

MORSE CODE
COMMUNICATION AID
FOR
THE HANDICAPPED

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by

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ABSTRACT

A microprocessor driven display was built and programmed for the storage and reproduction of Bliss symbols. An explanation is offered for the success of the symbol language in teaching the handicapped.

The hardware was designed to be inexpensive enough for classroom use, but still deliver adequate flexibility and resolution. Due to the complexity and variety of the symbols a method of data compaction was developed to reduce the required storage space.

Initial tests are presented and suggestions are made for continuing the work.

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CHAPTER I

INTRODUCTION

1.0 Communication Aids for the Handicapped

The inability to communicate with other people can make life difficult and unrewarding. Since the most common form of communication is speech, lack of speech can be seen as a serious handicap.

Communication aids enable a person with a speech, visual, auditory or motor disorder to communicate with other people. Those with handicaps in these areas are often downgraded, as a lack of words is unfairly equated with a lack of ideas.

The purpose of this project was to provide a non-vocal cerebral palsied person with a portable means of communication using Morse code. An appropriate review of cerebral palsy can be found in Day (4). Briefly, it is a neurological disorder, present at birth and nonincreasing, which causes widespread motor disabilities. Its effects vary from slight tremors and weakness to a complete lack of control over voluntary movements, including speech.

In this case, the subject was a young man confined to a wheelchair and possessing extremely limited speech, but of

normal intelligence. He had some control over his feet and was able to turn the pages of a book and type slowly with them. For practical reasons of size and weight a typewriter could not be permanently attached to his wheelchair. An electronic keyboard with an eye-level display might have worked, but it would have still required a ruggedized QWERTY keyboard small enough to be unobtrusive, but large enough to be worked with the feet. For use outdoors either unit would be unsuitable as they could not be operated with a shoe.

1.1 The Morse Code as a Communication Aid

It was decided to try the Morse code as a method of communication. Unlike a typewriter, a single switch can be used to send the whole code, thus solving the problems of size and ruggedness. Morse code is fairly compact, with common letters such as E or T being a single dot or dash, while Z is much longer. This fortuitous arrangement is only valid for English.

The main disadvantage is that most people do not understand the code. It was felt that a simple and reliable decoding device could be built using microprocessor technology.

1.2 Previous Work and Basic Requirements

Schemes for the automatic decoding of Morse code have always been popular, but not necessarily cheap or effective.

Most of the research in this area is performed by amateur radio operators and various military organizations, each for their own purposes. Bell (2) presents the military view with its aim of increased surveillance with less manpower. For ham radio operators it becomes possible to communicate more enjoyably by removing some of the hard work. Recently there has been work on microprocessor based decoders, but they tend to require large volumes of code for initialization (1) or require fairly good code (12). The first requirement cannot be met due to the short conversations, and the second because of the nature of the sender. It was felt that a better algorithm could be developed for this use.

Unlike receiving, sending the code is fairly easy, and in fact the subject learned to send it within one week. By contrast, this author still cannot understand the code without the aid of the device developed here!

The items generated by this project were a footswitch interface to send the code, an adaptive software program to interpret it, and recommendations for the hardware required to produce a portable unit.

CHAPTER II

THE MORSE CODE

2.0 Basic Description

Morse code is an internationally recognized binary (on-off) coding scheme for the transmission of letters and numbers. Although it was originally designed to be sent and received by human operators, it can also be used with mechanical or electronic devices.

Each character consists of a unique sequence of one to six MARKS (on) separated by SPACES (off). The marks and spaces, or ELEMENTS, have several different lengths. Marks of time duration 1 are called DOTS, and those of lengths 3 are DASHES, SYMBOL SPACES of length 1 are used inside a letter, CHARACTER SPACES of length 3 separate characters, and WORD SPACES of duration 7 divide words.

The absolute timing of the elements does not matter as long as they maintain the 1:3:7 ratios. The dot is usually designated as the speed determining element. By not fixing the speed, operators can work from 5 to 60 words per minute (Guenther,9), often limited by skill of the receiver. Machine to machine transfers can proceed at higher rates.

Grappell and Hemenway (8) state that the reciprocal of

the speed in words per minute corresponds roughly to the dot length in seconds for normal text. Thus, at an average speed of 10 words per minute, the dots and symbol spaces would be 0.1 second long.

Bell (2) was unable to find a standard value for the bandwidth of Morse code, but used 3 times the reciprocal of the dot length as a working value. This figure is important when designing analog input circuitry to reduce noise while retaining information. The bandwidth of machine code could be rigorously calculated, but hand sent code is too dependent on human factors for meaningful results.

2.1 Code Sample

To send a message, such as "I am hot", it is necessary to find the correct code for each letter, and then join the individual codes with the correct spaces. A table of the various character codes is given in Appendix A. There is no provision for upper or lower case in the code.

Individual codes:

I	. .	a	. -	m	- -
h	o	- - -	t	-

Complete message:

I	a	m	h	o	t	(text)
· I	· a	· m	· h	· o	· t	(code)
i i	i 3	3 3	i i i i	3 3 3	3	(mark times)
1	7	1 3	7	1 1 1 3	1 1 3	(space times)

It can be seen that information is contained in both the marks and spaces.

2.2 Decoding

To decode a message it is first necessary to take each element and decide its relative length: 1, 3 or 7. For well sent, low noise code, the wide ratios (1:3 and 3:7) make this an easy task, even for a mediocre operator or simple machine.

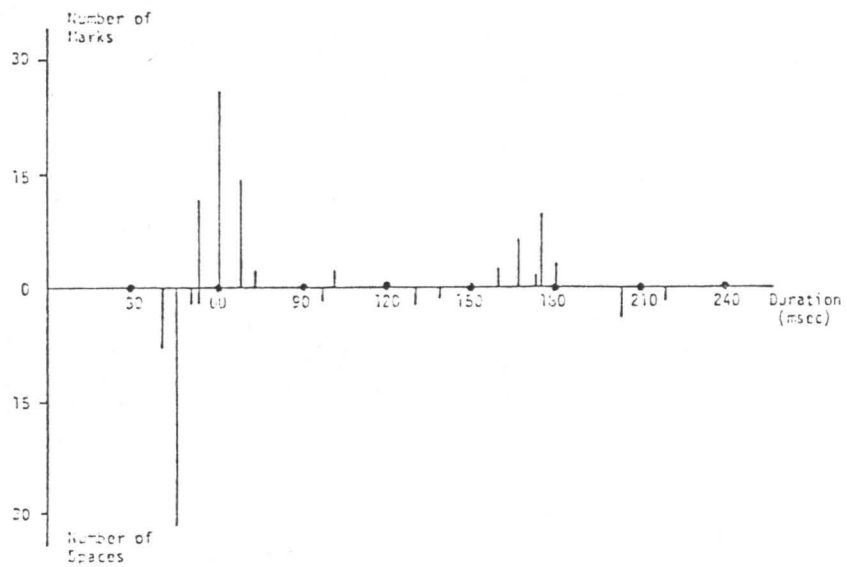
Because the spaces clearly define the beginning and end of letters and words, it merely remains to look up the individual letter codes and produce a copy of the message.

The above tasks can be implemented easily to produce a simple decoder at low cost. The Morse-A-Letter (11) is an example of such a device, available as a kit for \$150.

