sequence $\{e_i\}$ is now approximated by a low order linear process. The filter involved is the inverse of this process. This is similar to the method of "pre-whitening" used in spectral density estimation. Two possible processes are suitable for this purpose. They are the moving average process of the form

$$e_i = \xi_i + \sum_{r=1}^P m_r z^{-r} \xi_i \quad (4.26)$$

and the autoregressive process of the form

$$e_i + \sum_{r=1}^P f_r z^{-r} e_i = \xi_i \quad (4.27)$$

These processes are duals of each other as a moving average process filtered by an autoregressive filter becomes a white noise or vice versa.

In the present study, an autoregressive model is chosen. In particular, if we have cascaded filters of the type

$$\hat{e}_i' = [1 + \sum_{j=1}^P f_j z^{-j}] e_i' \quad (4.28)$$

we can approximate the true process to any degree of accuracy by choosing an appropriate sequence $\{f_j\}$. Using this principle, equation (4.25) can therefore be approximated by

$$e_i = \frac{\xi_i}{[1 + F_s(z^{-1})]} \quad (4.29)$$

where

$$F_s(z^{-1}) = f_1 z^{-1} + f_2 z^{-2} + \dots + f_s z^{-s}$$

This implies that

$$[1 + F_s(z^{-1})] = [1 + D(z^{-1})] [1 + A(z^{-1})]^{-1} [1 + G(z^{-1})]^{-1}$$

This is true if sufficient number of terms of the filter on the left hand side are used.

Substitute (4.29) into (4.22), we have

$$[1 + A(z^{-1})] y_i = [b_0 + B(z^{-1})] u_i + \frac{\xi_i}{[1 + F_s(z^{-1})]}$$

Multiplying throughout by $[1 + F_s(z^{-1})]$, we get

$$[1 + A(z^{-1})] [1 + F_s(z^{-1})] y_i = [b_0 + B(z^{-1})] [1 + F_s(z^{-1})] u_i + \xi_i$$

or

$$[1 + A(z^{-1})] \bar{y}_i = [b_0 + B(z^{-1})] \bar{u}_i + \xi_i \quad (4.30)$$

where

$$\begin{aligned} \bar{y}_i &= (1 + f_1 z^{-1} + f_2 z^{-2} + \dots + f_s z^{-s}) y_i \\ \bar{u}_i &= (1 + f_1 z^{-1} + f_2 z^{-2} + \dots + f_s z^{-s}) u_i \end{aligned} \quad (4.31)$$

ξ_i is a white noise sequence.

If the filter is known, the measurements y_i and u_i will be filtered in such a manner as in equation (4.31) to obtain the filtered input-output

pair \bar{u}_i and \bar{y}_i . These filtered quantities are then used in the algorithm to calculate the estimates. Comparing equations (4.22) and (4.31), the residual error sequence has now been changed to a white noise sequence. The final estimates will therefore be unbiased because the white residual error is not correlated with the input and output.

The next logical step is to find the parameters of the filter $F_s(z^{-1})$. This is where stochastic approximation comes into the picture and is the subject of discussion in the next section.

4.5 Application of Stochastic Approximation

At the k th iteration, the residual error is represented by

$$\hat{e}_k = F^T E_k + W_k \quad (4.32)$$

where $F^T = [f_1 \ f_2 \ \dots \ f_s]^T$, the filter parameter vector

$E_k = [-\hat{e}_{k-1} \ -\hat{e}_{k-2} \ \dots \ -\hat{e}_{k-s}]^T$, the error vector

W_k is a white noise sequence

s is the order of the noise filter

To obtain F , the stochastic approximation method of the form proposed by Kwantý [19] is employed, i.e.

$$\hat{F}_{k+1} = \hat{F}_k - \frac{\gamma}{k+1} \frac{[F_k^T E_k - \hat{e}_k] E_k}{||E_k||^2} \quad (4.33)$$

where γ is a positive gain constant chosen large enough to guarantee convergence.

F_k = kth estimate of F

$$= [\hat{f}_1(k) \ \hat{f}_2(k) \ . \ . \ . \ \hat{f}_s(k)] \quad (4.34)$$

Using the estimated filter \hat{F}_k , the input and output measurements u_k and y_k respectively are filtered to obtain

$$\bar{u}_k = u_k + \sum_{i=1}^s \hat{f}_i u_{k-i} \quad (4.35)$$

$$\bar{y}_k = y_k + \sum_{i=1}^s \hat{f}_i y_{k-i}$$

These filtered quantities are used in the updating of the information matrix A_k in the pseudoinverse algorithm. Since \hat{F}_k is only an approximation to the true F , using \hat{F}_k as filter will not remove all the bias but only part of it.

The results of experiments on a second order system are presented in chapter six.

CHAPTER FIVE

ANALOG SIMULATION AND HARDWARE

5.1 Introduction

The two algorithms we have just discussed are implemented and tested in real time as applied to the identification of a one-input one-output second order system. In the present chapter, a detailed description of the simulation and hardware will be provided.

Figure 5.1 shows the general layout of the hybrid set up of the experiment. It consists of three different types of equipments. The first one is the TR-20 analog computer. The integrators, summing operational amplifiers and the potentiometers on it provide the simulation of the system to be identified as well as the pseudorandom analog input signal. The second piece of major equipment is the PDP-11/45 minicomputer. This computer is of recent design with a unibus. It allows us to address any peripheral device as convenient as any other memory locations and thus facilitate the data acquisition procedures. The memory size is 20K with both the fixed-head and moving-head disks. It is also equipped with a hardwired floating-point arithmetic processor and a real time clock. Loading and running of programs can be done easily through the system monitor which is a software package provided by the manufacturer. The analog and digital computers are coupled together by the interface panel in between. Mounted on this panel are the sampling devices and control circuits to co-ordinate the sampling process.

Control of the analog computer is a manual mechanical switch. However, the program provides a feature to co-ordinate the on-off switching of the analog computer and the starting of the identification algorithm. This is done by having the computer to repeatedly check the data buffer of one channel at the beginning. If the analog computer is off, this data buffer is zero. As soon as a certain threshold value is being detected in this data buffer, it is understood by the program that the analog computer has been switched on and it will proceed with the rest of the program. The non-zero threshold value is necessary because of the presence of a small amount of noise in the system. Experience shows that about 0.075 volt is enough.

The control of program running itself is by means of the switch register console on the computer and the keyboard. Output device can either be the cathode ray screen or the teletype printer. Since each device works on a different speed, precise timing is necessary. This can be accomplished by utilizing the priority interrupt structure of the digital computer.

For experimental purposes, an external noise generator is installed to provide noise at different power levels. The noise is artificially introduced into the output terminal of the system. For all practical purposes, we can assume that this noise source gives us white gaussian noise with a very wide spectrum.

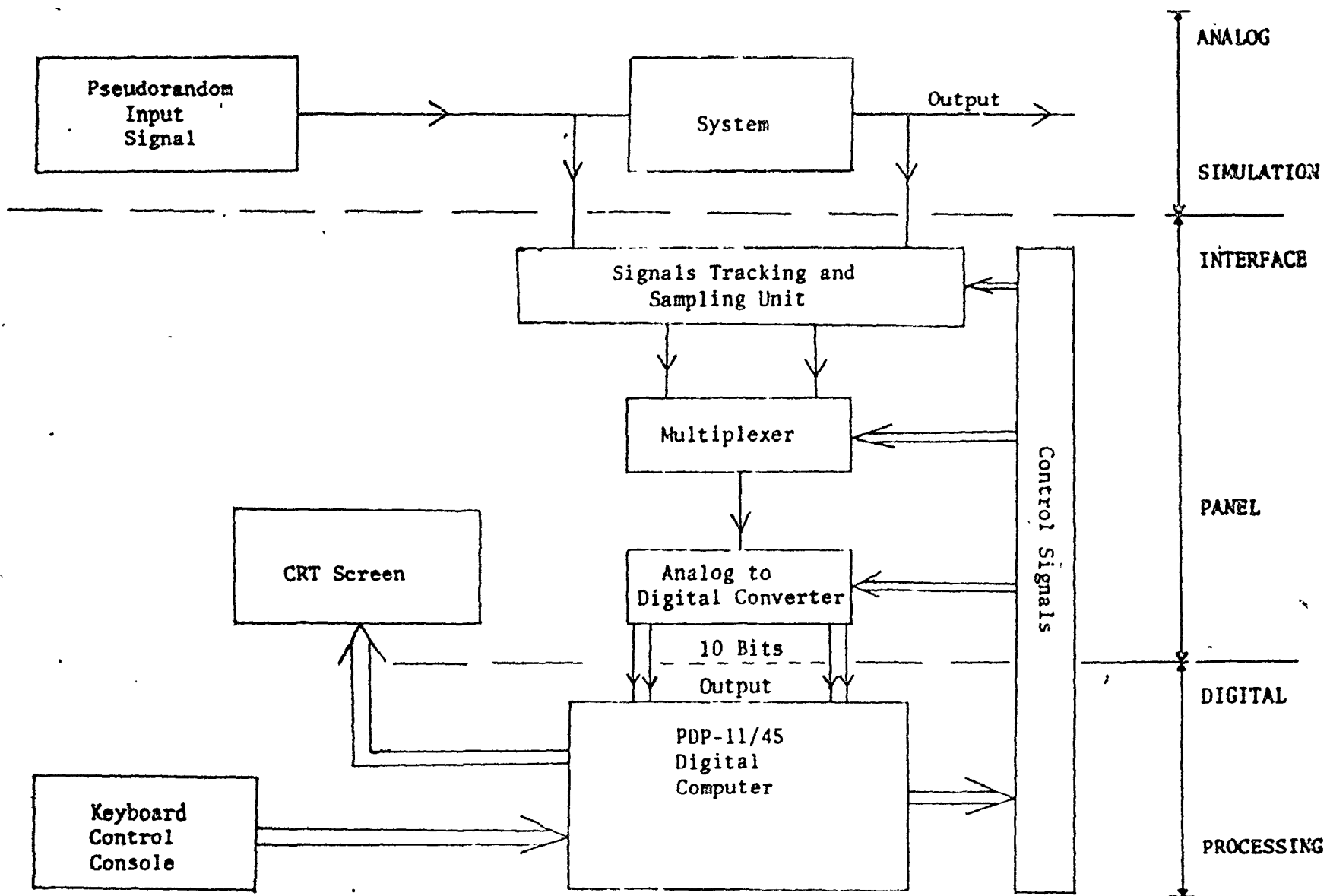


Figure 5.1 General Schematic of Experiment

5.2 Interface Hardware

The interfacing panel has been designed to handle two channels of input signals. The circuit diagram is shown in figure 5.2a.

The two channels of incoming signals are connected to a signal sampling unit which consists of two sample-and-hold modules so that simultaneous sampling of both channels is possible. This is necessary to obtain corresponding input and output measurements at the same time instant.

There is only one analog-to-digital converter capable of converting one voltage at a time. The multiplexer is situated in front of it to act as a switch. It can be programmed in such a way that different channels are presented to the input of the analog-to-digital converter individually in a pre-determined order. The binary coded digital outputs from the analog-to-digital converter are fed into a data buffer register which can be directly accessed by the digital computer to facilitate a data transfer into the core memory.

The sequence of actions of the sampling process are co-ordinated by programmed control signals together with hardwired control logic circuits. The circuit diagram of the control logic is shown in figure 5.2b.

The control bits from the channel selector register are gated to control the switching of the multiplexer through two flip-flops. The time delay circuit is used to delay the conversion trigger pulse so that the analog-to-digital converter is triggered to start the conversion just after the selected channel has been switched to the input of the converter by the

multiplexer. The delay introduced in the path travelled by the bus ready pulse is about fifty nanoseconds. The required delay for the trigger is then equal to this length of time plus the time for the analog voltage to settle in the multiplex output.

When the conversion is complete, it is signalled to the computer through a change in the logic level of the end-of-conversion output in the analog-to-digital converter. Now, the computer can transfer the data into core. After reading in the data, the computer is ready for another cycle of action.

The control of sampling frequency is done by setting an external frequency control switch on the interface panel. A detailed circuit diagram of the clock frequency generator is shown in figure 5.2c. The different frequencies are generated by allowing the clock pulses originated from the crystal clock inside the computer to pass through a different number of decade counters depending on the setting of the switch. There are a total of fifteen choices. At the lower end, we can either select one, two or five hertz. By bypassing one decade counter, we can generate pulses ten times faster. There are altogether four decade counters so that we can step up the above frequencies by four folds. However, for the purpose of process identification, we seldomly need such high frequencies. In the program written, the user can step down the sampling frequency set on the switch by any integral number of times by entering an integer constant from the keyboard. The program will make use of the integer supplied to determine the number of times it will loop through a delay loop in it.

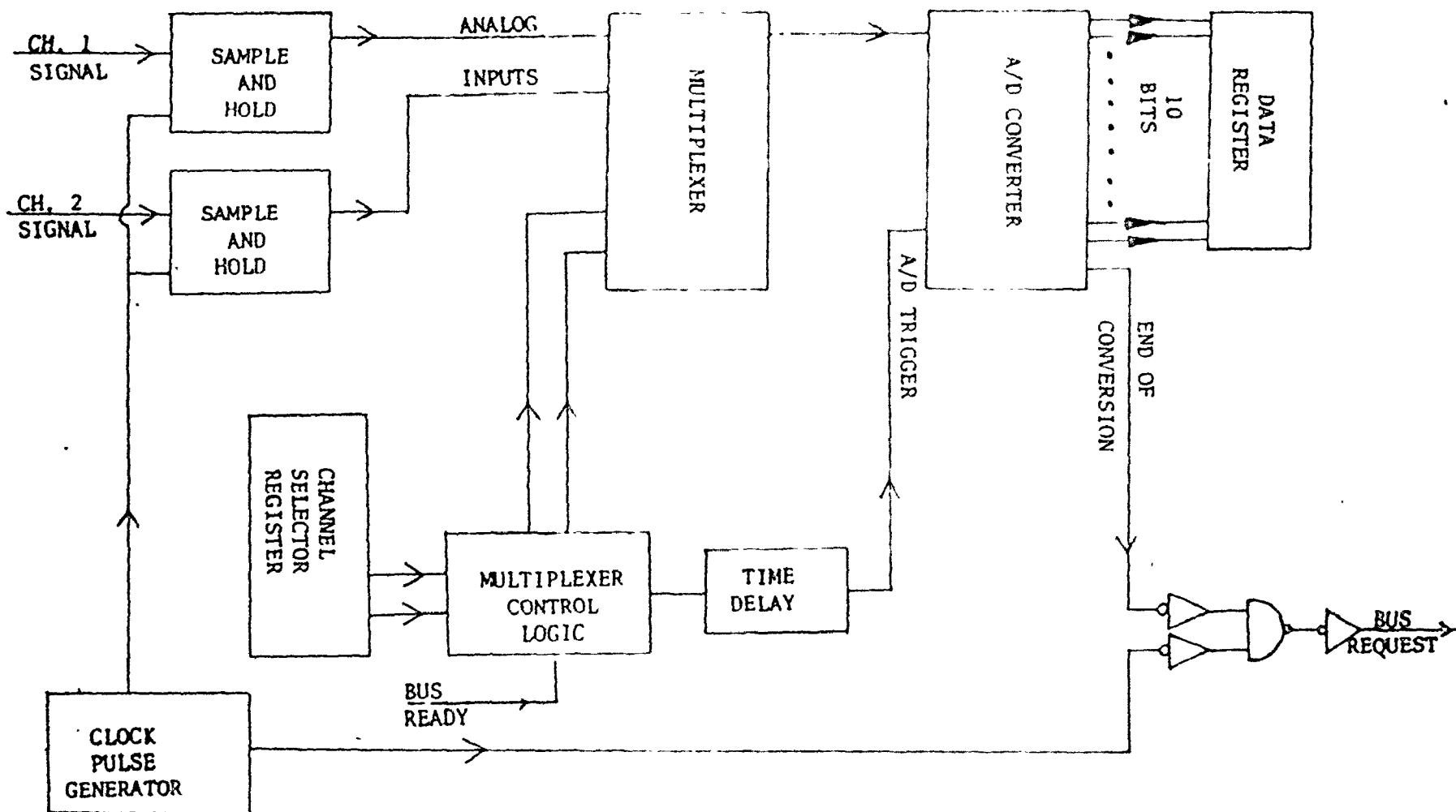


Figure 5.2a Circuit Diagram of Interface Panel

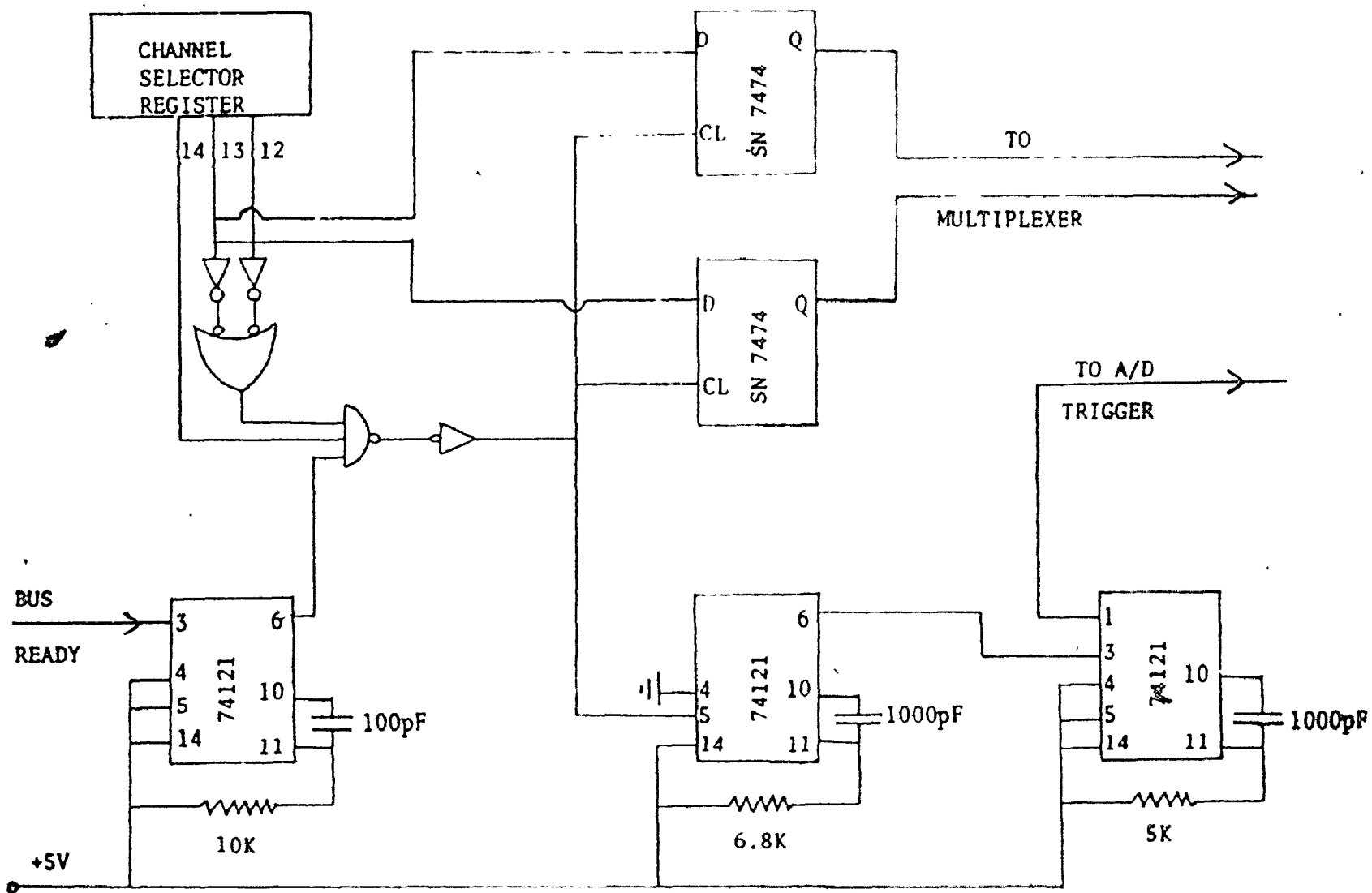


Figure 5.2b Circuit Diagram of Control Logic

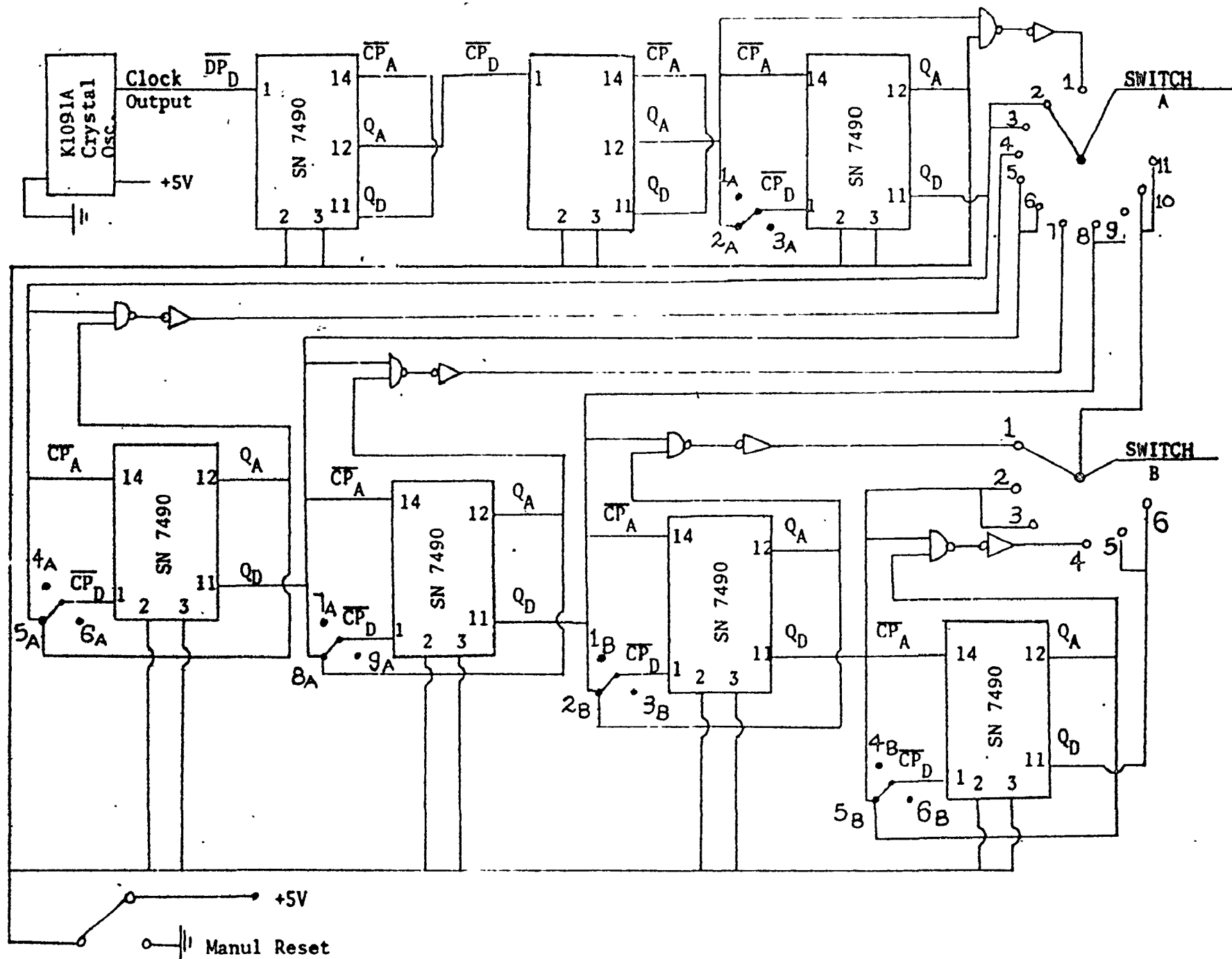


Figure 5.2c Circuit Diagram of Clock Pulse Generator

Referring back to figure 5.2a, the clock pulse is gated with the end-of-conversion output and is also fed into the sample-and-hold modules. The former connection is used for producing a bus request which is primarily responsible for the generation of an interrupt. The latter connection is used to switch the sample-and-hold modules to either one of two different modes of operation depending on the need. In the track mode, the sample-and-hold will be tracking the voltage level of the input signal. The second mode is the hold mode during which the voltage is being held at the level just before the switching signal comes in. Since both sample-and-hold modules are wired together, they go to the hold mode at the same time and hence simultaneous sampling of both channels. These samples will eventually be transmitted to the analog-to-digital converter for conversion. The digital outputs are received by a data buffer for final transfer to the computer.

Having briefly reviewed the functions of each piece of hardware in the interface panel, we shall describe the sequence of operations in the next section.

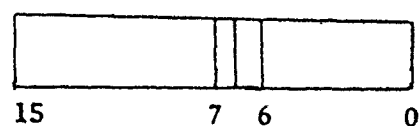
53 Operating Sequence

The whole interrupt sequence is effected by three sources of control working together, namely the external clock, the control logic circuits and the software.

There are three specialized registers in the computer central to the whole operation (see figure 5.3). The first one is the control-and-status register (ADCSR) at location 177520. Bit six is the interrupt enable bit and has to be set by the program to initiate the interrupt

sequence. Bit seven is the bus request bit set by the external circuits. When it is set, an interrupt will be generated. Depending on the priority of this interrupt request and that of the task being processed at that time, the computer will determine whether to honour the request immediately or wait until the higher priority jobs have been completed. The priority of the interrupt request is of course chosen by the programmer. The second register is the channel selector register (CHNSLR) at location 177522. Since there are two channels, we need a two-bit binary code to represent them. Bits twelve and thirteen of this register are used for this purpose. They decide which channel is being selected for conversion. The third register is the data buffer register (DATBUF). This is simply a data logging register serving as a temporary storage for the outputs of the analog-to-digital converter.

A timing diagram is shown in figure 5.4. There are six curves in the figure. The clock pulse curve determines the mode of the sample-and-hold modules. They are in track mode when the clock pulse is HI and in hold mode when the clock pulse is LO. The end-of-conversion (EOC) curve is normally at LO level except when conversion is taking place. The bus request curve is formed by ANDing the logical compliments of the first two curves. If the interrupt enable bit in the ADCSR is set, an interrupt will be generated whenever there is a positive logical transition from LO to HI in the bus request curve. In this situation, we have the sample-and-hold modules holding the signals and the previous conversion has been completed. Priority permitting and depending on the contents of a two-word interrupt vector at locations 110 and 112, the computer will honour the request by

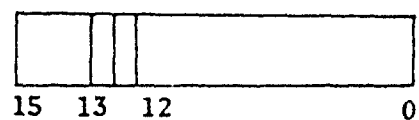


Control and Status Register (ADCSR)

Address = 177520

Bit 6 -- set for interrupt enable

Bit 7 -- set for bus request

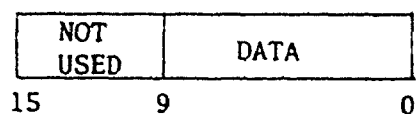


Channel Selector Register (CHNSLR)

Address = 177522

Bit 12 -- channel 1 selected if set

Bit 13 -- channel 2 selected if set



Data Register (DATREG)

Address = 177524

10 bits two's complement format

Figure 5.3 Special Registers

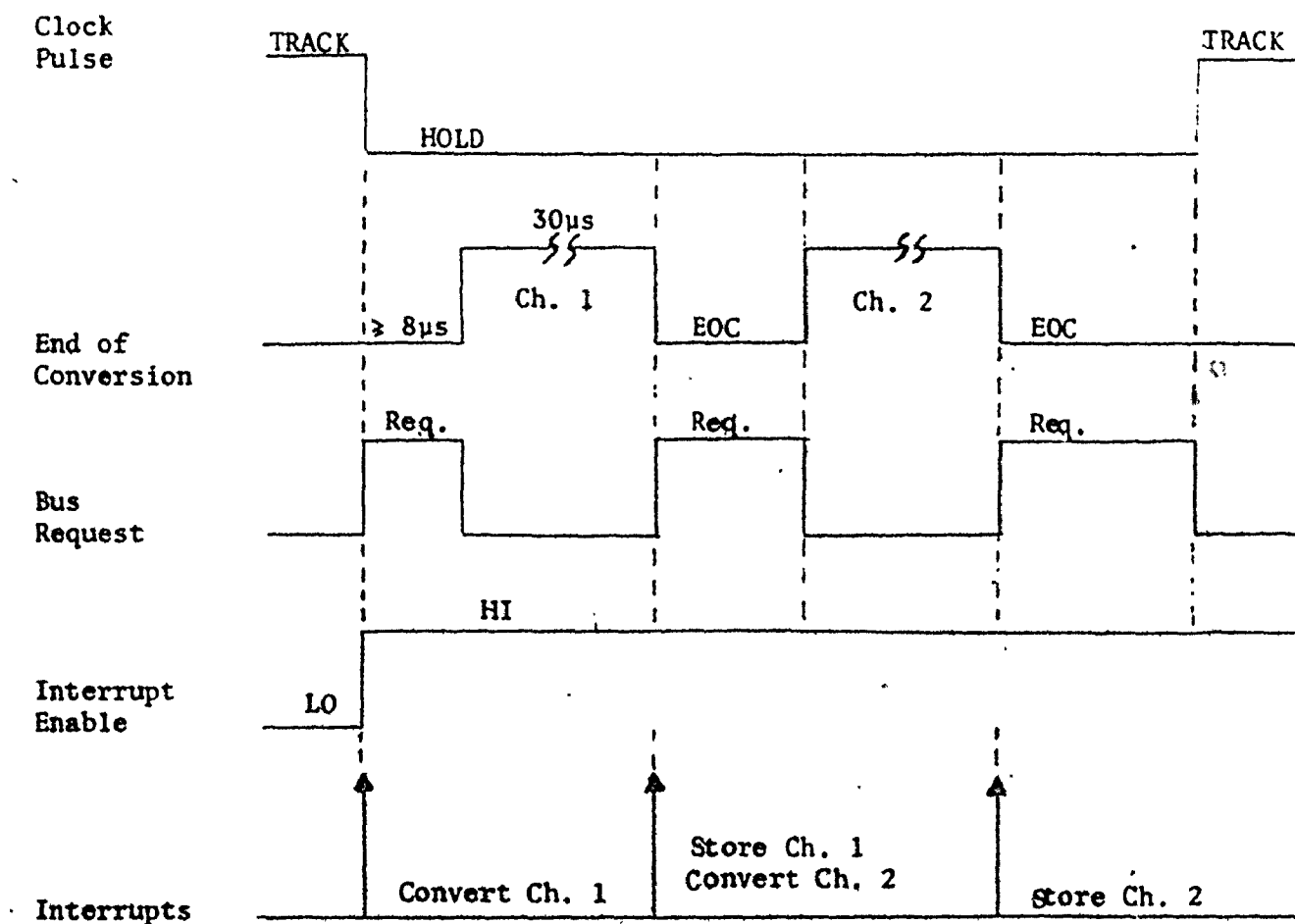


Figure 5.4 Timing Diagram

executing the subroutine pointed to by the first word of the vector with a priority according to the second word of the vector. The type of task performed is under program control. The program can change the contents of the vector to assign different tasks to different interrupt requests. As indicated in figure 5.4, the first interrupt service subroutine commands the interface panel to convert the first channel. The next one is for storing the data of the first channel in an assigned location in core followed by a command to convert the next channel. Finally, the third one is to store the data from channel two.

After servicing both channels, the cycle can be made to repeat itself for any desired length of time by simply keeping the interrupt enable bit in ADCSR set. If no further sampling is desired, the interrupt enable bit is cleared to inhibit further action.

5.4 Pseudorandom Input Signal

In section 3.4, we have noted that in order to make the assumption about the rank of the information matrix, the input signal has to satisfy one of several conditions.— One such condition is to excite the system by a random signal. In practice, true white noise signal cannot be realized physically. However, we can synthesize pseudorandom signals that would still satisfy our purpose. What we really need is a random sequence during the finite time interval when the identification process is taking place. We have at least two convenient methods at our disposal. The first method is to use a pseudorandom binary sequence (PRBS) generated from switch registers.

In our present study, we use the technique of summing many sinusoids with randomly selected phase angles to produce a pseudorandom signal. The use of sinusoids for this purpose has the important advantage that we can obtain whatever spectral density we desire. This is important because in identification problems, we can excite every mode of the system dynamics by a proper choice of the component sinusoids.

The basic principle can be understood by referring to figure 5.5 which represents a certain specified spectral density distribution of signal $y(t)$. The curve is divided into $2m$ parts of equal area. We can now replace the whole spectrum by pairs of impulse functions in both positive and negative frequencies as shown in figure 5.6. The frequency of each impulse function corresponds to the centre frequency of each of the $2m$ partitions. The magnitudes of the impulses are all equal to the area of each portion they replace. The pair of positive and negative spectra constitutes one sinusoid of randomly selected phase angles. That is, the approximation of the pseudorandom signal will be

$$y(t) = \sqrt{A} \sum_{k=1}^m \sin(\omega_k t + \phi_k)$$

where

$$A = \frac{2}{\pi} \times \text{area of each partition}$$

$$\phi = \text{randomly selected phase angle}$$

If m is increased to infinity, we would obtain a truly random signal that matches the specified spectral density exactly.

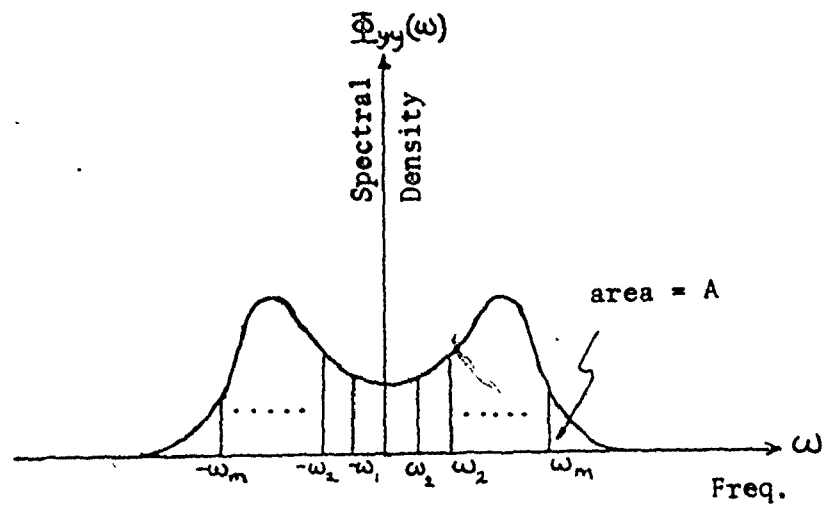


Figure 5.5 Spectral Density of $y(t)$

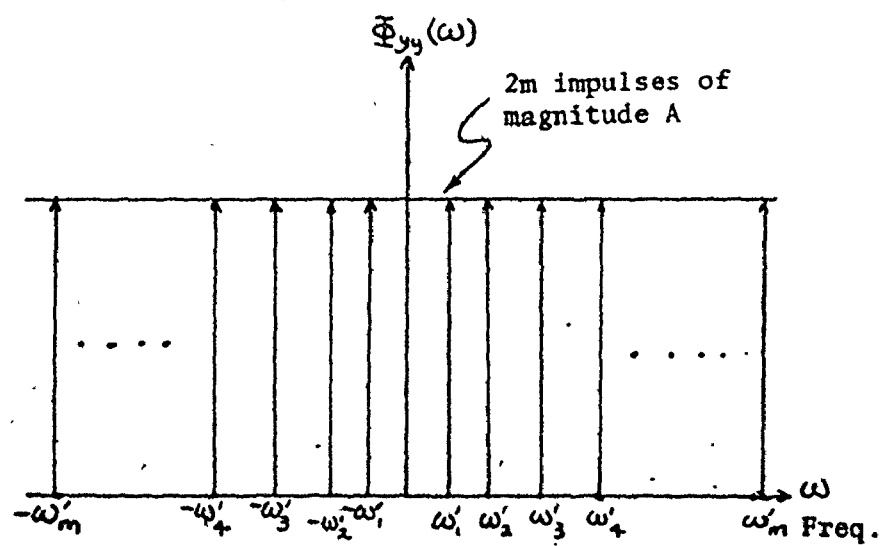


Figure 5.6 Equivalent Impulse Functions

5.5 Analog Simulations

Analog simulations of the system to be identified and the input signal we have just discussed are done on the TR-20 analog computer. A complete circuit diagram is shown in figure 5.7.