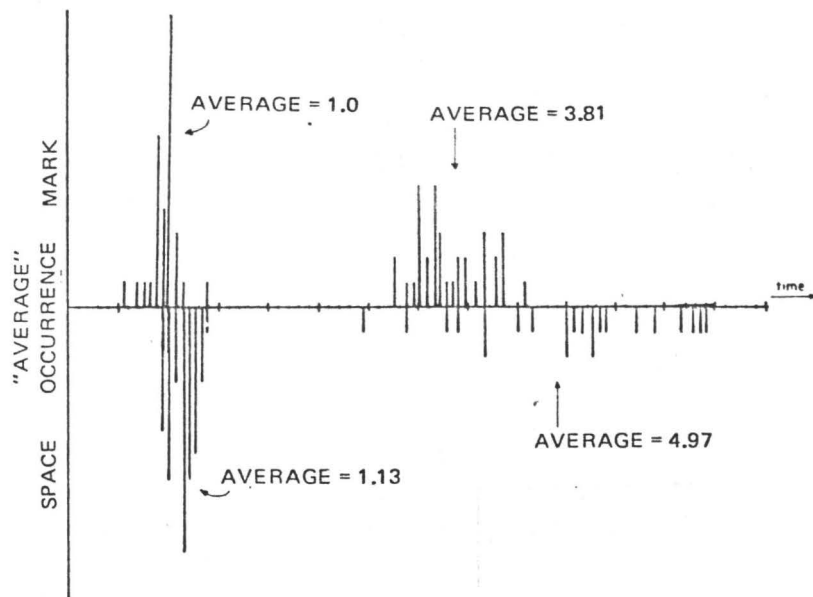
Unfortunately, Morse code is rarely received under ideal conditions. Human operators cannot exactly duplicate the 1:3:7 ratios, and during a long message both the ratios and overall speed will vary. Bell (2), Guenther (9), Freimar (5) and other detail some of these changes, including the variations which may occur inside a single character.

For most applications the signal has been received via a radio link, and is further corrupted by noise and fading.

For reliable decoding of Morse code, the machine must be adaptive in some way. It may only adjust its overall



a) Bell, Reference 2.



b) Hickey, Reference 10.

Figure 1. Histograms

speed, such as the Automatic Fist Follower (7), it may compensate for deviations in the 1:3:7 ratios (3), or it may even perform some textual analysis based on a rudimentary knowledge of English (10).

2.3 Code Characterization Displays

In order to analyze a Morse signal, most researchers have found it helpful to have a graphical representation of the main characteristics.

Hickey (10) and Bell (2) used histograms, shown in Figure 1. The horizontal axis is time, and the positive vertical axis shows the number of marks with that duration. The negative vertical axis displays the space information, inverted for clarity. In both cases, some of the important characteristics of Morse code are shown:

1. There are more dots than dashes, and more symbol spaces than character or word spaces.
2. The dot cluster is narrow with an obvious midpoint.
3. The dash cluster is wide, and the midpoint may not be clearly defined (flat top, rather than peak).
4. The symbol space cluster is as well-defined as the dot cluster, with nearly the same midpoint.
5. The character spaces are poorly grouped.
6. The word spaces are not grouped at all.

Point 1 is inherent in any Morse transmission of text,

but 2 to 6 are peculiar to hand sent code. It can be seen that the 1:3:7 ratios are not exactly followed.

Guenther (9) used scatter distribution plots which are able to show the changes when one element follows another. For the example shown in Figure 2, it can be seen that dashes following word or character spaces (groups D and E) are longer than dashes after symbol spaces (group F). This particular type of information cannot be easily seen on the histogram displays. The frequency of occurrence information is slightly obscured by the scatter plots as it changes from height to dot density.

For this project, the histogram method was chosen as it offers a good display of information with minimal hardware and software requirements. The output was first presented on a Teletype printer (Figure 3), but the program was changed to use a Tektronix graphics terminal (Figure 4). This gave a higher quality display and much faster output.

2.4 Decoder Functions

Any decoding device must perform several basic functions:

1. Separate the marks into dots and dashes.

This is the least difficult operation as the dots and dashes are usually well sent, preserving the 1:3 ratio or even making it larger (3).

DOT & DASH (ALL)

RECORDING SESSION 2

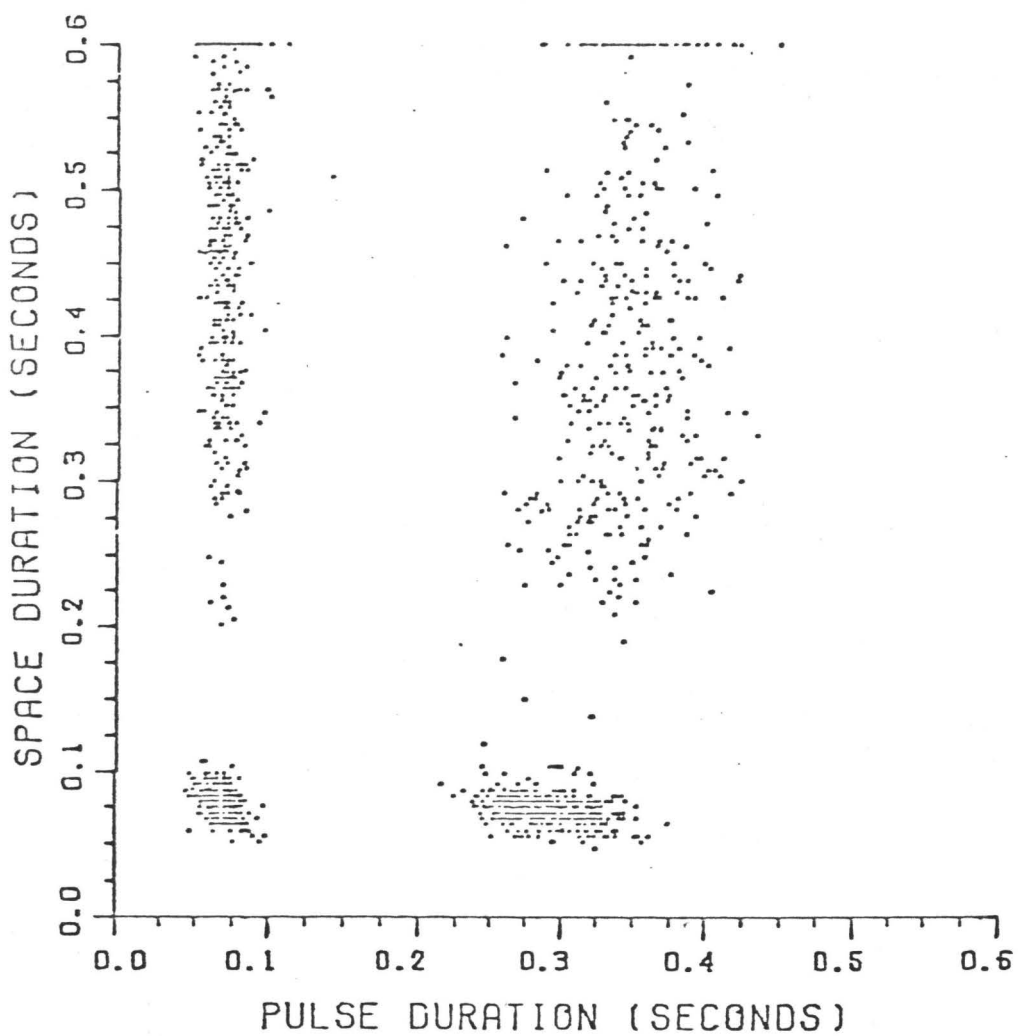
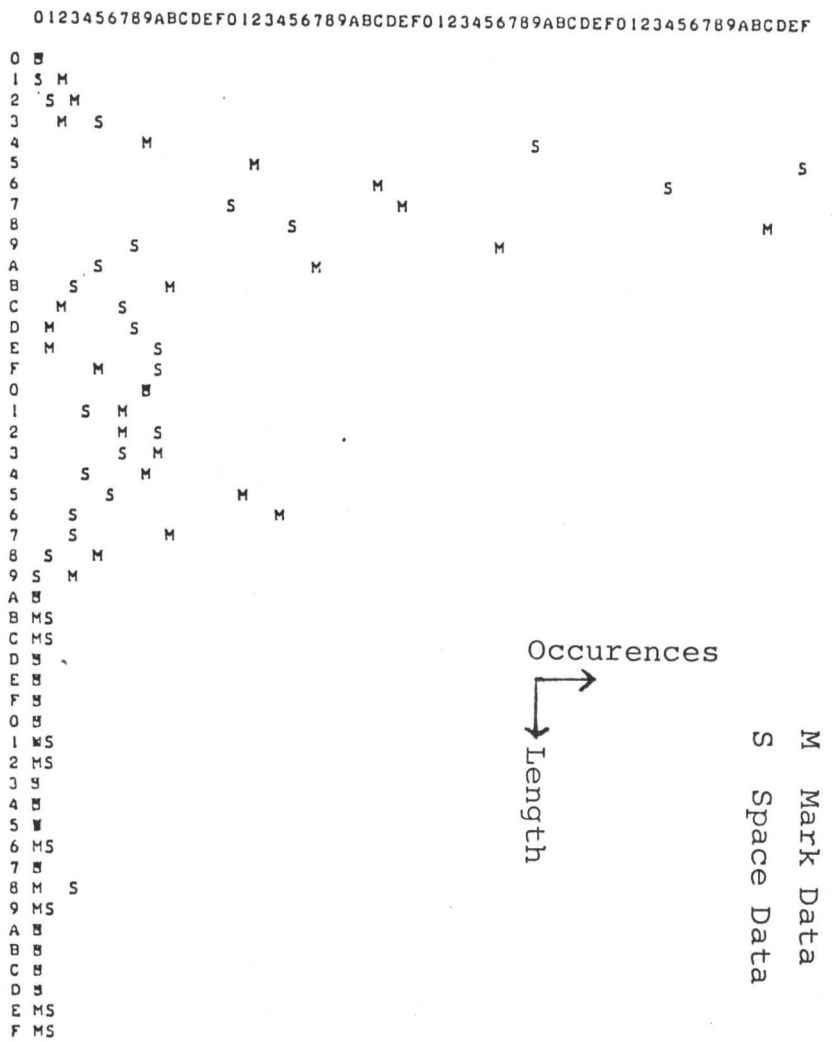