Simulation of the one-input one-output second order system is relatively straight forward requiring only two integrators and two summing amplifiers in addition to the potentiometers.

The input signal is constructed by superposing several sinusoids. Each sinusoid requires two integrators and an inverter. The choice of frequency for the construction of the pseudorandom signal must be such that they are not integrally related. This is to avoid pattern repetition by eliminating the presence of subharmonics. Fortunately, this can be easily satisfied when we use analog simulation.

Note that even though the phase angles are randomly selected, these sinusoids become deterministic signals once they are fixed. The resulting signal using a finite number of these components is therefore also deterministic and has a certain finite repetition frequency. This frequency is equal to the least common multiple of all the component frequencies and can be made small by proper adjustment. Again, using analog simulation, it poses no serious problem. Experience shows that at least five sinusoids are needed for our purpose.

The fact that we have a deterministic input signal is in effect an advantage in the experimental work because this signal is also repeatable.

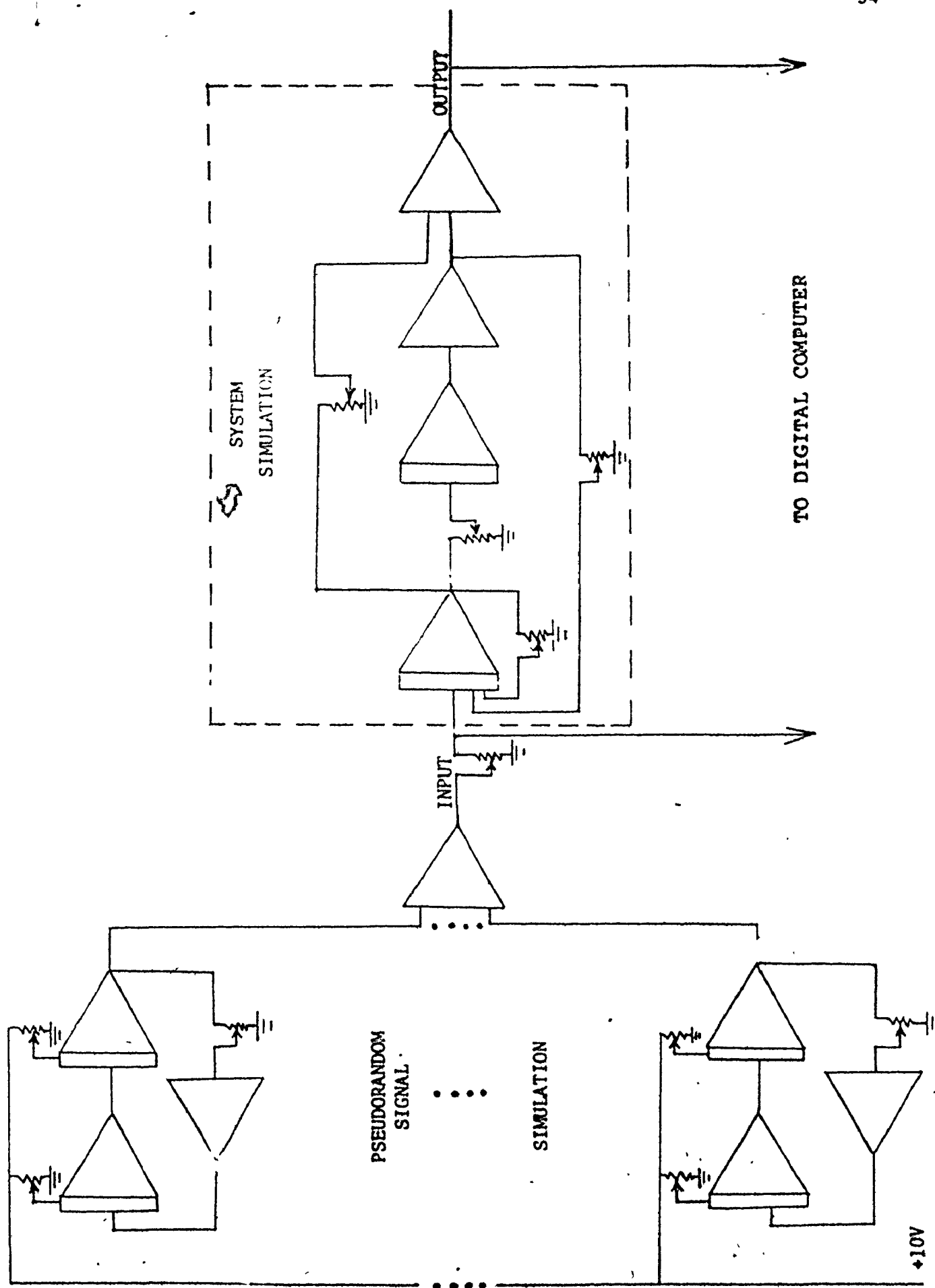


Figure 5.7 Analog Simulation

A repeatable input signal gives the same output signal every time. We can make comparisons of effects due to different noise levels. We can also compare different algorithms under essentially identical conditions. Should this input signal be truly random, we would have to make many runs of the same test in order to arrive at a statistical result.

5.6 Error Analysis

There are two main sources of errors in the system we have just described. Firstly, there are the random disturbances within the system that are completely unpredictable. This kind of disturbances are usually assumed to be white gaussian and are difficult to assess.

The second source of errors comes from the non-ideal components of the instruments used in the experiments. They are of systematic in nature. The following is a brief summary of these systematic errors:-

(A) Analog unit (non-ideal operational amplifiers)

- (i) drift
- (ii) zero off-set
- (iii) phase shifts

(B) Multiplexer

- (i) zero off-set
- (ii) non-infinite backward resistance
- (iii) non-zero forward resistance
- (iv) cross-talk due to imperfect isolation between channels

(C) Sampler

- (i) finite aperture time
- (ii) time delay due to finite tracking time and switching
- (iii) uncertainty in synchronization of simultaneous sampling
- (iv) due to finite word length

Temperature dependence of many electronic components also may introduce errors. However, the electronic components available nowadays are quite reliable and a 0.01 percent full scale deviation can usually be obtained. This is not at all severe in our present application. Since a 10-bit converter is used, it is accurate up to about ± 0.01 volt. Again, it will not cause severe degradation in the results.

CHAPTER SIX

SOFTWARE AND RESULTS

6.1 Introduction

This is a user-oriented program designed for convenient application to identify a second order one-input one-output system. It can easily be modified for higher order systems. The main bulk of the program is written in the MACRO-11 assembly language for the PDP-11/45 minicomputer. Compared with compiler language programs, the assembly language programs have the advantage of using less core memory and require less computation time since a lot of overheads can be eliminated. There is however one FORTRAN subroutine. This is being used for conducting the initial interactive dialog to obtain some essential data. It does not affect the efficiency of the identification algorithms because the dialog is being carried out at the beginning of the program before sampling is started. But it offers the convenience of flexible formats for the data to be read in from the keyboard.

As far as the size of the program is concerned, the assembly language portion of the program which consists of all computational and input/output aspects of both algorithms requires only 7632 words of core. The FORTRAN subroutine requires 1722 words of core. There is one common data block between the assembly language program and the FORTRAN subroutine occupying only 24 words of core. Modification of this program for higher order systems does not significantly change the numbers just quoted. If the output device is not fast enough to empty the output buffer for the intermediate estimates,

storage area has to be provided for storing all these intermediate results. The size of this temporary storage depends on the number of samples we want to take. Two words of storage are necessary to store one parameter value for each iteration.

6.2 Organization of the Program

A flowchart of the program is shown in figure 6.1 giving a general picture of its organization. The program is loaded into core by the system monitor in the usual manner. When it starts to run, it will first jump to the subroutine that conducts the initial dialog with the user.

During the interactive dialog, the user is asked to select either one of two algorithms i.e. the algorithm using matrix pseudoinverse only and the algorithm with filter. In the case of the second algorithm, the user has to supply a gain factor for use in the iterative stochastic approximation formula. He also has to specify when the filtering should begin. Since a third order filter is used and thirteen previous error terms are needed for updating the filter parameters, the filtering can only start after at least fourteen samples have been taken. Should the user direct the filtering to start at a even later time, the program would still update the noise filter after fourteen samples so that a more accurate filter would be available at that time.

We can also specify the maximum number of samples to be taken. However, the user reserves the right to abort the program any time during run time. This is done simply by raising bit 0 in the switch register console on the computer. The program checks this bit every time it enters a service

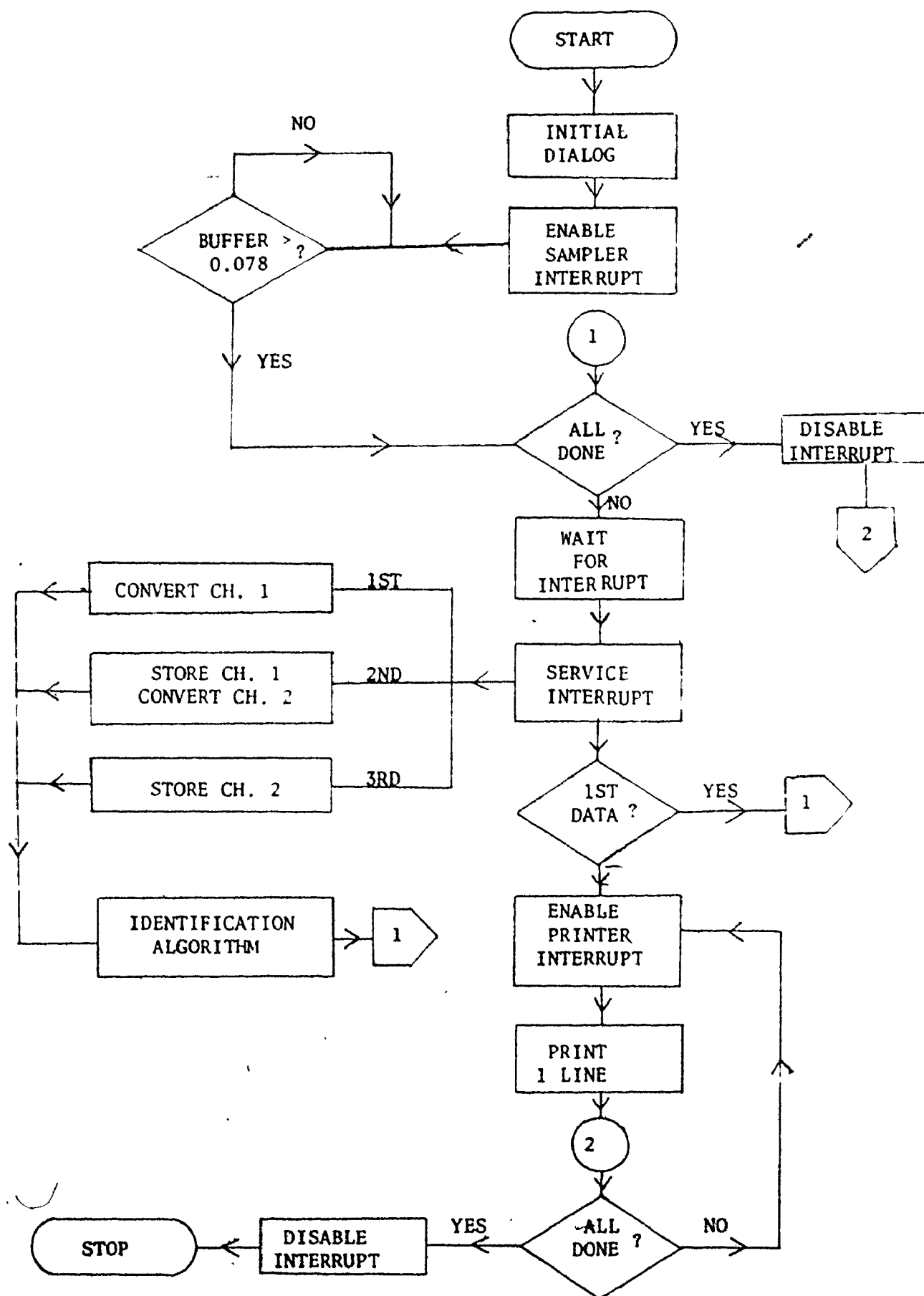


Figure 6.1 Flowchart of Program

subroutine. It would immediately clear all interrupts and then exit from the program to return to the system monitor. From thereon, the user can restart the program.

As we have noted in chapter five, the lowest sampling frequency available on the selector switch is one hertz. Here, we can enter an integer that will be used by the program as a multiplying factor of the sampling period. We can therefore greatly increase the number of choices in sampling frequency.

After completing the initial dialog, it will jump into a loop waiting for the analog computer to be switched on. The identification will start after the analog computer has been started.

Results of each iteration are printed out as soon as they are ready. A sample print-out of a typical run can be found on figure 6.2. Because of the different speeds of the many devices and the need to sample at equal intervals, priorities have been assigned to the different interrupt driven service subroutines. The sampling service subroutines have the highest priority (priority 5) because samples need to be taken at precise instances. The data acquired through sampling also need to be transferred into core from the data buffer before the next data comes in. By arranging the priority of the printing interrupt (priority 4) below that of the sampling and analysis but above that of the processor (priority 3), they can swap control of the computer until the prescribed maximum number of iterations has been reached without interfering with each other.

Figure 6.3 tabulates the functions of the major subroutines together

DO YOU WANT FILTERING ?
IF YES,TYPE 1; IF NO,TYPE 0

1
WHEN DO YOU WANT FILTERING TO START ? [13]

050
ENTER THE GAIN TERM FOR STOCHASTIC APPROXIMATION. [F5.1]

1.0
ENTER THE MULT. FACTOR FOR THE SAMPLING PERIOD. [12]

01
ENTER THE MAX. NO. OF SAMPLES YOU WANT. [13]

300

THANK YOU. TO START,STRIKE ANY KEY

	PHI1	PHI2	PHI3	PHI4	PHI5
1	2.22E-01	0.00E-00	0.00E-00	0.00E-00	0.00E-00
2	2.21E-01	5.89E-02	0.00E-00	-1.33E 00	0.00E-00
3	2.21E-01	5.90E-02	-1.61E-01	-1.33E 00	4.48E-01
4	2.21E-01	5.90E-02	-1.62E-01	-1.33E 00	4.52E-01
5	2.21E-01	5.90E-02	-1.61E-01	-1.33E 00	4.51E-01

Figure 6.2 Sample Print-out of the Results

Figure 6.3 Table of Subroutines

Subroutine Name	Functions and calling conventions
START	Main Program.
DIALOG	Obtains information from user through a series of questions.
ANAL	Dispatch subroutine for data analysis, stores the resulting estimations into buffers.
METHOD	Minimum norm and least squares algorithm for pseudoinvers.
STAPR	Stochastic approximation.
FILTER	Filtering of measurements.
CONVSN	Conversion of A/D outputs into floating point format. Calling convention: MOV DATA, R3 JSR R5, CONVSN Result is on top two words of the stack
CLOCK	Enable the sampling interrupt.
AD	Gives command to convert channel 1; interrupt driven.
STR1	Stores data of channel 1; gives command to convert channel 2, interrupt driven.
STR2	Stores data of channel 2; jumps to subroutine CONVSN: interrupt driven.
DELAY	Dispatch subroutine for analysis of data; stores resulting estimates in buffers.
DELAY	Modifies sampling period by an integer multiplicative factor $k \geq 1$.
PRINT	Enables the printing interrupt; monitors the progress of sampling, data analysis and printing.
IO	ASCII conversion of results before transferring to the printing buffers.
IOF	Generates a 3-digit ASCII coded line counter in ascending order.
PRN	Printing subroutine; interrupt driven.

Figure 6.3 cont'd.

Mathematical Subroutines

MULFP	Floating point multiplication, addition or subtraction
ADDFP	of A and B.
DIVFP	Calling convention:
SUBFP	JSR R5, XXXFP
	.WORD A, B, C
	Result is in C
MMUL	Matrix multiplication, addition or subtraction of A and B.
MADD	Calling convention:
MSUB	JSR R5, MXXX
	.WORD A, B, C
	.WORD ROW, COL
	Result is in C.
DSC	Matrix division or multiplication by a scaler.
MSC	Calling conventions:
	JSR R5, XSC
	.WORD A, B, SC
	.WORD RWO, COL
	Result is in B.
MTRN	Matrix transposition.
	Calling convention:
	JSR R5, MTRN
	.WORD A, AT
	.WORD ROW, COL
	Result is in AT.
SHIFT	Shifts all elements of vector A by one position downwards.
	Calling convention:
	JSR R5, SHIFT
	.WORDS A, N

with their calling conventions whenever they are appropriate. This program is adaptable to computers without a hardwired floating point processor without making any changes. Macro definitions are liberally used in the mathematical subroutines to facilitate easy checking. The complete heavily commented program listing is in the appendix.

6.3 Results

The results are based on experiments identifying the second order one-input one-output system given by

$$H(s) = \frac{0.04+0.28s}{0.04+0.4s+s^2} \quad (6.1)$$

According to Sinha [21], the sampled-data equivalent of equation (5.1) can be obtained by the bilinear transformation

$$s = \frac{2(1-z^{-1})}{T(1+z^{-1})} \quad (6.2)$$

subject to the condition that

$$p_k T \leq 0.5 \quad (6.3)$$

where T = sampling period

p_k = system poles.

The system represented by equation (6.1) has poles

$$p_1 = p_2 = 0.2 \quad (6.4)$$

Therefore, the transformation given by (6.2) is valid if $T=2$ seconds is used.

Applying (6.2) to (6.1), we have

$$H(z) = \frac{0.222 + 0.0556z^{-1} - 0.167z^{-2}}{1 - 1.333z^{-1} + 0.444z^{-2}} \quad (6.5)$$

However, due to the presence of a small amount of disturbance in the system even when no external noise is added, the experimental results of the coefficients of equation (5.5) are found to be

$$\phi^T = [0.221 \quad 0.059 \quad 0.162 \quad -1.33 \quad 0.452] \quad (6.6)$$

Values in equation (6.6) are being used as a reference for subsequent comparisons.