Figure 2. Scatter Plot. (Guenther, Ref. 9)

Figure 3. Teletype Histogram



SPEED 0236

MARK THRESHOLD 07

SPACE THRESHOLD 09

Computed Parameters

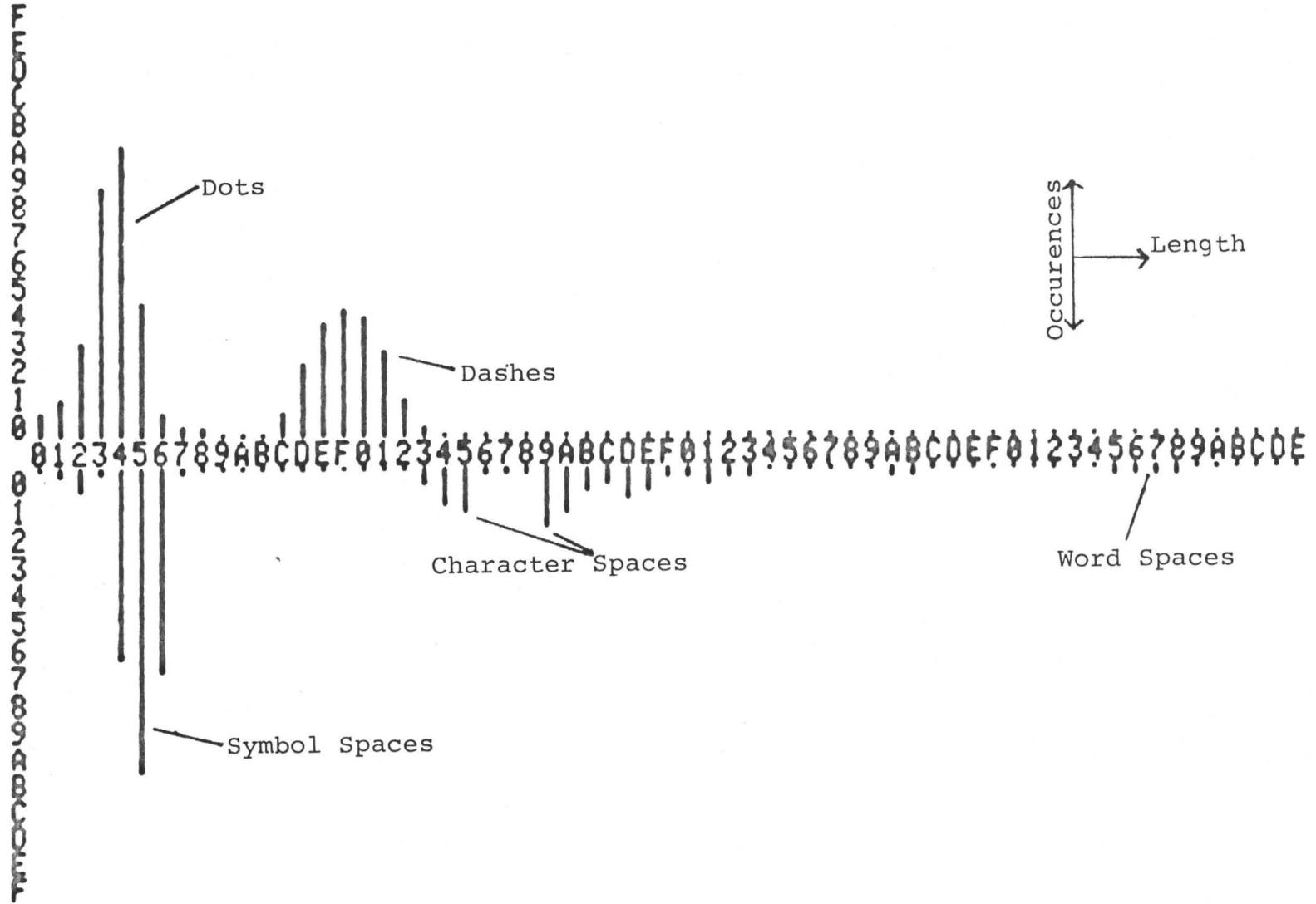


Fig. 4 Annotated Tektronix Histogram

2. Separate the spaces into symbol, character and word. The spaces are never sent with as much care as the marks. Blair (3) and Guenther (9) agree that a more complicated decision process is needed for the spaces.

3. Based on the character spaces which have been found, divide the data stream into characters, decode and output them.

If the character spaces have been correctly found, and if the marks were properly identified, the final conversion from Morse to text is a simple operation using a look-up table. When the above processes produce an invalid character, there is a choice between printing an error message, and reassigning the various elements using a modified decision technique in an attempt to find a valid character. This correction process may be quite involved in the larger devices (11).

4. Based on the word spaces, output spaces between words.

Although word spaces are the most variable type of space, their identification is not critical, as a missing space between words rarely causes a loss of understanding.

2.5 General Decision Strategies

Hickey (10) noted that there are three basic approaches to the manual Morse problem:

1. Macro: You can accumulate statistical information on an

operator and use this data to make decoding decisions.

2. Micro: You can make your decisions on a mark to mark basis, sometimes called the "idiot dot" method.
3. Hybrid: You can compromise these two methods and come up with a hybrid algorithm.

Method 1 requires a memory large enough to contain all the timing information in a message, but should give the best results. Method 2 requires very little computer time or space, but is the least reliable. Successful devices, such as MAUDE (6), use the hybrid approach.

A common method is to keep an average speed parameter which varies slowly, and an average dot time which uses only the last few characters. The decision thresholds are calculated from empirical formulas using the dot average. The short "time constant" on the dot average allows tracking of fast speed changes, while the slower responding speed parameter inhibits locking on to noise (8).

2.6 Decision Problems and Approaches

1. Decoding Delays:

The methods which build up a large body statistical information about an operator require a large sample of code to be effective, thus delaying decoding. Guenther (9) tried to quicken the process by decoding initially with one set of rules, then switching to a better set when more code

had been received.

For use as a conversational communication aid, the long initialization times of most methods are unacceptable. Since the code will be sent very slowly (possibly 1 w.p.m.) it is important to have each letter decoded as quickly as possible.

2. Non-Code Activation:

Due to the nature of the sender, there are periods when the code switch is involuntarily activated. If some of the simpler averaging techniques are used, after several non-code "words", the program will have locked on to that style, losing its memory of the sender's true characteristics. This is similar to the problems caused by noise and interference in the radio case, and is why "idiot dot" strategies fail.

Since a non-code sequence is obvious to an observer, (because of the sender's physical involvement) no attempt was made to block decoding during one. It would require a complex program to have the computer recognize invalid input. The decoding algorithm was designed so that inputs which are not similar to previously received code do not affect the decoding thresholds unless prolonged.

3. Space Problems:

Problems arise when there is poor distinction between the symbol and character spaces. If a long group of marks

is received with no obvious character space in the string, a good receiver can separate the string into characters using his knowledge of valid letter codes, combined with his expectations based on context.

Machines find this difficult unless they contain rules and vocabulary from the language being decoded. An easier method is to mix timing and simple language rules such as in the MAUDE decoder. Rules such as "The longest of 6 successive spaces is almost always a character space" allowed the machine to divide a string of poor code into letters. If the division resulted in an invalid character, a modified set of rules was tried. (6)

The bad grouping of the various space types not only causes incorrect decoding (due purely to the bad sending), but also makes it difficult to calculate best values for the space thresholds, compounding the errors. In this program the threshold setting process does disregard poorly grouped information as mentioned above, but in addition the initially calculated symbol/character threshold is averaged with the dot/dash threshold, which is more reliable. For perfectly sent code these two thresholds should be equal.

4. Correction Routines:

McElwain and Evens (11) reported good success with a "degarbling" program to correct machine received code. It was given a vocabulary list of every word which might be

used in the message, along with the received, partially decoded text, and some timing information. Based on this it could correct many of the errors in the received text. This would be useful in a military situation where the context is unknown but the vocabulary prescribed. It allows large volumes of readable text to be produced with little operator attention.

In ordinary conversation the vocabulary is large, but the context is known - such as replies to questions. It was found in preliminary tests that a reply never had to be completely spelled out, as the listener could correctly guess the full statement after a few letters or words. For this reason no error correction routines were included in this design. This works well in a conversational environment, but might be inadequate for other uses.

5. Effects of Incorrect Decisions:

Most of the schemes examined use variations of the following method. The element being processed (on line or from memory) is examined and a decision made as to its nature (dot/dash, symbol/character/word). This decision is used as part of the character currently being assembled for decoding. If the element was classified incorrectly, the current character will be incorrect. Error correction schemes may help if an invalid character was produced.

One incorrect character is not a problem, but each element is then used to update a running average, as in:

$$\text{DOT AVG.} = \text{DOT AVG.} + \text{NEW DOT}/8 - \text{DOT AVG.}/8$$

$$\text{DASH AVG.} = \text{DASH AVG.} + \text{NEW DASH}/8 - \text{DASH AVG.}/8$$

(This keeps running averages of the typical dot and dash, with each new element given a weight of 0.125.)

The various averages are used to calculate the thresholds, as in:

$$\text{MARK AVG.} = \text{DOT AVG.}/4 + \text{DASH AVG.}/2$$

(For perfectly weighted 1:3 code, this produces a threshold at 1.75, an empirically better value than value 2, which might have been expected.) (after 1)

Unfortunately, any incorrectly classified element is averaged into the wrong place.

The effect of this is that a period of bad code or noise is not only incorrectly decoded (which may be unavoidable), but is also used to calculate incorrect thresholds which prolong decoding errors into periods of good code.

Most of the element classification errors occur when an element is received with a timing value near a threshold value. These elements are poorly grouped with respect to the average of their intended value. If these elements are included in the average calculations, even if they are correctly classified, poor threshold values must result.

Elements near present threshold values cannot however be explicitly and permanently excluded from calculations as they may be important if the operator has changed speed.

The decoding algorithm outlined in the next chapter tries to use only well-grouped data, ignoring data which may represent noise or odd code.

CHAPTER III

THE DECODING PROGRAM

Some of the main features of this program are:

3.0 Timing

A software timer is used instead of external counters. This gives maximum adaptability and minimizes the hardware requirements. At the slow code speeds anticipated, the non-timing portions of the program do not significantly distort the timing, so the delays through various paths were not equalized, but ignored. A single speed parameter is used, with an estimated useful range of 300 to 1.

All timing in the program is in the arbitrary, dynamically variable units determined by the speed parameter.

3.1 Input

The code input is sampled under program control, instead of using an interrupt. It was felt that an interrupt based system was too susceptible to noise, as it must respond to every input change, however short. When fully adapted to a code style, an average length dot is sampled 8 times, a dash 24 times. While adapting, the program will decode correctly if a dot is sampled from 2 to 20 times.

The program will lock on to code over a much wider range.

3.2 Histograms

As each mark or space finishes, its length (in arbitrary timing parameter units) is stored in a histogram type table. There are separate tables for marks and spaces. The tables allow for a length from 1 to 64 timing units, with a maximum count in each entry bin of 64 occurrences. When any bin in either the mark or space table reaches 64 counts, the data in both tables is divided by 2. This prevents software overflows and allows new data to slowly take over from the old. The histograms are smoothed with a moving window and used for threshold calculations.

The smoothed data can be displayed on a graphics terminal as an actual histogram. This is useful for checking the overall quality of the code and the validity of the threshold setting algorithm.

3.3 Morse/ASCII Conversion

As each element comes in, it is shifted into a character holding register, using "0" for dots and "1" for dashes. When a character space is found, the data in the character holding register is used directly as a memory address pointing to the equivalent ASCII in a look-up table. This method is simple and fast while requiring a minimum of space for the table.

Two 6 mark characters produce ambiguous addresses and are dealt with separately.

If it is desired to put out more than one character per Morse input, such as a "CR/LF", then a special case subroutine must be written.

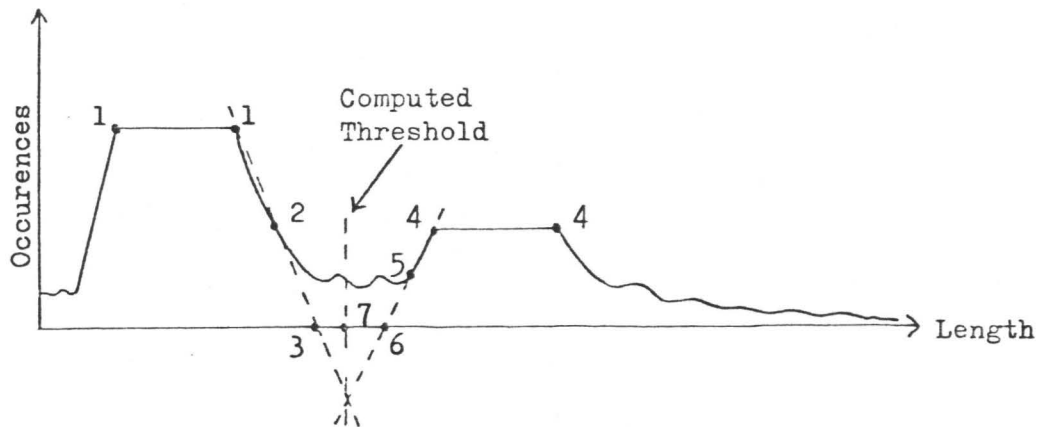
3.4. Slow Adaptability

After the thresholding process has decided dot or dash each mark time is compared to 8 or 24 respectively. If the time was short, the overall speed is increased by a small increment, and vice versa. This action slowly forces the program into the correct speed range. As it works slowly, noise or bad code will not really affect the speed unless prolonged for dozens of marks, an unlikely event. Due to its poor quality, the space data is not used.