Whenever it is appropriate, experimental results from both algorithms are presented together. Both algorithms are tested under different amount of externally introduced white noise into the output of the simulated system. For the second algorithm using combined matrix pseudoinverse and stochastic approximation, a third order noise filter has been used.

Figure 6.4 summarizes the estimates and the resulting normalized errors for both algorithms after three hundred iterations. We can conclude from these results that the error of estimation is directly dependent on the level of noise present. It is also observed that the second algorithm gives better estimates in all instances. The extent of improvement, however, varies.

Sample plots are shown in figures 6.5 and 6.6 displaying the behaviour of the error in the estimates during the course of identification for the pseudoinverse algorithm and the one with filtering respectively. Note that in figure 6.6, filtering starts only after fifty iterations in order to avoid the initial region with large fluctuations in the estimation.

Despite the good convergence properties clearly exhibited, there always exists a bias in the estimation as expected because the residuals are correlated. Considerable amount of the bias is successfully removed by the filtering method.

It is of interest to recall that parameter estimation problems can be considered as optimization problems. In our case, the objective for minimization is the residual error sequence. This is our criterion to define the "goodness" of estimation. There are other criteria that can be used. For example, the time domain errors of step response or the impulse response are both valid criteria to evaluate estimations. Since there is no direct functional relationship among different criteria, a good estimation according to one criterion does not necessarily imply that it is also good when judged by other criteria. A case in point to illustrate this fact is to compare the step response and impulse response of the tests we have performed.

The two plots figure 6.7 show the step response and impulse response respectively of the two algorithms at a noise ratio of 25%. From the table of figure 6.4, we notice the big difference in the values of the parameters. We also notice that, the first algorithm has a normalised error of 88.72% compared with the much smaller 11.38% in the second algorithm. The difference

of almost seven hundred percents is also show graphically in figure 6.6b. Despite all these dramatic differences, a reference to figure 6.7 shows that their responses are not too far apart considering what we have seen from the difference in parameter values. The table in figure 6.8 tabulates the corresponding mean square errors in step and impulse response of all the tests we have performed. There is no direct correlation between this set of values and the normalised errors. While filtering unfailingly reduces the parameter error, no such conclusion can be drawn about the errors in step response and impulse response.

As far as computation time is concerned, no more than 0.05 second per iteration is required for the matrix pseudoinverse algorithm. Twice as much time is required for the second algorithm i.e. 0.1 second.

	NOISE RATIO (%)	ESTIMATION					NORMALIZED ERROR
		$\hat{\phi}_1$	$\hat{\phi}_2$	$\hat{\phi}_3$	$\hat{\phi}_4$	$\hat{\phi}_5$	
PSEUDOINVERSE	5	0.2188	0.2596	0.04155	-0.4156	-0.07847	0.5487
	10	0.2176	0.2928	0.07486	-0.2561	-0.1682	0.8035
	25	0.2137	0.3011	0.08354	-0.1934	-0.1879	0.8872
	50	0.2076	0.2992	0.08611	-0.1603	-0.1546	0.9044
	75	0.2025	0.2976	0.08969	-0.1305	-0.1081	0.9130
	100	0.1982	0.2961	0.09082	-0.1056	-0.0654	0.9199
WITH STOCHASTIC APPROX.	5	0.2126	0.1407	-0.07118	-0.9829	0.2575	0.08469
	10	0.1835	0.1252	-0.04103	-0.9497	0.2355	0.1035
	25	0.2167	0.1960	-0.06142	-0.9101	0.2858	0.1137
	50	0.1878	0.1088	-0.04592	-0.5546	-0.05058	0.4255
	75	0.1973	0.06805	0.03884	-0.5079	0.02452	0.4387
	100	0.1841	0.01295	0.02091	-0.4835	0.005042	0.4716

Figure 6.4 Results in Parameter Estimation and the Normalized Errors

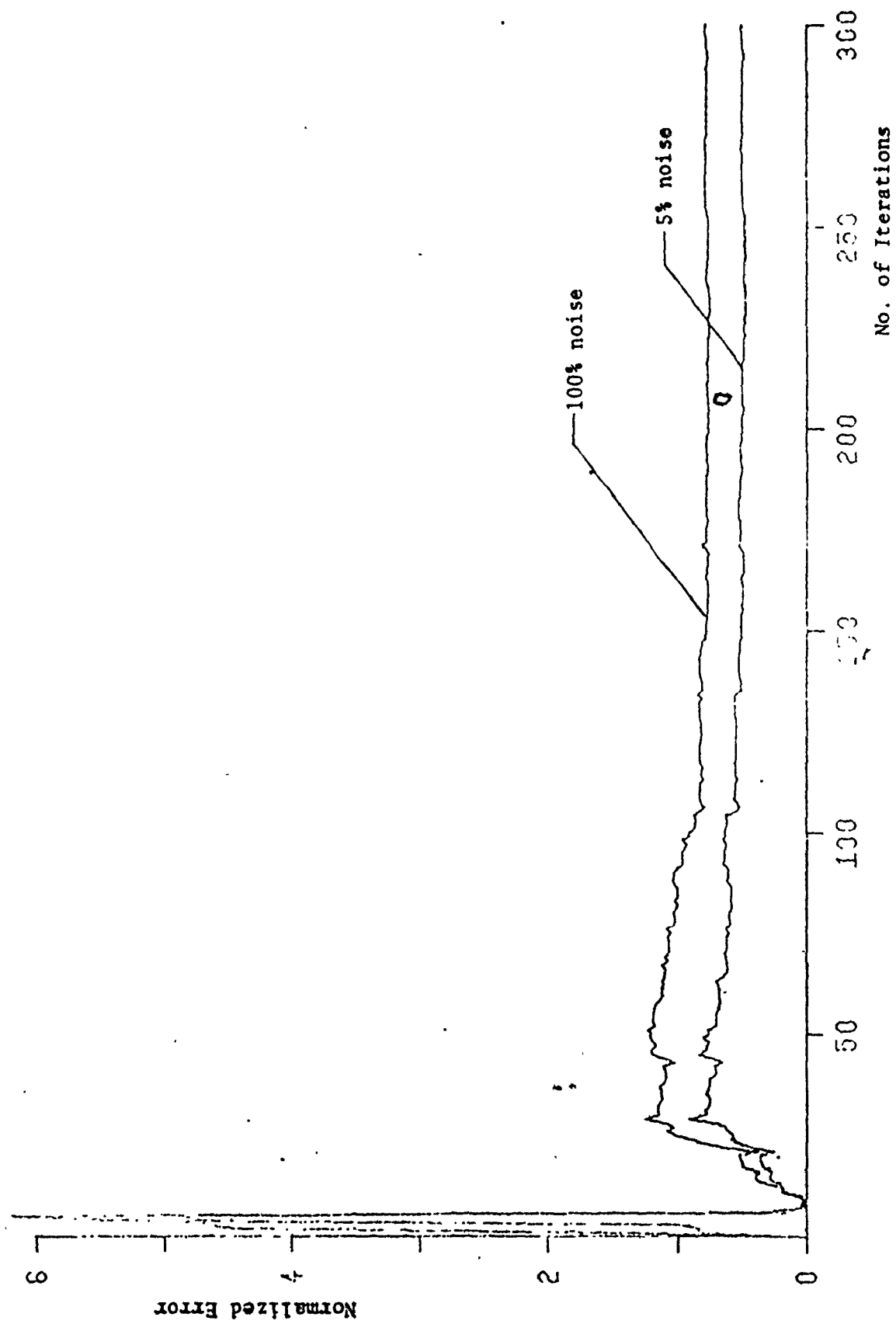


Figure 6.5 Pseudoinverse Algorithm

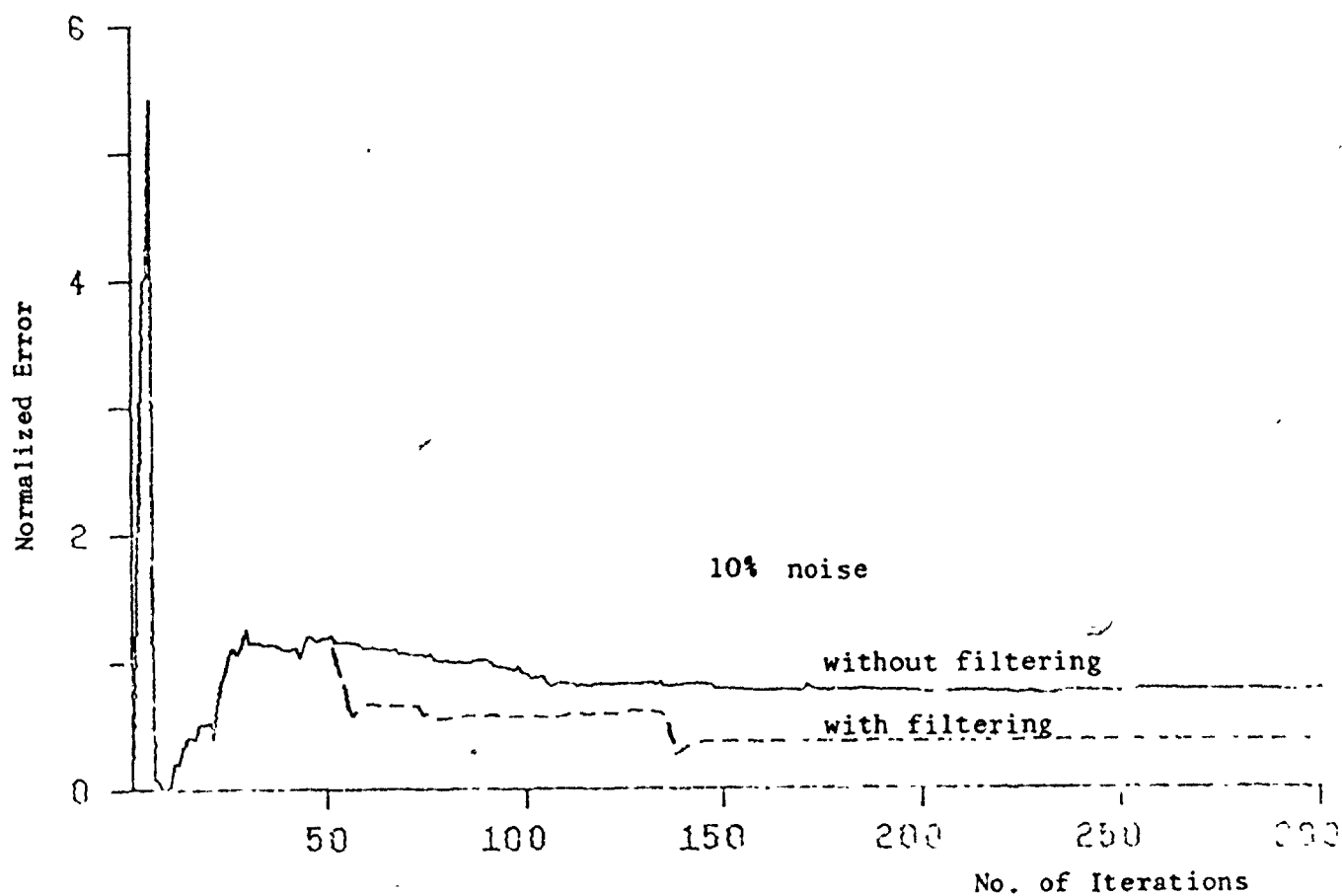
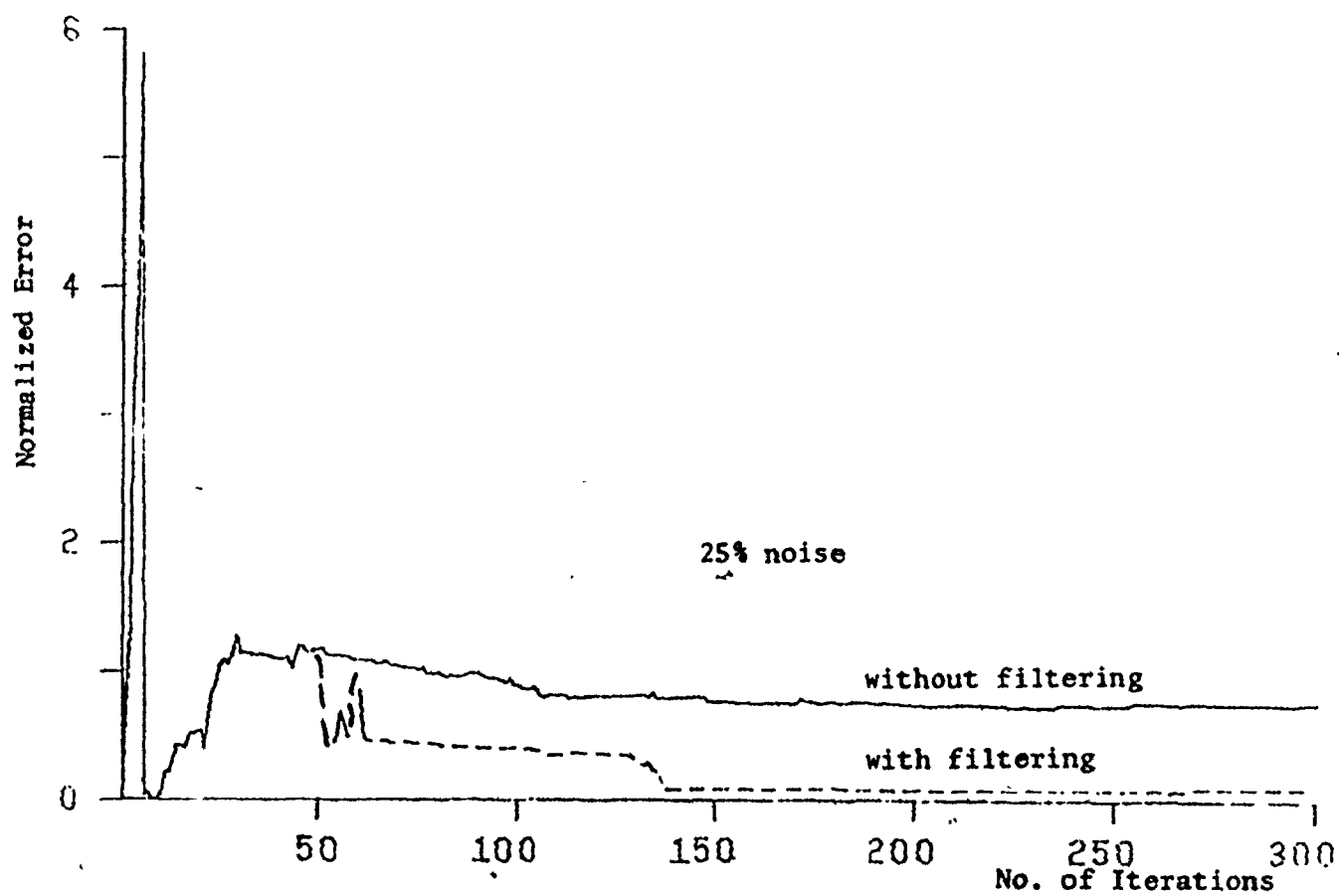