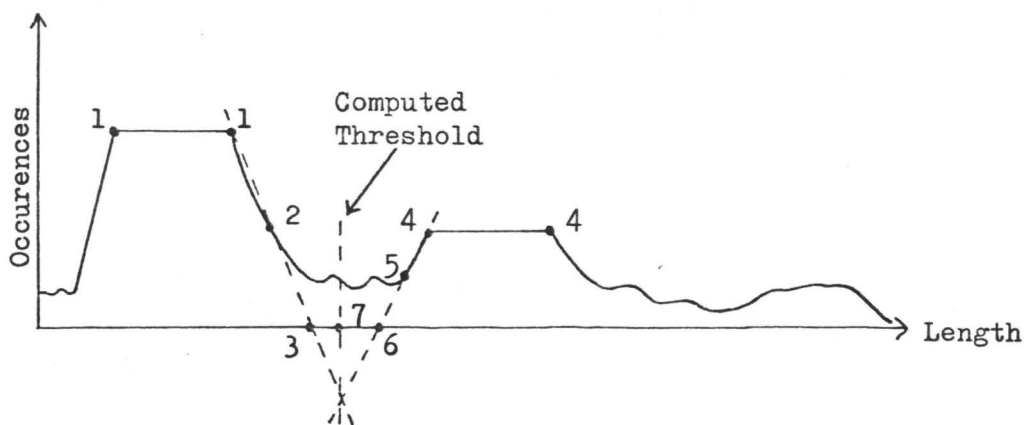
Initially this was the only adaptation mechanism in the program, but it could not be made fast-responding without the program becoming unstable during bad code. It also required perfect 1:3 ratio code, which is rarely sent.

3.5 Threshold Setting

Figure 4 shows a typical histogram and 5 shows an idealized version. If point 7 can be reliably found, then the critical dot/dash and element/character



a) Marks



b) Spaces

Figure 5. Threshold Histograms.

decisions can be properly made.

Blair (3) used histogram data, by assuming an initial threshold value, then calculating an empirical "goodness of separation" based on various statistical parameters. The threshold estimate was moved until the "goodness of separation" was maximized. This method requires a lot of computation and would not be suitable for real time application, but the basic idea is excellent.

The stylized histogram in Figure 5 provides the basis for the method chosen. It was found that code from different sources always had similarly shaped histograms. The most important similarities are that the dot peak is consistently higher than the dash peak, and the symbol space peak higher than the character space peak. The work space peak is not well-defined and often cannot be found. The "noise level" between peaks is lower than the peaks, and consists of elements which were poorly timed. The peaks are shown with flat tops, but this is not essential.

To find the threshold (dot/dash or symbol/character):

1. Find point 1:

An entire histogram is searched for the largest value, point 1. If two points have this value (1 and 1'), then the rightmost one is taken because it should be closer to the threshold. This largest peak is always the dot or symbol space peak. The position and value of point 1 are noted.

2. Find point 2:

The histogram is searched from point 1 to the right for the data point which is $1/2$ or less than the value at point 1. If this point does not occur before the dash peak the code is of very bad quality and cannot be decoded. Only its position is noted, as its value was used to find it.

3. Find point 3:

The first estimate of the threshold is found by calculating point 3, which is the position as far to the right of point 2 as point 2 is to the right of point 1. Its value is ignored, because it contains the bad data that this method was designed to avoid.

4. Find point 4:

The histogram is now searched for the largest data value from point 3 to the end. This will be the dash or character peak. If two bins contain this value the leftmost one is chosen as this will be closer to the threshold. The position and value are noted.

5. Find point 5:

The histogram is searched from point 4 to the left for the first bin whose value is $1/2$ or less than that of point 4. If the code is very bad the noise floor could be higher than this $1/2$ value, which would extend the search to the left side of the first peak. To avoid this, the search will stop at point 3, even if the desired value was not found. Code

that bad would be unreadable anyway, but a good threshold value must still be chosen to aid the slow adaptation. Only the position of this point is used.

6. Find point 6:

The second estimate of the threshold is found by calculating point 6, which is as far to the left of point 5 as point 5 is to the left of point 4. Just the position is taken as the value contains questionable data.

7. Find point 7:

Point 7 is the arithmetic average of points 3 and 6. This is the estimate of the threshold. If either line 1-2-3 or 4-5-6 has a shallow slope it is possible for point 3 to be to the right of point 6. This does not affect the choice of point 7, as it just means that 1-2-3 crosses 4-5-6 above rather than below the horizontal axis.

8.8. Space Correction:

The above seven operations are performed separately for the mark and space data, so that the only tie between the thresholds should be the overall program speed. This allows for different distributions of the marks and spaces. It also is useful for very slow code (less than 1 w.p.m.) so that long spaces can be used without long marks.

It was found however that the space data is never as well-grouped as the mark data. This becomes a problem when small code samples are obtained, as it may take a long time

to build up the space peaks properly. For this reason the space threshold estimate obtained from steps 1 to 7 is averaged with the value for the mark threshold to obtain the final space value. It was found that this helped decoding in almost all cases.

9. Character/Word Threshold:

The space between words in regular Morse code is 7 dot times long. These word spaces should appear as a third peak in the space histogram, but rarely do, so a continuation of the above methods cannot be used to find it.

The character/word threshold is set at 3.5 times the symbol/character threshold, or 7 dot times (assuming a symbol/character value of 2 dot times). Normal code would give a value between 3 and 7, but it was found that for this application the character spaces were as long as ordinary word spaces. This operation is not critical for conversational use.

The threshold setting method has been shown to find a good threshold value. No math other than integer addition, subtraction, and division by shifting is used, giving a high speed. It can be seen that a simple curve fit is actually done to find the horizontal position of the intersection of lines 1-2-3 and 4-5-6, which should be at the low point between the two distributions. Only very basic assumptions are made about the shape of the histograms,

without fixing the time ratios at 1:3 or the valley at a certain distance between peaks. Any time bins which do not have counts to a value of at least 50% of the closest peak are automatically excluded from the calculations, as are all bins on the far sides of the peaks which also do not contain information about the threshold position.

The most important objective has been achieved, and that is any elements which were incorrectly classified when first received are not averaged into the wrong place to perpetuate the initial error. Instead they are put into the histogram where they belong in time-relation to previously received elements. If they were part of a legitimate speed change, new peaks will form in the correct place to give a correct threshold, and if they were errors of low occurrence they will be completely ignored.

If reasonable quality code is received, a useable threshold is usually found after the receipt of less than a dozen marks, giving the required fast initialization time.

3.6 Out of Range Correction

If the program is running far too quickly, then the dashes will time out to the maximum of 64, rather than the ideal 24. If a mark of length 64 is detected, the timing constant is doubled (to slow the program and halve the counts) and the program is restarted with clear histograms.

This allows one to start without having to make an initial guess at the code speed. The program will restart on each mark until the timing constant is within range. This mechanism can be fooled by holding down a key to produce a long mark, so an automatic keyer should be used which cannot produce marks longer than those required. Spaces are not checked in this manner as a long space could be just a pause in sending, not a change of speed.

If the program is far too slow, a mark may only be sampled once. There is no explicit mechanism to deal with this case, as a typical key bounce would also be sampled once and a speed halving mechanism (as above) would cause the program to lock onto that noise. The normal threshold setting process can deal with this condition, for if all the marks go into the 1 bin rather than being spread around 8 and 24, the histogram fills up and empties quickly. This lets the slow adaptation and threshold mechanisms function effectively by quickly removing data stored at an earlier speed. This is a slower correction than for the opposite case above, but it is stable unless the number of noise marks greatly exceeds the number of true marks. Code of such bad quality is not expected in this application.

Flowcharts are given in Appendix B and the assembled code is in Appendix D.

CHAPTER IV

HARDWARE - AIMS AND DEVELOPMENT

4.0 Hardware Requirements

The hardware falls into three main categories:

1. Patient Interface - physical movement to audible code
2. Digital Interface - audio signal to digital input
3. Processor and Display - digital input to visible letters

The ideal system would consist of a single switch connected directly to an input port on a one chip microprocessor. An output port would drive a speaker with clean code for audio feedback, while another drove a display with the reconstructed characters. Very little other hardware would be required, as all the timing and switch debouncing would be performed in software. The computer and the display would be lightweight and low powered. The system used for development was more complex, but at all stages decisions were made with the above criteria in mind.