Figure 6.6a Stochastic Approx. Filtering



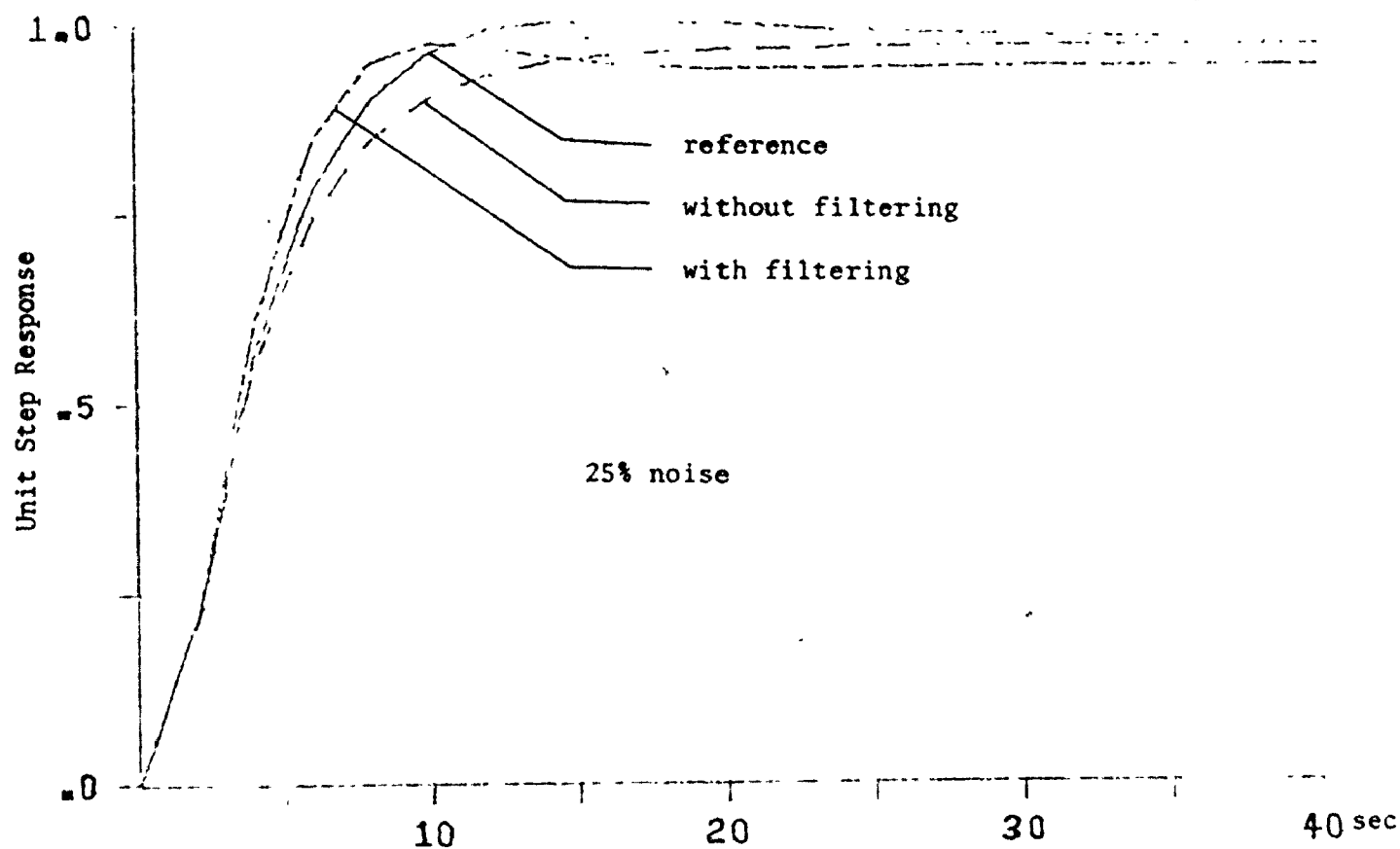


Figure 6.7a Comparison of Step Responses

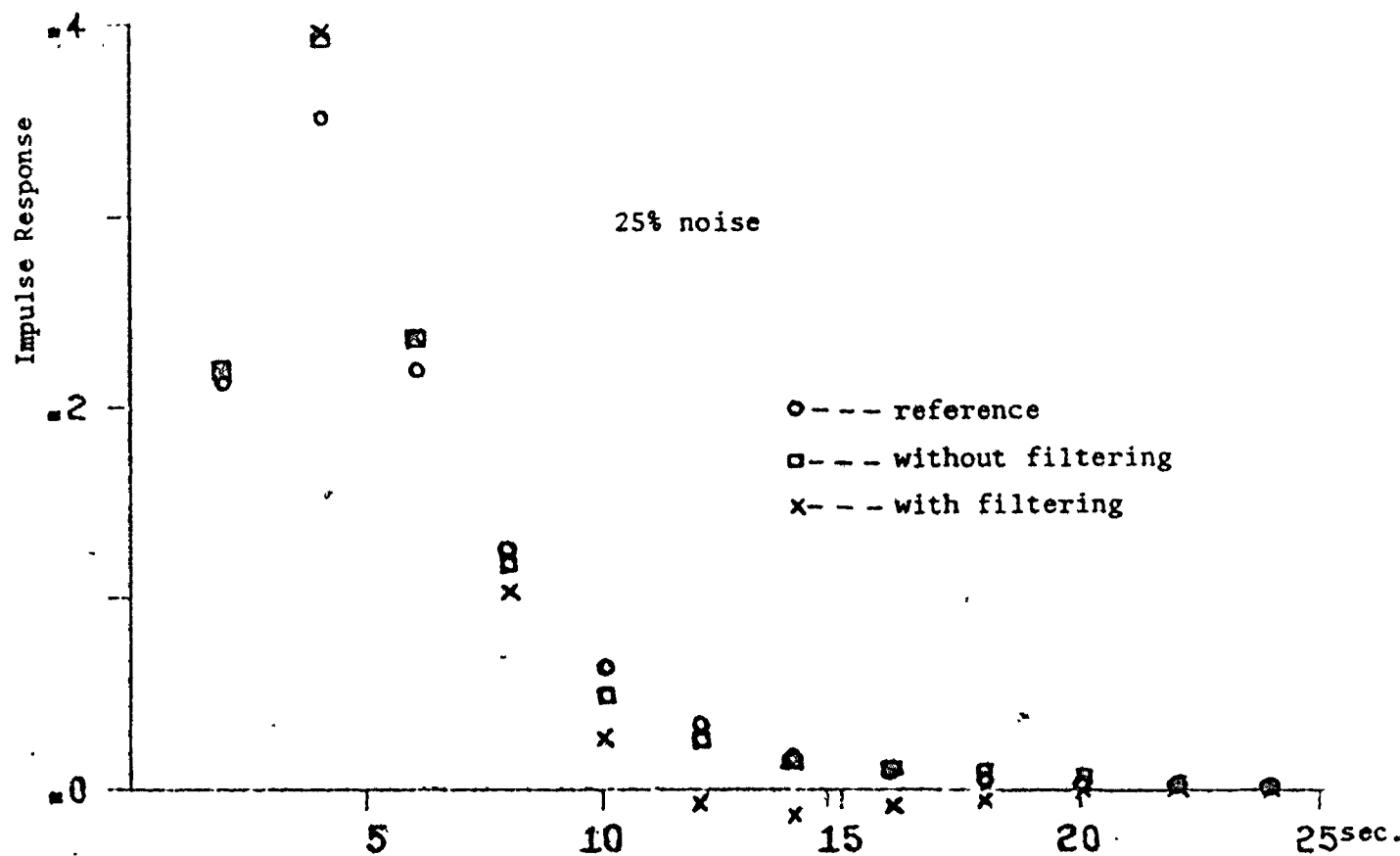


Figure 6.7b Comparison of Impulse Responses

NOISE RATIO (%)		MEAN SQUARE ERROR	
		STEP RESPONSE ($\times 10^{-2}$)	IMPULSE RESPONSE ($\times 10^{-4}$)
PSEUDO INVERSE	5	0.1122	0.2367
	10	0.07293	0.2759
	25	0.1108	0.5449
	50	1.3490	2.1173
	75	3.8398	5.0463
	100	6.6374	8.4712
WITH STOCHASTIC APPROX.	5	0.1186	0.2367
	10	0.4288	0.2318
	25	0.2245	3.0333
	50	11.7404	20.811
	75	14.086	24.0609
	100	28.6127	46.859

Figure 6.8 Mean Square Errors in Step Responses and Impulse Responses

CHAPTER SEVEN

CONCLUSIONS

The theories presented in this thesis result in two algorithms for on-line system identification. The first one is the matrix pseudoinverse algorithm. A fundamental property of this method is that the estimates are optimal in the sense that the residual error is minimized.

The second algorithm is an extension of the first one aiming at further improving the estimation by removing the bias in the estimates due to measurement noise. The mechanism employed is the introduction of a filter obtained from the properties of the noise present in the system. The estimation of the noise properties and hence the construction of the filter is itself an on-line estimation process being carried out in parallel with the pseudoinverse method. The tool to this end is the use of stochastic approximation which is also computationally simple. As a result, incorporation of the stochastic approximation into the original scheme does not result in serious degradation in efficiency.

The upgrading of the program to accommodate higher order systems can be easily done by simply expanding the data block to provide enough room as a working area. All other instructions remain unchanged. However, high order models are usually not necessary in many engineering applications. Very complex systems can sometimes be approximated by second or third order models. For example, Sinha and Bereznoi[22] have modelled the dynamics of

the nuclear reactor power generating station in Douglas Point by a second order model. This second order model which is being updated on-line is subsequently used for controller design. This example may serve as a justification for choosing only a second order example in the present study. Also of importance as a reminder, simulation of the system on the analog computer bears close resemblance to the actual situations where identification might be employed. In the industrial setting, transducers are used to make such measurements as temperature and speed in terms of electrical voltages. These voltages are being sampled just as we have done in the present study.

We can see from the results that the convergence property of both algorithms are clearly and nicely exhibited even in high noise environment. The accuracy in estimation are quite naturally deteriorating with the increase in the amount of noise. Introduction of filtering, however, greatly enhances this accuracy at the expense of a relatively small amount of computation time.

However, there remains room for further improvement. A common problem in applying stochastic approximation is the difficulty in choosing an appropriate gain factor that would best suit a particular situation. No systematic method yet exists. It would be a worthwhile research area to devise an iterative scheme by which an optimal prediction of this gain factor can be done on-line together with what we have already developed. A point of caution is in order. We have to watch out for the amount of extra computation thereby introduced. A time consuming method would destroy the major appeal of our present approach which is computational simplicity.

APPENDIX
Program Listing

```

1      .TITLE IDENT
2      .GLOBL DIALOG
3      .GLOBL $ADR,$SBR,$MLR,$DVR,$POLSH
4      000000 R0=Z0
5      000001 R1=Z1
6      000002 R2=Z2
7      000003 R3=Z3
8      000004 R4=Z4
9      000005 R5=Z5
10     000006 SP=Z6
11     000007 PC=Z7
12     ;
13     .SBTTL MACRO DEFINITIONS
14     ;
15     .MACRO MMUL A,B,C,ROW,COLA,COLB
16         JSR      R5,MMUL
17         .WORD    A,B,C
18         .WORD    ROW,COLA,COLB
19         .ENDM
20     .MACRO MADD A,B,C,ROW,COL
21         JSR      R5,MADD
22         .WORD    A,B,C
23         .WORD    ROW,COL
24         .ENDM
25     .MACRO MSUB A,B,C,ROW,COL
26         JSR      R5,MSUB
27         .WORD    A,B,C
28         .WORD    ROW,COL
29         .ENDM
30     .MACRO MSC A,B,SC,ROW,COL
31         JSR      R5,MSC
32         .WORD    A,B,SC
33         .WORD    ROW,COL
34         .ENDM
35     .MACRO DSC A,B,SC,ROW,COL
36         JSR      R5,DSC
37         .WORD    A,B,SC
38         .WORD    ROW,COL
39         .ENDM
40     .MACRO MTRN A,AT,ROW,COL
41         JSR      R5,MTRN
42         .WORD    A,AT
43         .WORD    ROW,COL
44         .ENDM
45     .MACRO SAVE RA,RB,RC,RD,RE,RF
46         MOV      RA,-(SP)
47         .IIF DF RB,      MOV      RB,-(SP)
48         .IIF DF RC,      MOV      RC,-(SP)
49         .IIF DF RD,      MOV      RD,-(SP)
50         .IIF DF RE,      MOV      RE,-(SP)
51         .IIF DF RF,      MOV      RF,-(SP)
52         .ENDM

```

```

53      .MACRO UNSAVE RA,RB,RC,RD,RE,RF
54          .IIF DF RF,      MOV      (SP)+,RF
55          .IIF DF RE,      MOV      (SP)+,RE
56          .IIF DF RD,      MOV      (SP)+,RD
57          .IIF DF RC,      MOV      (SP)+,RC
58          .IIF DF RB,      MOV      (SP)+,RB
59          MOV      (SP)+,RA
60      .ENDM
61      .MACRO POPF A
62          MOV      (SP)+,A
63          MOV      (SP)+,A+2
64      .ENDM
65      .MACRO PUSHF A
66          MOV      A+2,-(SP)
67          MOV      A,-(SP)
68      .ENDM
69      .MACRO FMUL A,B,C
70          JSR      R5,MULFP
71          .WORD    A,B,C
72      .ENDM
73      .MACRO FDIV A,B,C
74          JSR      R5,DIVFP
75          .WORD    A,B,C
76      .ENDM
77      .MACRO FADD A,B,C
78          JSR      R5,ADDFP
79          .WORD    A,B,C
80      .ENDM
81      .MACRO FSUB A,B,C
82          JSR      R5,SUBFP
83          .WORD    A,B,C
84      .ENDM
85      .MACRO MOVF A,B
86          MOV      A,B
87          MOV      A+2,B+2
88      .ENDM
89      .MACRO FSHIFT A,NUM
90          JSR      R5,SHIFT
91          .WORD    A,NUM
92      .ENDM

```

```

1      .SBTTL MAIN PROGRAM
2      ;
3      ;
4      ;
5      ;      THIS IS THE MAIN PROGRAM
6      ;
7      ;
8      ;      IT FIRST JUMPS TO SUBROUTINE DIALOG TO
9      ;      OBTAIN THE FOLLOWING PARAMETERS :-
10     ;
11     ;      IFTR-- - TO SEE IF FILTERING IS DESIRED
12     ;      IFIRST-- IF SO, WHEN SHOULD IT BEGIN
13     ;      GAIN--   AND WHAT WOULD BE THE VALUE OF
14     ;      THE GAIN FACTOR FOR STOCH. APPROX.
15     ;      IDELAY-- A PARAMETER TO MODIFY THE SAMPLING
16     ;      FREQUENCY
17     ;      MAXAD--- MAX. NO. OF SAMPLES TO BE TAKEN
18     ;
19     ;      NEXT, SUBROUTINE CLOCK INITIATES THE
20     ;      SAMPLING PROCESS AND PROCEED WITH THE
21     ;      ALGORITHM SELECTED I.E.
22     ;      ALGORITHM 1 -- MATRIX PSEUDOINVERSE
23     ;      ALGORITHM 2 -- COMBINED PSEUDOINVERSE
24     ;      AND STOCHASTIC APPROX.
25     ;
26     ;      SUBROUTINE PRINT MONITORS THE PROGRESS OF
27     ;      THE PROGRAM AND PRINTS OUT THE ESTIMATES
28     ;      WHEN THEY ARE READY
29     ;
30     ;
31     ;
32     ;
33 00300 004567 START: JSR      R5,DIALOG;      INITIAL DIALOG
34      003000G
35 00004 004567      JSR      R5,CLOCK;        STARTS SAMPLING
36      003532
37 00010 000001      WAIT
38 00012 004567      JSR      R5,PRINT;        STARTS PRINTING
39      004620
40 00016 000005      RESET
41 00020 000000      HALT
42 00022 104060 15:   EMT      60
43      ;
44      ;
45      ;
46      ;      000000 .CSECT FTRN
47      ;
48      ;      DATA OBTAINED FROM INITIALIZATION SUBROUTINE
49      ;
50 00000 000000 IFTR:   .WORD 0
51 00002 000000 IFIRST: .WORD 0
52 00004 000000 IDELAY: .WORD 0
53 00006 000000 MAXAD:  .WORD 0
54 00010 000000 GAIN:   .BLKW 2.

```

```

1      000024'.CSECT
2      .SBTTL PSEUDOINVERSE ALGORITHM
3      ;
4      ;
5      ;
6      ;      COMMON SECTION
7      ;      CAL. ATPHI,Y-ATPHI,PA,ATPA,1+ATPA
8      ;
9      000024      METHOD: MMUL      AT,PHI,DUM1,ONE,FIVE,ONE
10     00044      FSUB      Y,DUM1,DUM1
11     00056      MMUL      P,A,PA,FIVE,FIVE,ONE
12     00076      MMUL      AT,PA,DUM2,ONE,FIVE,ONE
13     00116      FADD      F.ONE,DUM2,DUM2
14     00130      026727      CMP      AN.K,#4;          >5 ITERATIONS ?
           004500
           000004
15     00136      003156      BGT      PSEUDO
16     ;
17     ;      MINIMUM NORM SECTION
18     ;
19     00140      MMUL      Q,A,QA,FIVE,FIVE,ONE
20     00160      MMUL      AT,QA,DUM3,ONE,FIVE,ONE
21     00200      DSC      QA,T5.1,DUM3,FIVE,ONE
22     ;
23     ;      UPDATE Q
24     ;
25     00216      MMUL      T5.1,QAT,T5.5,FIVE,ONE,FIVE
26     00236      MSUB      Q,T5.5,Q,FIVE,FIVE
27     ;
28     ;      UPDATE PHI
29     ;
30     00254      MSC      T5.1,T5.1,DUM1,FIVE,ONE
31     00272      MADD      PHI,T5.1,PHI,FIVE,ONE
32     ;
33     ;      UPDATE P
34     ;
35     00310      DSC      T5.5,T5.5,DUM3,FIVE,FIVE
36     00326      MSC      T5.5,T5.5,DUM2,FIVE,FIVE
37     00344      MADD      P,T5.5,P,FIVE,FIVE
38     00362      MMUL      PA,QAT,PAQAT,FIVE,ONE,FIVE
39     00402      MTRN      PAQAT,T5.5,FIVE,FIVE
40     00416      MADD      PAQAT,T5.5,T5.5,FIVE,FIVE
41     00434      DSC      T5.5,T5.5,DUM3,FIVE,FIVE
42     00452      MSUB      P,T5.5,P,FIVE,FIVE
43     00470      000167      JMP      FINSH
           000126

```

```

44      ;
45      ;
46      ;      LEAST SQUARES SECTION (K>P)
47      ;
48      ;
49      ;      UPDATE P
50      ;
51 00474 PSEUDO: MMUL    PA,PAT,T5.5,FIVE,ONE,FIVE
52 00514      DSC      T5.5,T5.5,DUM2,FIVE,FIVE
53 00532      MSUB     P,T5.5,P,FIVE,FIVE
54      ;
55      ;      UPDATE PHI
56      ;
57 00550      MSC      PA,T5.1,DUM1,FIVE,ONE
58 00566      DSC      T5.1,T5.1,DUM2,FIVE,ONE
59 00604      MADD     PHI,T5.1,PHI,FIVE,ONE
60      ;
61      ;      ALGORITHM 1 OR 2
62      ;
63 00622 022767 FINSH: CMP      #0,IFTR;      NEED FILTERING ?
        000000
        000000
64 00632 001402      BEQ      1$;      NO,SKIP
65 00632 004567      JSR      R5,STAPR;  STOCHASTIC APPRO.ALGO.
        001576
66 00636 000205 1$:   RTS      R5

```



```

1      .SBTTL SUBROUTINE UPDATE
2      ;
3      ;      UPDATE--SUBROUTINE TO UPDATE THE
4      ;      INFORM. MATRIX AS NEW SAMPLES COME IN
5      ;      IT ALSO JUMPS TO FILTERING IF NEEDED
6      ;
7      ;
8      000640      UPDATE:
9      000640 026767      CMP      IFIRST,AN.K;      START FILTERING ?
               000302
               003766
10     00646 002002      BGE      21$;      NO,SKIP
11     00650 004567      JSR      R5,FILTER;      YES
               002574
12     00654      21$:      FSHIFT  A,N3;      UPDATING OF
13     00664      FSHIFT  A+14,N2;      INFORMATION
14     00674      MOVF     U,A;      MATRIX
15     00710      MOVF     YOLD,A+14;      HERE
16     00724      MOVF     Y,YOLD
17     00740 022767      CMP      #0,YOLD;      CHECK IF
               000000
               001026
18     00746 001005      BNE      10$;      YOLD=0
19     00750 022767      CMP      #0,YOLD+2
               000000
               001020
20     00756 001413      BEQ      2$
21     00760 000407      BR       1$
22     00762 005767 10$:      TST      YOLD
               001006
23     00766 100004      BPL      1$
24     00770 042767      BIC      #100000,YOLD
               100000
               000776
25     00776 000403      BR       2$
26     01000 052767 1$:      BIS      #100000,YOLD
               100000
               000766
27     01006 000205 2$:      RTS      R5
28     ;
29     ;
30     ;      DATA BLOCK FOR BOTH ALGORITHMS
31     ;
32     ;
33     01010      A:
34     01010      AT:      .BLKW 10.
35     01034      P:      .BLKW 50.
36     01200      PA:
37     01200      PAT:     .BLKW 10.