4.1 Patient Interface

In any manual Morse system a key is required to translate the operator's movements into a string of dots and dashes, providing an electrical signal for transmission and

an audible one for the operator. The audio feedback may come from a directly connected buzzer, monitoring of the transmitted signal, or the mechanical action of the key. Keys vary from a simple switch to dual switches with mechanisms for automatically timing perfect dots, dashes and spaces.

Normally code is sent with the hand, but in this case the subject only had reliable control over his right foot, so this point was chosen for the interface. During the largest muscle spasms it was unuseable, but good action was available at most times.

At an early point in the project, it was decided not to use a direct electrical connection between the subject and the processor. In its place the audio signal from the keyer was picked up by a microphone and conditioned for use by the computer. This approach allowed all test sessions to be recorded on a cassette recorder in any location suitable to the subject. Whether the code was "live", taped in the subject's home, or from a shortwave radio, it could all be processed the same way and stored easily for reevaluation. It also solved the problem of electrically isolating the subject from the equipment, which was line-powered except for the keyer.

One final benefit was that verbal comments on the progress of the session were recorded at the appropriate times

and were thus more valuable than written notes. This was especially useful in analyzing the earlier sessions, as the character sent could rarely be recognized but was often framed by comments which described the true intent and the mistake.

The first key to be tried was a "Rancho"¹ footswitch switch connected directly to a 12-volt battery and a Mallory Sonalert (Figure 6). The Sonalert was chosen because of its low power consumption, high efficiency and sine wave audio output. The last quality was found to be useful in triggering the phase locked loop detector described below. Ordinary buzzer sources produced complex audio waveforms which were hard to detect against background noise.

The switch could be placed in a loose-fitting shoe or sandal and was cosmetically the most pleasing, but it had a large amount of contact bounce. This was not a problem in gait studies due to the higher pressure available for contact closure, but the effect was intolerable for machine read Morse code.

It became apparent that the subject was making a valiant effort to send correct code, but he could not time out the dots and dashes properly with the single switch. It was decided that a dual action switch was needed, with

¹Developed at Rancho Los Amigos for gait studies

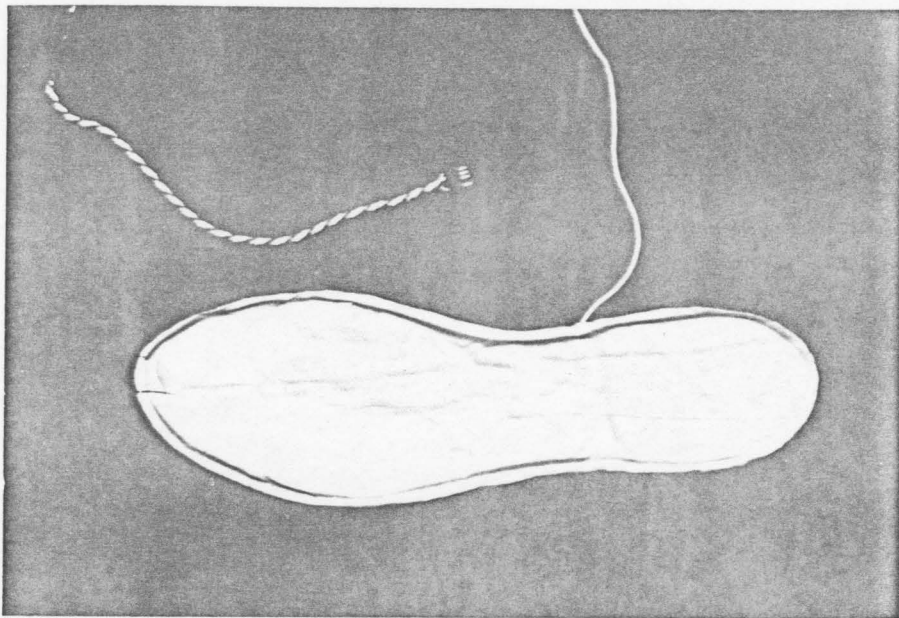


Figure 6. Rancho Footswitch

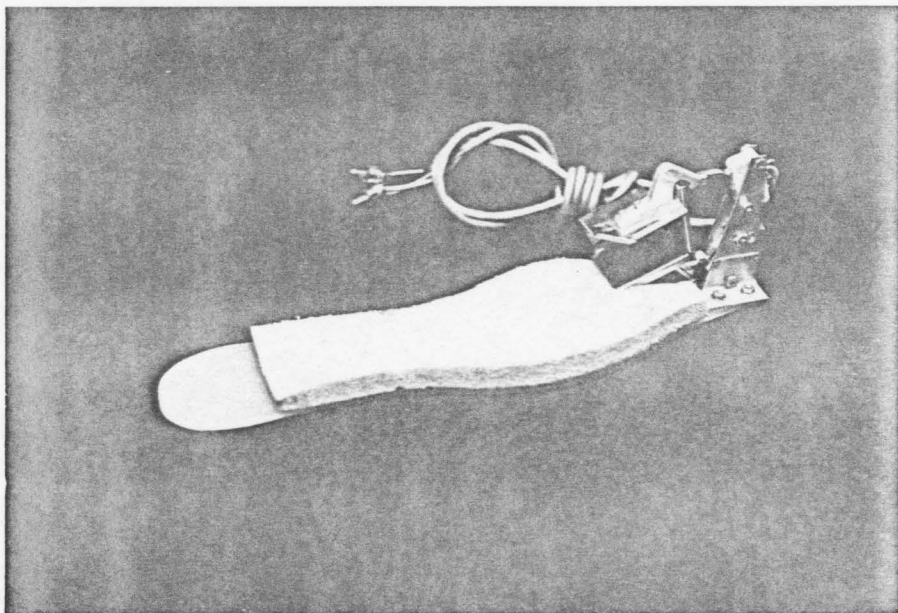


Figure 7. Metal Footswitch.

one position for dots and another for dashes. The author designed and built the footswitch in Figure 7. The metal loop enclosed the big toe of the right foot and pivoted in a similar way. Two microswitches sensed motion up or down from the rest position to trigger a dash or dot respectively. This orientation was chosen as the subject could press down more easily and more dots are used than dashes. The restoring force from each direction could be varied independently by changing the microswitches.

To complement this, a keyer was designed by Tony Wallace² with two separate inputs which would result in an audio tone of dot or dash length when activated.

Further trials were conducted with this equipment and new problems appeared. The footswitch required the foot to remain motionless while the toe pivoted up and down. The subject however found it easier to curl his toes down and push his foot up and down, which did not activate the switch properly. The switch also had to be flexible like the Rancho switch or else much stronger, as it was easily damaged.

The keyer had no automatic space delays, so although perfect length marks were being sent, a dot and dash could be sent simultaneously. Holding a switch down produced one mark, but another often appeared on release due to contact bounce. The square wave oscillator and small speaker pro-

²Electronic Technologist, Chedoke Rehabilitation Centre

duced an excellent signal for the ear, but it was not recognized over background noise by the equipment. The length of the tones, and hence the overall speed, could not be easily varied.

Another footswitch was designed and constructed. (Figure 8) Two microswitches were mounted into the wooden frame and activated by pressing the clear plastic cover plates. This design was more rugged and easier to activate.

4.2 Automatic Keyer

The author designed and built a fully automatic keyer which can send perfect dots or dashes, but with the addition of a perfect symbol space between them. If the footswitch is held down the keyer will produce a string of dots and spaces. If it is then suddenly pushed up, the dot in progress is finished, the symbol space timed out, and the switches sampled. If the dash switch is still activated, a dash is begun and the switch can then be released. The dashes are also auto-repeating if required.