```

```

38 01224 040200 Q:      .FLT2 1.E0
   01226 000000
39                      .BLKW 10.
40 01254 040200      .FLT2 1.E0
   01256 000000
41                      .BLKW 10.
42 01304 040200      .FLT2 1.E0
   01306 000000
43                      .BLKW 10.
44 01334 040200      .FLT2 1.E0
   01336 000000
45                      .BLKW 10.
46 01364 040200      .FLT2 1.E0
   01366 000000
47 01370              QA:
48 01370              QAT:      .BLKW 10.
49 01414              PHI:      .BLKW 10.
50 01440              PAQAT:    .BLKW 50.
51 01604              T5.5:     .BLKW 50.
52 01750              T5.1:     .BLKW 10.
53 01774              YOLD:     .BLKW 2.
54 02000              DUM1:     .BLKW 2.
55 02004              DUM2:     .BLKW 2.
56 02010              DUM3:     .BLKW 2.
57 02014 000001 ONE:      .WORD 1.
58 02016 000002 TWO:      .WORD 2
59 02020 000003 THREE:    .WORD 3
60 02022 000005 FIVE:     .WORD 5
61 02024 000002 N2:       .WORD 2
62 02026 000003 N3:       .WORD 3
63 02030 000012 N10:      .WORD 10.
64 02032 000014 N12:      .WORD 12.
65 02034 000015 N13:      .WORD 13.
66 02036 000017 N15:      .WORD 15.
67 02040 041710 N100:     .FLT2 100.
   02042 000000
68 02044              INDEX:    .BLKW 2.
69 02050              COUNT:    .BLKW 2.
70 02054              F:        .BLKW 6.
71 02070              FU:       .BLKW 24.
72 02150              FUU:      .BLKW 6.
73 02164              FY:       .BLKW 24.
74 02244              FYY:      .BLKW 6.
75 02260              EV:       .BLKW 20.
76 02330              EVV:      .BLKW 6.
77 02344              EV2:      .BLKW 20.
78 02414              EV22:     .BLKW 6.
79 02430 040200 F.ONE:     .FLT2 1.E0
   02432 000000

```

```

1      ; SBTL STOCHASTIC APPROXIMATION ALGORITHM
2      ;
3      ;
4      ; STAPR-- THIS SUBROUTINE UPDATES THE
5      ; PARAMETERS OF THE NOISE FILTER VECTOR
6      ; F(3) BY ITERATING TEN TIMES PER PASS
7      ; FU=VECTOR WITH 15 PREVIOUS VALUES OF U
8      ; FY=VECTOR WITH 15 PREVIOUS VALUES OF Y
9      ; FUU=VECTOR WITH 3 PREVIOUS VALUES OF U
10     ; FYY=VECTOR WITH 3 PREVIOUS VALUES OF Y
11     ;
12     ;
13     STAPR:  FADD      COUNT,F.ONE,COUNT;      K=K+1
14             FSHIFT   FU,N15;                  UPDATE FU & FY
15             FSHIFT   FY,N15;                  EACH STORING
16             MOVF      U,FU;                    15 PREVIOUS
17             MOVF      Y,FY;                    U & Y VALUES RESP.
18             CMP       #17,AN.K;                > 15 SAMPLES ?
19             BLT       IS;                      YES
20             JMP       STA.GO;                  NO,SKIP
21             ;
22             ; CAL. ERROR VECTOR [EV] & [EV]**2
23             ;
24     02532 016700 1$:  MOV      N13,R0;
25             177276
26     02536 016701      MOV      N12,R1
27             177270
28     02542 006301      ASL      R1
29     02544 006301      ASL      R1
30     02546      STA.E:  MMUL      FUU,PHI,DUM1,ONE,THREE,ONE
31     02566      MMUL      FYY+4,PHI+14,DUM2,ONE,TWO,ONE
32     02606      FSUB      FYY,DUM1,DUM1
33     02620      FADD      DUM1,DUM2,DUM1
34     02632 016761      MOV      DUM1,EV(R1)
35             177142
36             002260
37     02640 016761      MOV      DUM1+2,EV+2(R1)
38             177136
39             002262
40     02646      FMUL      DUM1,DUM1,DUM1
41     02660 016761      MOV      DUM1,EV2(R1)
42             177114
43             002344
44     02666 016761      MOV      DUM1+2,EV2+2(R1)
45             177110
46             002346
47     02674 024141      CMP      -(R1),-(R1)
48     02676      PUSHF     FUU+10

```

```

39 02706      PUSHF      FYY+10
40 02716      FSHIFT     FU,N15
41 02726      FSHIFT     FY,N15
42 02736      POPF        FU
43 02746      POPF        FY
44 02756 005300      DEC      R0
45 02760 020027      CMP      R0,#0
      000000
46 02764 001402      BEQ      115
47 02766 000167      JMP      STA.E
      177554

48
49      ;RESTORE FU,FY
50      ;
51 02772      115:      PUSHF      FUU+10
52 03002      PUSHF      FYY+10
53 03012      FSHIFT     FU,N15
54 03022      FSHIFT     FY,N15
55 03032      POPF        FU
56 03042      POPF        FU
57 03052      FSHIFT     FU,N15
58 03062      FSHIFT     FY,N15
59
60      ;KWANTY FORMULA FOR STOCHASTIC APPROX.
61      ;
62 03072 016700      MOV      N10,R0
      176732
63 03076      STA.F:    FADD      INDEX,F.ONE,INDEX
64 03110      FADD      INDEX,COUNT,DUM1
65 03122      FDIV      GAIN,DUM1,DUM1
66 03134      MMUL      EVV,F,DUM2,ONE,THREE,ONE
67 03154      FADD      EVV-4,DUM2,DUM2
68 03166      FADD      EV22,EV22+4,DUM3
69 03200      FADD      EV22+10,DUM3,DUM3
70 03212      FDIV      DUM2,DUM3,DUM2
71 03224      FMUL      DUM1,DUM2,DUM1
72 03236      MOVF      DUM1,DUM2
73 03252 000167      TST      DUM2;          IF DUM1>100.
      176520
74 03256 100000      BPL      13;          SKIP
75 03260 002167      RLC      #100000,DUM2
      100000
      176516
76 03266      13:      FSUB      N100,DUM2,DUM2
77 03300 005767      1ST      DUM2
      176500
78 03304 100436      BMI      STA.A

```

```

79      ;
80      ; NEW (F) CALCULATED HERE
81      ;
82 03306      FMUL      DUM1,EVV+10,DUM3
83 03320      FMUL      DUM1,EVV+4,DUM2
84 03332      FMUL      DUM1,EVV,DUM1
85 03344      FSUB      F,DUM1,F
86 03356      FSUB      F+4,DUM2,F+4
87 03370      FSUB      F+10,DUM3,F+10
88 03402      STA.A:    FSHIFT  EV,N13
89 03412      FSHIFT    EV2,N13
90 03422 005300      DEC      R0
91 03424 020027      CNP      R0,#0
          000000
92 03430 001402      BEQ      11$
93 03432 000167      JMP      STA.F
          177440
94 03436 005067 11$:  CLR      _INDEX
          176402
95 03442 005067      CLR      INDEX+2
          176400
96 03446 000205 STA.GO: RTS      R5
97      ;
98      ; SUBROUTINE FILTER FILTERS INCOMING DATA
99      ;
100 3450      FILTER:  SAVE      R5
101 3452      MMUL      F,FU,DUM1,ONE,THREE,ONE
102 3472      MMUL      F,FY,DUM2,ONE,THREE,ONE
103 3512      FADD      U,DUM1,U
104 3524      FADD      Y,DUM2,Y
105 3536      UNSAVE R5
106 3540 000205      RTS      R5

```

```

1          .SBTTL SAMPLING
2          ;
3          ;
4          177564 TPS=177564; PUNCH STATUS REG
5          177566 TPB=177566; PUNCH BUFFER REG
6          177520 ADCSR=177520; ADDR.OF CONTROL & STATUS REG
7          177522 CHNSLR=177522; ADDR. OF CHANNEL SELECTOR REG
8          177524 DATREG=177524; ADDR. OF DATA BUFFER REG.
9          ;
10         ;CLOCK---SUBROUTINE TO START SAMPLING PROCESS
11         ;
12 03542 005037 CLOCK: CLR      @#TPS; INITIALIZE SAMPLING & PUNCH
13         177564
14 03546 005037          CLR      @#ADCSR; CONTROL & STATUS REGS.
15         177520
16 03552 012737          MOV      #AD,@#110; SET UP INT. VECTOR
17         003632
18         000110
19 03560 012737          MOV      #240,@#112; FOR SAMPLING; PRY.=>
20         000240
21         000112
22 03566 013767          MOV      @#64,C.V1; SAVE OLD CONTENTS OF
23         000064
24         000032
25 03574 013767          MOV      @#66,C.V1+2; PRINTER INTR. VECTOR
26         000066
27         000026
28 03602 012737          MOV      #10,@#64; INTR. VECTOR FOR
29         004772
30         000064
31 03610 012737          MOV      #4200,@#66; PRINTING;PRY. <
32         004200
33         000066
34         ; REG. SET NO. 2
35 03616 012737          MOV      #100,@#ADCSR; HERE WE GO BABY
36         000100
37         177520
38 03624 000205          RTS      R5
39 03626          C.V1: .BLKW 2.
40         ;
41         ;
42         ;AD-----SERVICE ROUTINE FOR THE FIRST A/D
43         ; INTERRUPT TO CONVERT CH. 1
44         ;
45         ;
46 03632 000235 AD:      SPL      5; MASK OUT PRINTING INTERRUPT
47 03634 032737          BIT      #1,@#177570; ABORT PROGRAM ?
48         000001
49         177570
50 03642 001405          BEQ      25; NO,SKIP
51 03644 005037          CLR      @#ADCSR; YES,QUIT
52         177520
53 03650 005267          INC      QUIT; SET FLAG
54         000050

```

```

34 03654 000002      RTI
35 03656 012737 25:  MOV      #10000,@#CHNSLR;  CONVERT CH.1
      010000
      177522
36 03664 012737      MOV      #STRI,@#110;  SERVE STR1 NEXT
      003732
      000110
37 03672 022767      CMP      #0,AD.CHK;  CHECK PT. CLR ?
      000000
      000026
38 03700 001010      BNE      1$;  NO,SKIP
39 03702 005267      INC      AD.A;  INC COUNT OF SAMPLES
      000022
40 03706 026767      CMP      AD.A,MAXAD;  =MAX. COUNT?
      000016
      000006
41 03714 003402      BLE      1$;  NO,SKIP
42 03716 005037      CLR      @#ADCSR;  YES,FINISH
      177520
43 03722 000002 1$:  RTI
44 03724 000000 QUIT:  .WORD 0
45 03726 000001 AD.CHK: .WORD 1
46 03730 000000 AD.A:  .WORD 0
47
48      ;STRI--- SERVICE ROUTINE FOR THE 2ND A/D INTERRUPT
49      ; TO STORE CH.1 DATA & CONVERT CH.2
50      ;
51 03732 000235 STR1:  SPL      5;  MASK OUT PRINTING INTERRUPT
52 03734 013767      MOV      @DATREG,BUF1;  STORE CH.1 DATA
      177524
      000262
53 03742 032737      BIT      #1,@#177570;  ABORT PROGRAM
      000001
      177570
54 03750 001405      BEQ      4$;  NO,SKIP
55 03752 005037      CLR      @#ADCSR;  YES,QUIT
      177520
56 03756 005267      INC      QUIT;  SET FLAG
      177742
57 03762 000002      RTI
58 03764 022767 4$:  CMP      #0,AD.CHK;  CHECK PT. CLR?
      000000
      177734
59 03772 001416      BEQ      1$;  YES,SKIP
60 03774 012737      MOV      #AD,@#110;  NO
      003632
      000110
61 04002 016700      MOV      BUF1,R0;  TEST ABS(BUF1)
      000216
62 04006 005700      TST      R0
63 04010 100001      BPL      2$
64 04012 005400      NEG      R0

```

```

65 04014 020027 2$:      CMP      R0,#7;          >0.078 VOLT ?
      000007
66 04020 101402      BLOS      3$;          NO, SKIP
67 04022 005067      CLR      AD.CHK;          CLR CHECK PT.
      177700
68 04026 000002 3$:      RTI
69 04030 012737 1$:      MOV      #20000,@#CHNSLR; CONVERT CH.2
      020000
      177522
70 04036 012737      MOV      #STR2,@#110;      SERVE STR2 NEXT
      024046
      000110
71 04044 000002      RTI
72
73      ;STR2-----SERVICE ROUTINE TO STORE CH.2 DATA
74      ;      AND PROCESS THEM IN THE SELECTED METHOD
75      ;
76 04046 000235 STR2:    SPL      5; MASK OUT PRINTING INTERRUPT
77 04050 013707      MOV      @#DATHEG,BUF2; STORE CH.2 DATA
      177524
      000150
78 04056 032737      BIT      #1,@#177570;      ABORT PROGRAM ?
      000001
      177570
79 04064 001405      BEQ      1$;          NO,SKIP
80 04066 005037      CLR      @#ADCSR;          YES,QUIT
      177520
81 04072 005267      INC      QUIT;          SET FLAG
      177626
82 04076 000002      RTI
83 04100 022767 1$:      CMP      #0,1DELAY;      NEED DELAY ?
      000000
      000004
84 04106 001004      BNE      2$
85 04110 012737      MOV      #AD,@#110;          NO,SERVE AD NEXT
      003632
      000110
86 04116 000403      BR      3$
87 04120 012737 2$:      MOV      #DELAY,@#110;      YES,SERVE DELAY NEXT
      004232
      000110
88 04126 022767 3$:      CMP      #0,ANFLG; FINISH LAST ANALYSIS ?
      000000
      000074
89 04134 001402      BEQ      4$
90 04136 000236      SPL      6
91 04140 104060      EMT      60;          IF NO,ERROR EXIT
92 04142 005267 4$:      INC      ANFLG;          SET FLAG
      000062
93 04146 016703      MOV      BUF1,R3;          CONVERT CH.1
      000052
94 04152 004567      JSR      R5,CONVSN;          TO FLT.PT.
      000132

```



```

95 04156 012667      MOV      (SF)+,U;          STORE IN U
      000032
96 04162 012667      MOV      (SP)+,U+2
      000030
97 04166 016703      MOV      BUF2,R3;          CONVERT CH.2
      000034
98 04172 004567      JSR      R5,CONVSN;        TO FLT. PT.
      000112
99 04176 012667      MOV      (SP)+,Y;          STORE IN Y
      000016
100 4202 012667      MOV      (SP)+,Y+2
      000014
101 4206 004567      JSR      R5,ANAL;          GO TO PROCESS DATA

```

IDENT MACRO VR05A 01-JAN-72 05:23 PAGE 6+
SAMPLING

```

      000270
102 4212 000002      RTI
103
104 4214      U:      .BLKW 2.
105 4220      Y:      .BLKW 2.
106 4224 000000 BUF1:  .WORD 0
107 4226 000000 BUF2:  .WORD 0
108 4230 000000 ANFLG: .WORD 0
109
110      ; DELAY----SERVICE ROUTINE AS A DELAY LOOP
111      ; TO MODIFY THE SAMPLING FREQ.
112
113 4232 000235 DELAY: SPL      5; MASK OUT PRINTING INTERRUPT
114 4234 032737      BIT      #1,0#177570; ABORT PROGRAM ?
      000001
      177570
115 4242 001405      BEQ      25; NO,SKIP
116 4244 005037      CLR      @#ADCSR; YES,QUIT
      177520
117 4250 005267      INC      QUIT; SET FLAG
      177450
118 4254 000002      RTI
119 4256 005267 2$:   INC      DE.D; INC DELAY COUNT
      000024
120 4262 026767      CMP      DE.D,IDELAY; ENOUGH ?
      000020
      000004
121 4270 002405      BLT      15; NO HURRY, SON
122 4272 012737      MOV      #AD,0#110; BACK TO CH.1
      003632
      000110
123 4300 005067      CLR      DE.D; RESET DE.D
      000002
124 4304 000002 1$:   RTI
125 4306 000000 DE.D: .WORD 0

```

```

126      .SBTTL  DECODING OF OUTPUT OF A/D CONVERTER
127      ;
128      ;      SUBROUTINE TO CONVERT A BINARY CODED OUTPUT
129      ;      FROM A/D CONVERTER INTO NORMALISED
130      ;      FLOATING POINT FORMAT
131      ;
132 4310 012667 CONVSN: MOV      (SP)+,IMPSP;      SAVE STACK
                        000154
133 4314 005703      IST      R3;                  POSITIVE?
134 4316 100004      BPL      1$;                  YES,SKIP
135 4320 012767      MOV      #100000,SIGN;      SET SIGN BIT
                        100000
                        000152
136 4326 005403      NEG      R3;                  GET 2'S COMPLIMENT
137 4330 010301 1$:  MOV      R3,R1;                  SAVE ON R1
138 4332 022703      CMP      #0,R3;                  ZERO ?
                        000000
139 4336 001436      BEQ      3$
140 4340 006303      ASL      R3;
141 4342 006303      ASL      R3;
142 4344 060103      ADD      R1,R3;                  *5
143 4346 006303 2$:  ASL      R3;                  SHIFT LEFT
144 4350 105367      DECB     EXP;                  DECREMENT EXPONENT
                        000120
145 4354 005703      IST      R3;                  MSB SET ?
146 4356 100373      BPL      2$;                  NO,GO BACK
147 4360 042703      BIC      #100100,R3;          CLEAR MSB
                        100000
148 4364 110067      MOVB     R3,TEMBUF+3;          STORE 2ND WORD
                        000105
149 4370 016746      MOV      TEMBUF+2,-(SP);      PUSH ON STACK
                        000100
150 4374 000303      SWAB     R3;                  GET 2ND BYTE
151 4376 110367      MOVB     R3,TEMBUF;          STORE MANTISA
                        000070
152 4402 000367      SWAB     EXP;
                        000070
153 4406 006067      ROR      EXP;                  ALIGN EXPONENT BYTE
                        000064
154 4412 056767      BIS      EXP,TEMBUF;          COPY EXPONENT
                        000060
                        000052
155 4420 056767      BIS      SIGN,TEMBUF;          COPY SIGN
                        000054
                        000044
156 4426 016746      MOV      TEMBUF,-(SP);      PUSH 1ST WORD
                        000040
157 4432 000402      BR       4$;                  ON STACK

```

```

158 4434 005046 3$: CLR -(SP); FLOATING ZERO
159 4436 005046 CLR -(SP)
160 4440 005067 4$: CLR TEMBUF; READY TO QUIT
    000026
161 4444 005067 CLR TEMBUF+2
    000024
162 4450 005067 CLR SIGN
    000024
163 4454 012767 MOV #000210,EXP; RESTORE INITIAL EXP
    000210
    000014
164 4462 016746 MOV IMPSP,-(SP); RESTORE STACK
    000002
165 4466 000205 RTS R5
166 4470 000000 IMPSP: .WORD 0
167 4472 TEMBUF: .DLKW 2.
168 4476 210 EXP: .BYTE 210,0
    4477 000
169 4500 000000 SIGN: .WORD 0

```

```

1      .SBTTL DATA ANALYSIS
2      ;
3      ;ANAL---SUBROUTINE TO JUMP TO IDENT. ALGORITHM
4      ;      AND STORE RESULTS IN DATA BUFFER
5      ;
6      ;
7 004502 004567 ANAL:   JSR      R5,UPDATE;   UPDATE INFORM. MATRIX
           174132
8 004506 004567       JSR      R5,METHOD;   GO TO ALGORITHM
           173312
9 004512 016700       MOV      AN.K1,R0;      LOAD POINTER
           000114
10 04516 016760       MOV      PHI,FI1(R0);   STORE RESULTS
           174672
           000000
11 04524 016760       MOV      PHI+2,FI1+2(R0)
           174666
           000002
12 04532 016760       MOV      PHI+4,FI2(R0)
           174662
           003720
13 04540 016760       MOV      PHI+6,FI2+2(R0)
           174656
           003720
14 04546 016760       MOV      PHI+10,FI3(R0)
           174652
           007640
15 04554 016760       MOV      PHI+12,FI3+2(R0)
           174646
           007642
16 04562 016760       MOV      PHI+14,FI4(R0)
           174642
           013560
17 04570 016760       MOV      PHI+16,FI4+2(R0)
           174636
           013562
18 04576 016760       MOV      PHI+20,FI5(R0)
           174632
           017500
19 04604 016760       MOV      PHI+22,FI5+2(R0)
           174626
           017502
20 04612 022020       CMP      (R0)+,(R0)+;   INC R0 BY 4
21 04614 010067       MOV      R0,AN.K1;      STORE POINTER
           000012
22 04620 005267       INC      AN.K;   INC. ITERATION COUNT
           000010
23 04624 005067       CLR      ANFLG;         CLR FLAG
           177400
24 04630 000205       RTS      R5
25 04632 000000 AN.K1:  .WORD 0
26 04634 000000 AN.K:  .WORD 0

```

```

1      .SBTTL SUBROUTINES FOR PRINTING
2      ;
3      ;PRINT-----SUBROUTINE TO MONITOR PROGRESS OF
4      ;      ANALYSIS AND PRINTING
5      ;
6 004636 032737 PRINT: BIT      #1,@#177570;      ABORT PROGRAM ?
          000001
          177570
7 004644 001401      BEQ      2$;      NO,GO AHEAD
8 004646 000426      BR      5$;      YES,QUIT
9 004650 026767 2$:   CMP      10.K,AN.K;      IS PRINTING LAGGING ?
          000114
          177756
10 04656 002402      BLT      3$;      IF YES, GO AHEAD
11 04660 000501      WAIT;      IF NO,WAIT AND
12 04662 000765      BR      PRINT;      CHECK AGAIN, SON
13 04664 012737 3$:   MOV      #100,@#TPS; PRINT NEXT CHARACTER
          000100
          177564
14 04672 000001 4$:   WAIT
15 04674 022767      CMP      #0,PR.FI;      FINISH 1 LINE ?
          000000
          000064
16 04702 031373      BNE      4$;      NO,CHECK AGAIN
17 04704 005267      INC      10.K;      INC I/O COUNT
          000060
18 04710 005267      INC      PR.FI;      RESET FLAG
          033052
19 04714 026767      CMP      10.K,MAXAD;      ALL DONE ?
          000050
          000006
20 04722 002745      BLT      PRINT;      NO,CARRY ON SON
21 04724 005037 5$:   CLR      @#ADCSR
          177520
22 04730 012700      MOV      #77777,R0
          077777
23 04734 016737      MOV      C.V1,@#64; RESTORE INT. VECTOR
          176666
          000064
24 04742 016737      MOV      C.V1+2,@#66
          176662
          000066
25 04750 012737      MOV      #100,@#TPS
          000100
          177564
26 04756 000005      RESET
27 04760 000000      HALT
28 04762 104060      EMT      60
29 04764 000205      RTS      R5
30 04766 000001 PR.FI: .WORD 1
31 04770 000000 10.K: .WORD 0

```

```

32      ;
33      ;ECO----MACRO DEFINITION TO FETCH A FLOATING
34      ;      WORD AT NUM TO BE CONVERTED TO 9-BYTES
35      ;      ASCII CODES STARTING FROM ASCII
36      ;
37      .MACRO ECO ASCII,NUM
38          MOV      #ASCII,-(SP)
39          MOV      #11,-(SP)
40          MOV      #2,-(SP)
41          MOV      #1,-(SP)
42          MOV      NUM+2(R2),-(SP)
43          MOV      NUM(R2),-(SP)
44          JSR      PC,$ECO
45          .ENDM
46      .GLOBL $ECO
47      ;
48      ;
49      ;IO----SERVICE SUBROUTINE TO DO ASCII CONVERSION
50      ;      AND FILL OUTPUT PRINTER BUFFER
51      ;
52 04772 016702 IO:      MOV      IO.K,R2;      R2=POINTER
53      177772
54 04776 006302      ASL      R2
55 05000 006302      ASL      R2
56 05002 010267      MOV      R2,IO.R2
57      000346
58 05006      ECO $F11,F11;      CONVERT 1ST WORD
59 05042 016702      MOV      IO.R2,R2
60      000306
61 05046      ECO $F12,F12;      CONVERT 2ND WORD
62 05102 016702      MOV      IO.R2,R2
63      000246
64 05106      ECO $F13,F13;      CONVERT 3RD WORD
65 05142 016702      MOV      IO.R2,R2
66      000206
67 05146      ECO $F14,F14;      CONVERT 4TH WORD
68 05202 016702      MOV      IO.R2,R2
69      000146
70 05206      ECO $F15,F15;      CONVERT 5TH WORD
71 05242 012703      MOV      #BUFST,R3
72      005470
73 05246      JSR      R5,IOF;      3-DIGIT LINE NO.
74      000106
75 05252 012737      MOV      #PRN,@#64
76      005266
77      000064
78 05260 012767      MOV      #BUFST,PONTR
79      005470
80      000070
81 05266 032737 PRN:   BIT      #1,@#177570;      ABORT PROGRAM ?
82      000001
83      177570
84 05274 001403      BEQ      1$;      NO,GO AHEAD

```

```

71 05276 005037      CLR      @#ADCSR;      YES,QUIT
      177520
72 05302 000423      BR        2$
73 05304 105767 1$:  TSTB      TPS;          PRINTER READY ?
      177564
74 05310 100366      BPL       PRN
75 05312 016700      MOV       PONTR,R0;     YES,LOAD POINTER
      000040
76 05316 112037      MOVB      (R0)+,@#TPB;  LOAD PRINTER BUFFER
      177566
77 05322 010067      MOV       R0,PONTR
      000030
78 05326 020027      CMP       R0,#BUFEND;   FINISH 1 LINE ?
      005571
79 05332 001007      BNE       2$
80 05334 005037      CLR       @#TPS;       YES
      177564
81 05340 005067      CLR       PR.F1;       CLR FLAG
      177422
82 05344 012737      MOV       #IO,@#64;     RESTORE VECTOR
      004772
      000064
83 05352 020002 2$:  RTI
84 05354 000000 IO.R2:  .WORD 0
85 05356 000000 PONTR: .WORD 0
86
87      ; ICF--- SUBROUTINE TO FORM 3-DIGIT ASCII CODE
88      ; IN ASCENDING ORDER EACH TIME IT IS CALLED
89      ; ASCII BUFFER MUST FIRST BE INITIALISED
90      ; BY A NO.
91      ;
92 05360 126327 IOF:  CMPB      2(R3),#71;    DIGIT 3=0 ?
      000002
      000071
93 05366 001403      BEQ        1$
94 05370 105263      INCB       2(R3);        NO,INC BY 1
      000002
95 05374 000434      BR         5$
96 05376 112763 1$:  MOVB      #60,2(R3);    DIGIT 3=0
      000060
      000002
97 05404 126327      CMPB      1(R3),#60;    DIGIT 2=0 OR SPACE ?
      000001
      000060
98 05412 003004      BGT        2$
99 05414 112763      MOVB      #61,1(R3);    DIGIT 2=1
      000061
      000001
100 5422 000421      BR         5$
101 5424 126327 2$:  CMPB      1(R3),#71;    DIGIT 2=9 ?
      000001
      000071
102 5432 001403      BEQ        3$

```

```

103 5434 105263      INCB      1(R3);          NO, INC BY 1
      000001
104 5440 000412      BR        5$
105 5442 112763 3$:  MOV      #60,1(R3);      DIGIT 2=0
      000060
      000001
106 5450 121327      CMPB      (R3),#40;      DIGIT 1=SPACE ?
      000040
107 5454 001003      BNE       4$
108 5456 112713      MOV      #61,(R3);      DIGIT 1=1
      000061
109 5462 000401      BR        5$
110 5464 105213 4$:  INCB      (R3);          DIGIT 1=1
111 5466 000205 5$:  RTS       R5
112      ;
113      ;OUTPUT ASCII BUFFER FOR 5 PARAMETERS
114      ;
115 5470      BUFST:
116 5470      040      .BYTE 40,40,60,40,40,40
      5471      040
      5472      060
      5473      040
      5474      040
      5475      040
      5476      $F11: .BLKB 11
118 5507      040      .BYTE 40,40,40
      5510      040
      5511      040
119 5512      $F12: .BLKB 11
120 5523      040      .BYTE 40,40,40
      5524      040
      5525      040
121 5526      $F13: .BLKB 11
122 5537      040      .BYTE 40,40,40
      5540      040
      5541      040
123 5542      $F14: .BLKB 11
124 5553      040      .BYTE 40,40,40
      5554      040
      5555      040
125 5556      $F15: .BLKB 11
126 5567      015      .BYTE 15,12
      5570      012
127      005571 'BUFEND=.
128      .EVEN
129      ;
130      ;DATA BUFFER FOR 5 PARAMETERS IN 2-WORD F.P.
131      ;500 EACH
132      000000 .CSECT DATA
133 0000      F11:      .BLKW 1000.
134 3720      F12:      .BLKW 1000.
135 7640      F13:      .BLKW 1000.
136 3560      F14:      .BLKW 1000.
137 7500      F15:      .BLKW 1000.

```



```

138      005572'      .CSECT
139      .SBTTL ARITHMATICAL SUBROUTINES
140      ;
141      ;
142      ; SUBROUTINE MULFP => A*B=C
143      ;
144 5572      MULFP:  SAVE R0,R1,R2
145 5600 012500      MOV      (R5)+,R0;      ADDR.OF
146 5602 012501      MOV      (R5)+,R1;      A,B&C
147 5604 012502      MOV      (R5)+,R2
148 5606      SAVE      R5
149 5610 016046      MOV      2(R0),-(SP)
      000002
150 5614 011046      MOV      (R0),-(SP)
151 5616 016146      MOV      2(R1),-(SP)
      000002
152 5622 011146      MOV      (R1),-(SP)
153 5624 004467      JSR      R4,$POLSI.
      000000G
154 5630 000000G      .WORD      $MLR
155 5632 005634'      .WORD      .+2
156 5634 012612      MOV      (SP)+,(R2);      RESULT
157 5636 012662      MOV      (SP)+,2(R2)
      000002
158 5642      UNSAVE R0,R1,R2,R5
159 5652 000205      RTS      R5
160      ;
161      ; SUBROUTINE ADDFP => A+B=C
162      ;
163 5654      ADDFP:  SAVE R0,R1,R2
164 5662 012500      MOV      (R5)+,R0;      ADDR.OF
165 5664 012501      MOV      (R5)+,R1;      A,B&C
166 5666 012502      MOV      (R5)+,R2
167 5670      SAVE      R5
168 5672 016046      MOV      2(R0),-(SP)
      000002
169 5676 011046      MOV      (R0),-(SP)
170 5700 016146      MOV      2(R1),-(SP)
      000002
171 5704 011146      MOV      (R1),-(SP)
172 5706 004467      JSR      R4,$POLSH
      000000G
173 5712 000000G      .WORD      $ADR
174 5714 005716'      .WORD      .+2
175 5716 012612      MOV      (SP)+,(R2);      RESULT
176 5720 012662      MOV      (SP)+,2(R2)
      000002
177 5724      UNSAVE R0,R1,R2,R5
178 5734 000205      RTS      R5

```

```

179      ;
180      ; SUBROUTINE DIVFP => A/B=C
181      ;
182 5736  DIVFP:  SAVE R0,R1,R2
183 5744 012500      MOV      (R5)+,R0;      ADDR.OF
184 5746 012501      MOV      (R5)+,R1;      A,B &C
185 5750 012502      MOV      (R5)+,R2
186 5752      SAVE      R5
187 5754 016046      MOV      2(R0),-(SP)
           000002
188 5760 011046      MOV      (R0),-(SP)
189 5762 016146      MOV      2(R1),-(SP)
           000002
190 5766 011146      MOV      (R1),-(SP)
191 5770 004467      JSR      R4,$POLSH
           000000G
192 5774 000000G      .WORD      $DVR
193 5776 006000      .WORD      .+2
194 6000 012612      MOV      (SP)+,(R2);      RESULT
195 6002 012662      MOV      (SP)+,2(R2)
           000002
196 6006      UNSAVE R0,R1,R2,R5
197 6016 000205      RTS      R5
198      ;
199      ; SUBROUTINE SUBFP => A-B=C
200      ;
201 6020  SUBFP:  SAVE R0,R1,R2
202 6026 012500      MOV      (R5)+,R0;      ADDR.OF
203 6030 012501      MOV      (R5)+,R1;      A,B &C
204 6032 012502      MOV      (R5)+,R2
205 6034      SAVE      R5
206 6036 016046      MOV      2(R0),-(SP)
           000002
207 6042 011046      MOV      (R0),-(SP)
208 6044 016146      MOV      2(R1),-(SP)
           000002
209 6050 011146      MOV      (R1),-(SP)
210 6052 004467      JSR      R4,$POLSH
           000000G
211 6056 000000G      .WORD      $SBR
212 6060 006062      .WORD      .+2
213 6062 012612      MOV      (SP)+,(R2);      RESULT
214 6064 012662      MOV      (SP)+,2(R2)
           000002
215 6070      UNSAVE R0,R1,R2,R5
216 6100 000205      RTS      R5

```

```

217      ;
218      ;MATRIX MULTIPLY (A)*(B)=(C)
219      ;
220 6102  MMUL:  SAVE R0,R1,R2,R3,R4
221 6114      MOV      (R5)+,A.MUL;  ADDR OF A
      012567
      000376
222 6120      MOV      (R5)+,B.MUL;  B AND C
      012567
      000374
223 6124      MOV      (R5)+,C.MUL
      012567
      000372
224 6130      MOV      @(R5)+,R0;  VALUES OF
225 6132      MOV      @(R5)+,R1;  ROWA, COLA
226 6134      MOV      @(R5)+,R3;  AND COLB
227 6136      SAVE R5
228 6140      ASL      R0
229 6142      ASL      R0;  4*ROWA
230 6144      MOV      R0,NXTA;  NEXT ELEM. OF A
      000354
231 6150      MUL      R0,R3;  4*ROWA*COLB =
232 6152      ADD      C.MUL,R3;  BASE +
      000344
233 6156      MOV      R3,NO.ELM;  NO. OF ELEM. IN ROWA
      000352
234 6162      MOV      R1,R2
235 6164      MUL      R0,R1;  4*ROWA*COLA
236 6166      CMP      -(R1),-(R1);  -4 =>
237 6170      MOV      R1,NXTCOL;  SWITCH COL IN B
      000332
238 6174      ADD      A.MUL,NXTCOL;  +BASE
      000316
      000324
239 6202      SUB      R0,R1;  -4*ROWA =>
240 6204      MOV      R1,NXTROW;  SWITCH ROW IN A
      000320
241 6210      DEC      R2;  COLA-1
242 6212      MOV      R2,COLA.1
      000320
243 6216      ASL      R2
244 6220      ASL      R2;  4*(COLA-1) =>
245 6222      MOV      R2,SAMCOL;  SAME COL IN B
      000304
246 6226      MOV      A.MUL,R0;  BASE ADDR OF A
      000264
247 6232      MOV      B.MUL,R1;  BASE ADDR OF B
      000262
248 6236      MOV      C.MUL,R3;  BASE ADDR OF C
      000260
249 6242      MOV      COLA.1,R2;  NO. OF ADDITIONS
      000270
250 6246      15:  SAVE R0,R1,R2,R3
251 6256      MOV      2(R0),-(SP)
      000002