There is a semi-automatic mode in which a switch has to be released and reactivated before a mark is repeated, although the minimum space timing is still produced.

In the manual mode either switch activates the tone for as long as it is closed with no automatic timing.

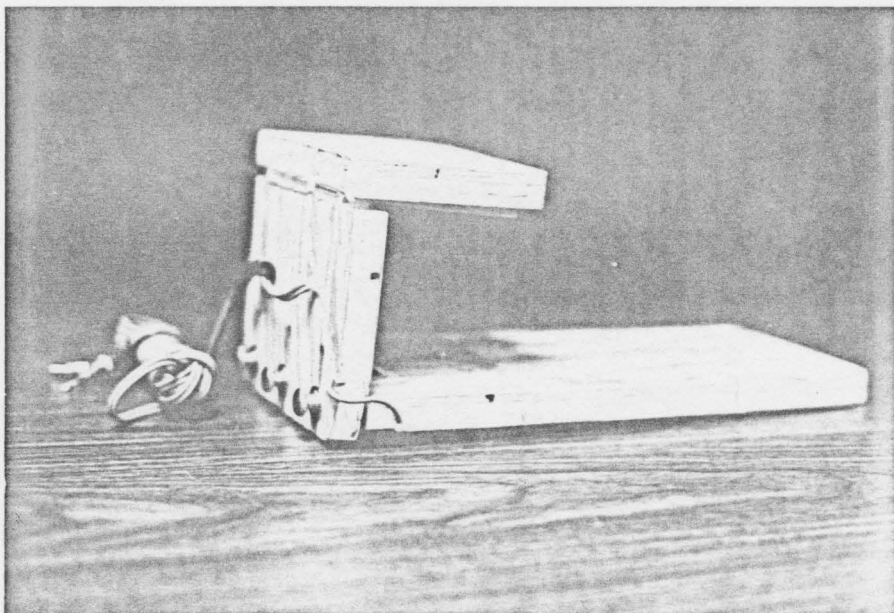


Figure 9. Wooden Footswitch.

Since the design is fully synchronous without using monostables, the speed can be varied over a wide range with all timings (dot, dash, space) in the correct ratio. The Sonalert was used again for the advantages outlined above. The unit was built with CMOS technology and has a quiescent power drain of microwatts. For this reason and to simplify operation no power switch was included. The battery can be from 4 to 9 volts with no change in operation.

4.3 Detailed Keyer Operation

In the manual mode either switch can directly activate the output through IC1a, D3 and Q1. The trigger signal to the timing circuits is shorted to ground through D1 and SW1b.

In the automatic mode IC6 is the timing oscillator which sets the code rate. IC5 is a decade counter with 10 active high decoded outputs, clocked by IC6. If no switches have been closed:

1. IC5 is in the "0" state, closing the trigger sampler formed by IC1b and IC1c.
2. IC3a is off turning off IC6, the clock.
3. IC3b is off, turning off the output.

When a switch is closed, the trigger signal gated through IC1b and IC1c:

1. Sets IC3a to turn on IC6 and start the clock.

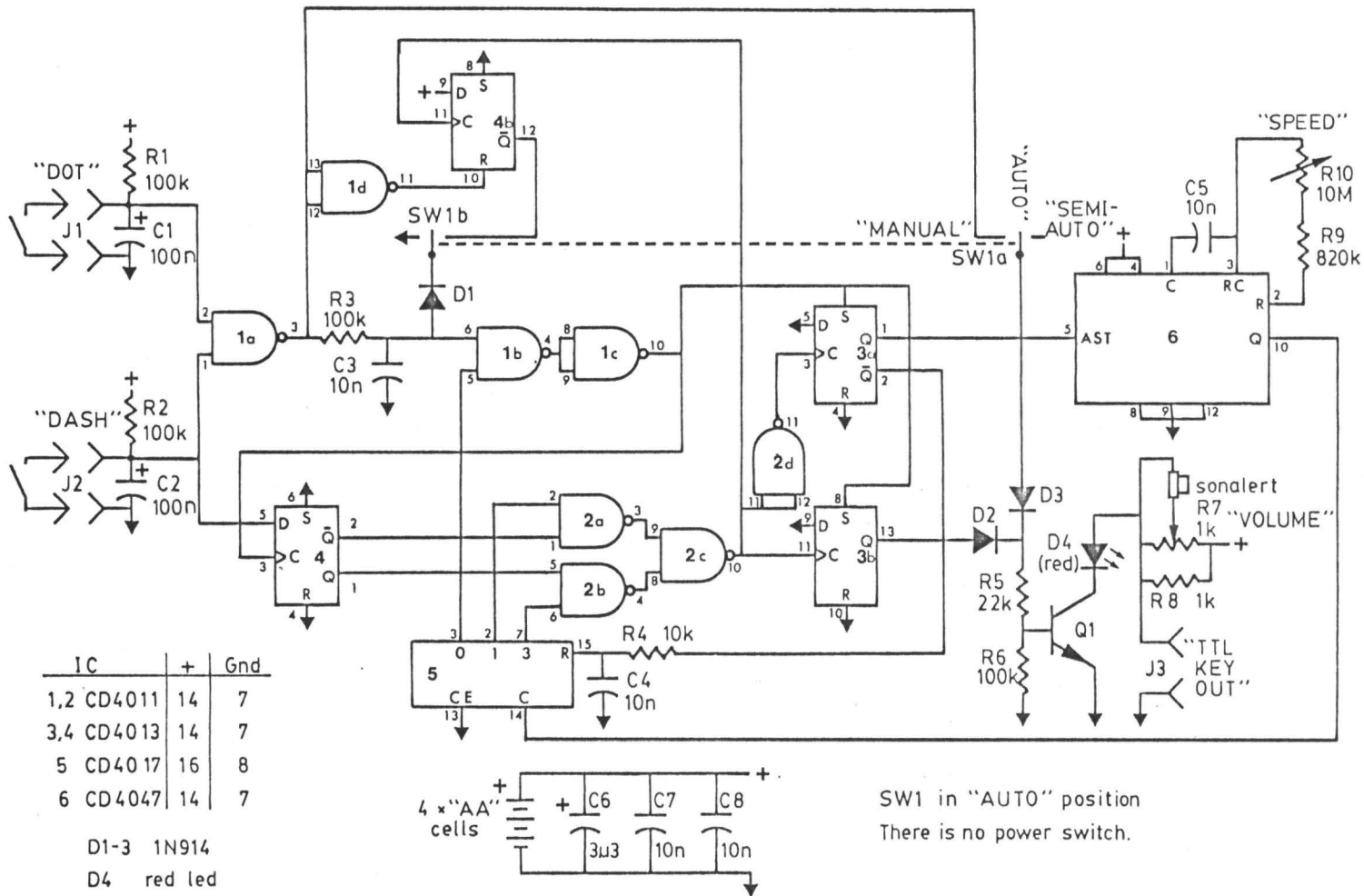


Figure 9. Keyer Schematic

2. Sets IC3b to turn on the output through D2 and Q1.

3. Clocks IC4a to latch the input - 1 for dash, 0 for dot.

Because the input is latched, a change in the switches will not change the timing.

IC2a, b and c select the count from IC5 which will stop the output, based on the switch latched in IC4a. A dot is stopped by the "1" count, a dash is stopped by the "3" count, giving the correct timing.

Using a dot as an example:

1. The output is on during the "0" count after triggering
This is one clock cycle long.
2. The rising edge of the "1" count, gated by IC2a, b and c, clocks IC3b to the zero state, turning off the output and ending the dot.
3. During the "1" count, the output is off and the input is locked out. This gives the required symbol space.
4. One clock cycle later, the falling edge of the "1" count clocks IC3a to the zero state, which turns off the clock and resets IC5 to the "0" state.
5. As IC5 is reset to "0" the switch sampler of IC1a and b opens again to check the switches, ending the cycle.

The dash sequence is identical, except the "3" count is used instead of the "1" count.