```

252	6262	011046	MOV	(R0),-(SP)	
253	6264	016146	MOV	2(R1),-(SP)	
		000002			
254	6270	011146	MOV	(R1),-(SP)	
255	6272	004467	JSR	R4,\$POLSH	
		000000G			
256	6276	000000G	.WORD	\$MLR	
257	6300	006302	.WORD	.*2	
258	6302		POPF	TMPMUL	
259	6312		UNSAVE	R0,R1,R2,R3	
260	6322	022702	CMP	#0,R2;	A = COL VECTOR ?
		000000			
261	6326	001437	BEQ	3\$;	YES,SKIP
262	6330	066700 2\$:	ADD	NXTA,R0; NXT ELM. OF A,SAME ROW	
		000170			
263	6334	022121	CMP	(R1)+,(R1)+;	NEXT ELM. OF B
264	6336		SAVE	R0,R1,R2,R3	
265	6346		PUSHF	TMPMUL;	PREVIOUS PARTIAL SUM
266	6356	016046	MOV	2(R0),-(SP)	
		000002			
267	6362	011046	MOV	(R0),-(SP)	
268	6364	016146	MOV	2(R1),-(SP)	
		000002			
269	6370	011146	MOV	(R1),-(SP)	
270	6372	004467	JSR	R4,\$POLSH	
		000000G			
271	6376	020000G	.WORD	\$MLR	
272	6400	000000G	.WORD	\$ADR	
273	6402	006404	.WORD	.*2	
274	6404		POPF	TMPMUL	
275	6414		UNSAVE	R0,R1,R2,R3	
276	6424	077237	SOB	R2,2\$	
277	6426	016702 3\$:	MOV	COLA.1,R2;	RESTORE R2
		000104			
278	6432	016723	MOV	TMPMUL,(R3)+;	STORE RESULT
		000102			
279	6436	016723	MOV	TMPMUL+2,(R3)+;	IN C
		000100			
280	6442	026703	CMP	NO.ELM,R3;	DONE ?
		000066			
281	6446	001414	BEQ	10\$;	YES,QUIT
282	6450	020067	CMP	R0,NXTCOL;	NEXT COL OF B?
		000052			
283	6454	001405	BEQ	4\$;	YES,SKIP
284	6456	166701	SUB	SAMCOL,R1;	NO,SAME COL
		000050			
285	6462	166700	SUB	NXTROW,R0;	NEXT ROW OF A
		000042			
286	6466	000667	BR	1\$;	CARRY ON, SON
287	6470	016700 4\$:	MOV	A.MUL,R0;	BACK TO ROW 1 OF A
		000022			
288	6474	022121	CMP	(R1)+,(R1)+;	NEXT COL OF B
289	6476	000663	BR	1\$;	CARRY ON, SON

```

290 6500      105:
291 6500      UNSAVE R0,R1,R2,R3,R4,R5
292 6514 000205      RTS      R5
293 6516 000000 A.MUL: 0
294 6520 000000 B.MUL: 0
295 6522 000000 C.MUL: 0
296 6524 000000 NXTA: 0
297 6526 000000 NXTCOL: 0
298 6530 000000 NXTROW: 0
299 6532 000000 SAMCOL: 0
300 6534 000000 NO.ELM: 0
301 6536 000000 COLA.1: 0
302 6540      TMPMUL: .BLKW 2.
303
304      ;MATRIX ADDITION [A]+[B]=[C]
305      ;
306 6544      MADD:      SAVE R0,R1,R2,R3,R4
307 6556 012500      MOV      (R5)+,R0;      ADDR OF A
308 6560 012502      MOV      (R5)+,R2;      B AND C
309 6562 012503      MOV      (R5)+,R3
310 6564 013501      MOV      @ (R5)+,R1;      NO. OF ROW
311 6566 070135      MUL      @ (R5)+,R1;      ROW*COL
312 6570      SAVE R5
313 6572 005301      ASL      R1
314 6574 006301      ASL      R1;      4*RC/*COL
315 6576 060001      ADD      R0,R1;      +BASE
316 6600      15:      SAVE R0,R1,R2,R3
317 6610 016046      MOV      2(R0),-(SP)
318 6614 011046      MOV      (R0),-(SP)
319 6616 216246      MOV      2(R2),-(SP)
320 6622 011246      MOV      (R2),-(SP)
321 6624 004467      JSR      R4,$POLSH
322 6630 000000G      .WORD $ADR
323 6632 006634      .WORD .+2
324 6634      POPF TMPADD
325 6644      UNSAVE R0,R1,R2,R3
326 6654 016723      MOV      TMPADD,(R3)+; STORE RESULT
327 6660 016723      MOV      TMPADD+2,(R3)+; IN C
328 6664 022020      CMP      (R0)+,(R0)+;      INC R0,R2
329 6666 022222      CMP      (R2)+,(R2)+;      BY 4
330 6670 020001      CMP      R0,R1;      DONE ?
331 6672 001342      BNE      15;      NO,GO BACK
332 6674      UNSAVE R0,R1,R2,R3,R4,R5
333 6710 000205      RTS      R5
334 6712      TMPADD: .BLKW 2.

```

```

335      ;
336      ;MATRIX SUBTRACTION [A]-[B]=[C]
337      ;
338 6716  MSUB:  SAVE R0,R1,R2,R3,R4
339 6730 012500  MOV      (R5)+,R0;      ADDR OF A
340 6732 012502      MOV      (R5)+,R2;      B AND C
341 6734 012503      MOV      (R5)+,R3
342 6736 013501      MOV      @(R5)+,R1;      NO. OF ROW
343 6740 070135      MUL      @(R5)+,R1;      ROW*COL
344 6742      SAVE R5
345 6744 006301      ASL      R1
346 6746 006301      ASL      R1;      4*ROW*COL
347 6750 060001      ADD      R0,R1;      +BASE
348 6752      IS:  SAVE R0,R1,R2,R3
349 6762 016046      MOV      2(R0),-(SP)
350      000002
350 6766 011046      MOV      (R0),-(SP)
351 6770 016246      MOV      2(R2),-(SP)
352      000002
352 6774 011246      MOV      (R2),-(SP)
353 6776 004467      JSR      R4,$POLSH
354      000000G
354 7002 000000G      .WORD $SBR
355 7004 007006      .WORD .+2
356 7006      POPF TMPSUB
357 7016      UNSAVE R0,R1,R2,R3
358 7026 016723      MOV      TMPSUB,(R3)+; STORE RESULT
359      000032
359 7032 015723      MOV      TMPSUB+2,(R3)+; IN C
360      000030
360 7036 022020      CMP      (R0)+,(R0)+; INC R0,R2
361 7040 022222      CMP      (R2)+,(R2)+; BY 4
362 7042 020001      CMP      R0,R1;      DONE ?
363 7044 001342      BNE      IS;      NO,GO BACK
364 7046      UNSAVE R0,R1,R2,R3,R4,R5
365 7062 000205      RTS      R5
366 7064      TMPSUB: .BLKW 2.
367      ;
368      ;MATRIX SCALAR MULTIPLICATION [A]*SC=[B]
369      ;
370 7070  MSC:  SAVE R0,R1,R2,R3,R4
371 7102 012500      MOV      (R5)+,R0;      ADDR OF A
372 7104 012501      MOV      (R5)+,R1;      ADDR OF B
373 7106 012502      MOV      (R5)+,R2;      ADDR OF SCALAR
374 7110 013503      MOV      @(R5)+,R3;      NO. OF ROW
375 7112 070335      MUL      @(R5)+,R3;      ROW*COL
376 7114      SAVE R5
377 7116 006303      ASL      R3
378 7120 006303      ASL      R3;      4*ROW*COL
379 7122 060003      ADD      R0,R3;      +BASE
380 7124      IS:  SAVE R0,R1,R2,R3
381 7134 016046      MOV      2(R0),-(SP)
382      000002

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382 7140 011046      MOV      (A0),-(SP)
383 7142 016246      MOV      2(R2),-(SP)
                        000002
384 7146 011246      MOV      (R2),-(SP)
385 7150 004467      JSR      R4,$POLSH
                        000000G
386 7154 000000G      .WORD $MLR
387 7156 007160      .WORD .+2
388 7160              POPF TMPMSC
389 7170              UNSAVE R0,R1,R2,R3
390 7200 016721      MOV      TMPMSC,(R1)+; STORE RESULT
                        000030
391 7204 016721      MOV      TMPMSC+2,(R1)+; IN B
                        000026
392 7210 022020      CMP      (R0)+,(R0)+; INC R0 BY 4
393 7212 020003      CMP      R0,R3; DONE ?
394 7214 001343      BNE      1$
395 7216              UNSAVE R0,R1,R2,R3,R4,R5
396 7232 000205      RTS      R5
397 7234              TMPMSC: .BLKW 2.
398
399
400
401
402 7252 012500      DSC:      SAVE R0,R1,R2,R3,R4
403 7254 012501      MOV      (R5)+,R0; ADDR OF A
404 7256 012502      MOV      (R5)+,R1; ADDR OF B
405 7260 013503      MOV      (R5)+,R2; ADDR OF SCALAR
406 7262 070335      MOV      @(R5)+,R3; NO. OF ROW
407 7264              MUL      @(R5)+,R3; ROW*COL
408 7266 006303      SAVE R5
409 7270 006303      ASL      R3
410 7272 060003      ASL      R3; 4*ROW*COL
411 7274              ADD      R0,R3; +BASE
412 7304 016046      1$:      SAVE R0,R1,R2,R3
                        MOV      2(R0),-(SP)
                        000002
413 7310 011046      MOV      (R0),-(SP)
414 7312 016246      MOV      2(R2),-(SP)
                        000002
415 7316 011246      MOV      (R2),-(SP)
416 7320 004467      JSR      R4,$POLSH
                        000000G
417 7324 000000G      .WORD $DVR
418 7326 007330      .WORD .+2
419 7330              POPF TMPDSC
420 7340              UNSAVE R0,R1,R2,R3
421 7350 016721      MOV      TMPDSC,(R1)+; STORE RESULT
                        000030
422 7354 016721      MOV      TMPDSC+2,(R1)+; IN B
                        000026
423 7360 022020      CMP      (R0)+,(R0)+; INC R0 BY 4
424 7362 020003      CMP      R0,R3; DONE ?
425 7364 001343      BNE      1$
426 7366              UNSAVE R0,R1,R2,R3,R4,R5
427 7402 000205      RTS      R5
428 7404              TMPDSC: .BLKW 2.

```

```

429      ;
430      ; MATRIX TRANSPOSITION
431      ;
432 7410   MTRN:  SAVE R0,R1,R2,R3,R4
433 7422 012567   MOV      (R5)+,A.TRN;      ADDR OF A
          000120
434 7426 012567   MOV      (R5)+,AT.TRN;     ADDR OF AT
          000116
435 7432 013500   MOV      @(R5)+,R0;        NO. OF ROW
436 7434 013502   MOV      @(R5)+,R2;        NO. OF COL
437 7436          SAVE R5
438 7440 010267   MOV      R2,CO.TRN;        R2=C
          000106
439 7444 010005   MOV      R0,R5;            R0=R
440 7446 070502   MUL      R2,R5;            R5=RC
441 7450 010503   MOV      R5,R3
442 7452 026303   ASL      R3
443 7454 006323   ASL      R3
444 7456 024343   CMP      -(R3),-(R3);      R3=4RC-4
445 7460 010304   MOV      R3,R4;            R4=4RC-4
446 7462 010001   MOV      R0,R1
447 7464 026301   ASL      R1
448 7466 036301   ASL      R1;              R1=4R
449 7470 066703   ADD      A.TRN,R3;         R3=4RC-4+BASE
          000052
450 7474 016346 15:   MOV      2(R3),-(SP);  GET ELEMENTS
          000002
451 7500 011346   MOV      (R3),-(SP);  IN REVERSE ORDER
452 7502 160103   SUB      R1,R3;  NEXT HIGHER ELEM.
453 7504 077205   SOB      R2,15;         FINISH 1 COL ?
454 7506 016702   MOV      CO.TRN,R2;      YES,RESTORE R2
          000040
455 7512 060403   ADD      R4,R3;          NEXT COL
456 7514 077011   SOB      R0,15;         DONE ?
457 7516 016700   MOV      AT.TRN,R0
          000026
458 7522 012620 25:   MOV      (SP)+,(R0)+; STORE TRANSPOSED
459 7524 012620   MOV      (SP)+,(R0)+; RESULTS
460 7526 077503   SOB      R5,25;         DONE ?
461 7530          UNSAVE R0,R1,R2,R3,R4,R5
462 7544 000205   RTS R5
463 7546 000000 A.TRN: 0
464 7550 000000 AT.TRN: 0
465 7552 000000 CO.TRN: 0

```



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466      ;
467      ;SHIFT---SUBROUTINE THAT SHIFTS ALL NUM ELEMENTS
468      ;          OF VECTOR A 1 PLACE DOWN
469      ;
470 7554 SHIFT: SAVE R1
471 7556 012567 MOV      (R5)+,S.A;          ADDR OF A
      000046
472 7562 013501 MOV      @(R5)+,R1;          VALUE OF NUM
473 7564      SAVE      R5
474 7566 005301 DEC      R1;
475 7570 006301 ASL      R1;          4*(NUM-1)
      006301      +BASE ADDR
476 7572 006301 ASL      R1
477 7574 066701 ADD      S.A,R1
      000030
478 7600 016111 15: MOV      -4(R1),(R1)
      177774
479 7604 016161 MOV      -2(R1),2(R1)
      177776
      000002
480 7612 024141 CMP      -(R1),-(R1)
481 7614 020167 CMP      R1,S.A
      000010
482 7620 001367 BNE      15
483 7622      UNSAVE R1,R5
484 7626 000205 RTS      R5
485 7630 000000 S.A:      .WORD 0
486      000000      .END START
487

```

SYMBOL TABLE

A	001010R	AD	003632R	ADCSR =	177520		
ADDFP	005654R	AD.A	003730R	AD.CHK	003726R		
ANAL	004502R	ANFLG	004230R	AN.K	004634R		
AN.K1	004632R	AT	001010R	AT.TRN	007550R		
A.MUL	006516R	A.TRN	007546R	BUFLND=	005571R		
BUFST	005470R	BUF1	004224R	BUF2	004226R		
B.MUL	006520R	CHNSLR=	177522	CLOCK	003542R		
COLA.1	006536R	CONVSN	004310R	COUNT	002050R		
CO.TRN	007552R	C.MUL	006522R	C.VI	003626R		
DATREG=	177524	DELAY	004232R	DE.D	004306R		
DIALOG=	***** G	DIVFP	005736R	DSC	007240R		
DUM1	002000R	DUM2	002004R	DUM3	002010R		
EV	002260R	EVV	002330R	EV2	002344R		
EV22	002414R	EXP	004476R	F	002054R		
FILTER	003450R	FINSH	000622R	FIVE	002022R		
FI1	000000R	003 FI2	003720R	003 FI3	007640R	023	
FI4	013560R	003 FI5	017500R	003 FU	002070R		
FUU	002150R	FY	002164R	FYY	002244R		
F.ONE	002430R	GAIN	000010R	002 IDELAY	000004R	002	
IFTR	000000R	002 IFIRST	000002R	002 INDEX	002044R		
IO	004772R	IOF	005360R	IO.K	004770R		
IO.R2	005354R	MADD	006544R	MAXAD	000006R	002	
METHOD	000024R	MMUL	006102R	MSC	007070R		
MSUB	006716R	MTRN	007410R	MULFP	005572R		
NO.ELM	006534R	NXTA	006524R	NXTCOL	006526R		
NXTROW	006530R	N10	002030R	N100	002040R		
N12	002032R	N13	002034R	N15	002036R		
N2	002024R	N3	002026R	ONE	002014R		
P	001034R	PA	001200R	PACAT	001440R		
PAT	001200R	PC	=Z000007	PHI	001414R		
PONTR	005356R	PRINT	004636R	PRN	005266R		
PR.FI	004766R	PSEUDO	000474R	Q	001224R		
QA	001370R	QAT	001370R	QUIT	003724R		
R0	=Z000000	R1	=Z000001	R2	=Z000002		
R3	=Z000003	R4	=Z000004	R5	=Z000005		
SAMCOL	006532R	SHIFT	007554R	SIGN	004500R		
SP	=Z000006	STAPR	002434R	START	000000R		
STA.A	003402R	STA.E	002546R	STA.F	003076R		
STA.GO	003446R	STR1	003732R	STR2	004046R		
SUBFP	006020R	S.A	007630R	TEMBUF	004472R		
THREE	002020R	TMPADD	006712R	IMPDSC	007404R		
TMPMSC	007234R	TMPMUL	006540R	IMPSP	004470R		
TMPSUB	007064R	TPB	= 177566	TPS	= 177564		
TWO	002016R	T5.1	001750R	T5.5	001604R		
U	004214R	UPDATE	000640R	Y	004220R		
YOLD	001774R	\$ADR	= ***** G	\$DVR	= ***** G		
\$ECO	= ***** G	\$FI1	005476R	\$FI2	005512R		
\$FI3	005526R	\$FI4	005542R	\$FI5	005556R		
\$MLR	= ***** G	\$POLSH=	***** G	\$SBR	= ***** G		
. ABS.	000000	000					
	007632	001					
FTRN	000014	002					
DATA	023420	003					

```

0001      SUBROUTINE DIALOG
0002      COMMON /FTRN/ IFTR,IFIRST,IDELAY,MAXAD,GAIN
0003      CALL SETFIL (5,'A',IERR,'KB')
0004      PRINT 1
0005      READ (5,2) IFTR
0006      IF(IFTR.EQ.0) GOTO 100
0007      PRINT 3
0008      READ(5,4) IFIRST
0009      PRINT 5
0010      READ(5,6) GAIN
0011      PRINT 7
0012      READ(5,8) IDELAY
0013      PRINT 11
0014      READ(5,12) MAXAD
0015      PRINT 9
0016      READ(5,10) ZZZZ
0017      PRINT 77
0018      1  FORMAT(' DO YOU WANT FILTERING ?'/,
*        ' IF YES,TYPE 1; IF NO,TYPE 0'//)
0019      2  FORMAT(I1)
0020      3  FORMAT(' WHEN DO YOU WANT FILTERING TO START',
X        ' 1X,' ? [I3]'//)
0021      4  FORMAT(I3)
0022      5  FORMAT(' ENTER THE GAIN TERM FOR STOCHASTIC',
X        ' 1X,' APPROXIMATION. [F5.1]'//)
0023      6  FORMAT(F5.1)
0024      7  FORMAT(' ENTER THE MULT. FACTOR FOR THE',
X        ' 1X,' SAMPLING PERIOD. [I2]'//)
0025      8  FORMAT(I2)
0026      11 FORMAT(' ENTER THE MAX. NO. OF SAMPLES YOU',
X        ' 1X,' WANT. [I3]'//)
0027      12 FORMAT(I3)
0028      9  FORMAT(' THANK YOU. TO START,STRIKE ANY KEY '//)
0029      10 FORMAT(A1)
0030      77  FORMAT(///,5X,'PHI1',8X,'PHI2',8X,'PHI3',
X        ' 8X,'PHI4',8X,'PHI5',///)
0031      END FILE 5
0032      RETURN
0033      END

```

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