In the semi-automatic mode, the signal which ends the mark also clocks IC4b to the "1" state. The trigger signal

is then grounded out through the \bar{Q} output and D1. When both keys have been released, IC4b is reset through IC1d. The effect of this is that a switch must have been released after the completion of the mark (i.e. during the symbol space or later) and then reclosed to send another mark.

4.4 Footswitch/Keyer Results

Sequences such as "dot dash dot" (R) became "down up down" with little regard for timing. The only requirement was that a switch had to be held until the mark began, and the switch for the next mark had to be activated (or the last one released) by the time the previous symbol space was finished. Due to the very slow speeds available and the fact that the switches were locked out when not being sampled, good "noise immunity" from involuntary movements was achieved.

This equipment was tried and the only improvements needed were a reorientation and height adjustment of the upper switch, but its overall performance was finally good enough to test the software.

A final footswitch was built incorporating the needed adjustments((Figure 10) and it performed satisfactorily.

4.5 Audio Input Board

The basis of this circuit, shown in Figure 11, is IC3, a phase locked loop tone detector. This approach was chosen

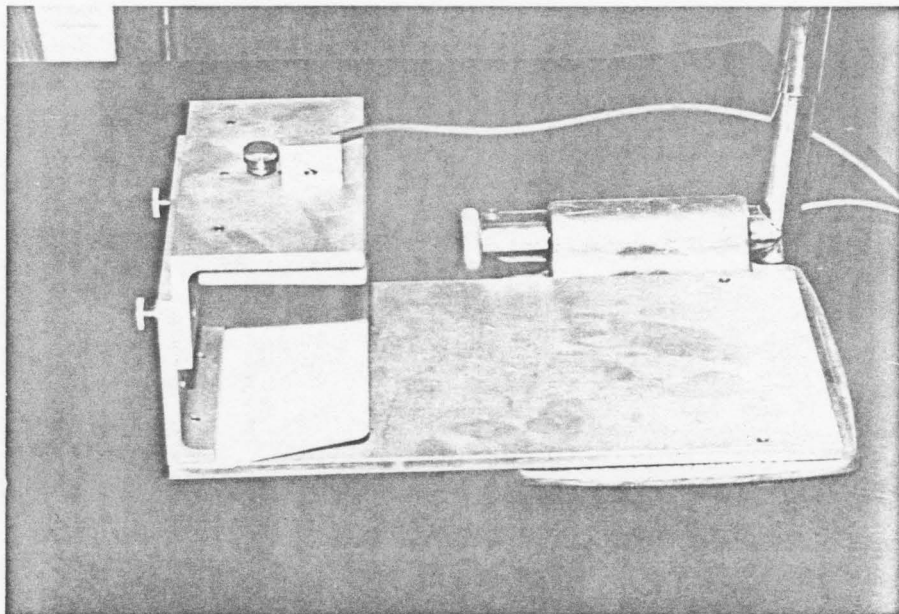
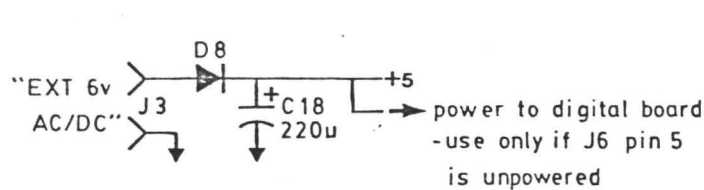
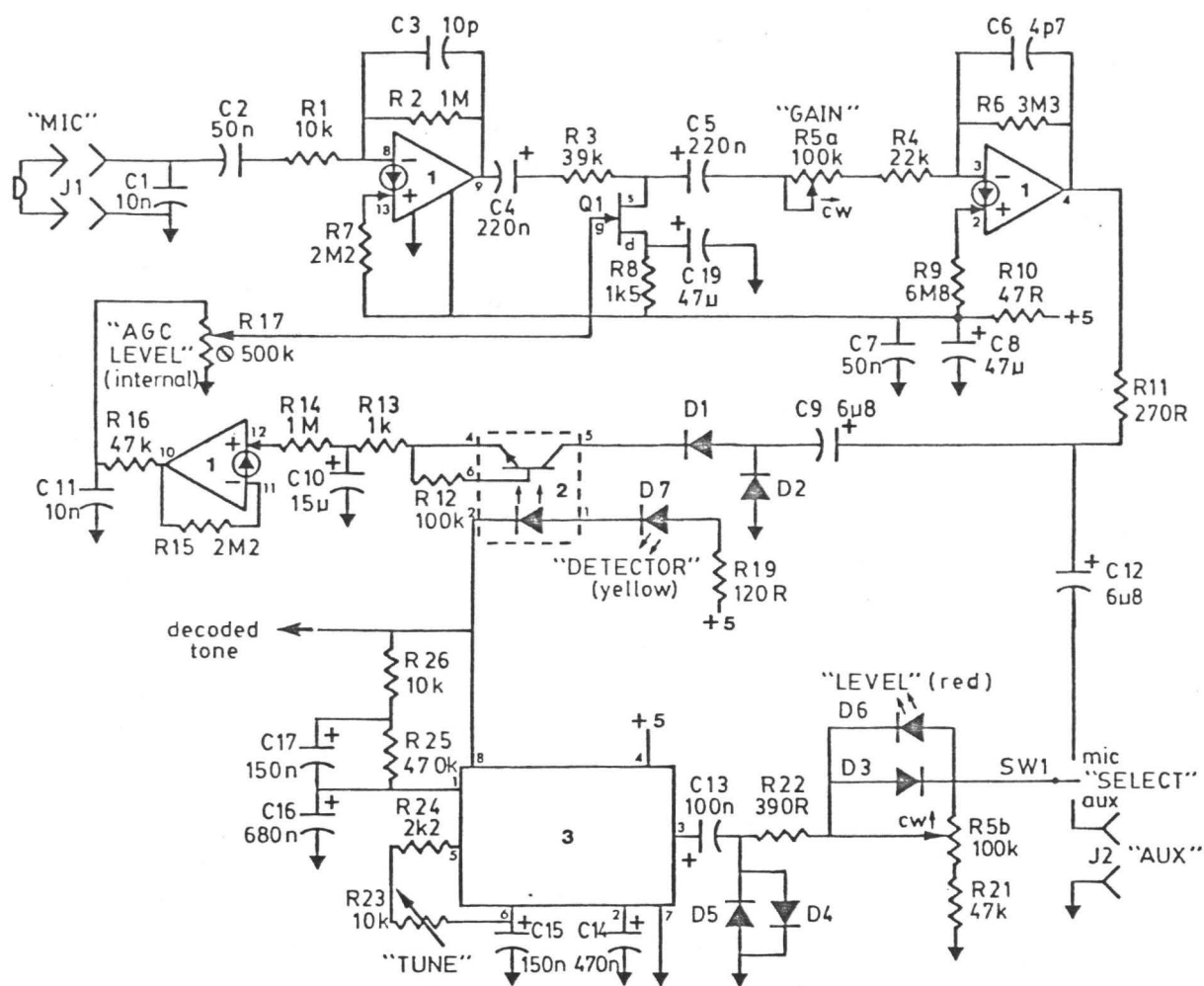


Figure 10. Final Footswitch (on wheelchair mounting).



Q1 2N3819 d
g
s bot

D1-5 1N914

D6,7 led

D8 1N4004

IC1 LM3900

IC2 4N26

IC3 LM567

Figure 11. Audio Input Schematic

over simple amplitude detection as it gives excellent rejection of background noise. As outlined in reference 13, there are design tradeoffs involving detection bandwidth, speed of response, and noise immunity.

Switch SW1 selects a high gain microphone input or a direct input for tape playback. Gain control R5b and D3, D6 form a nonlinear input network to aid noise rejection. At low gain levels (large inputs) the signal must overcome the forward voltage drop on the diodes, so low level noise is partially inhibited. D6 was also intended to function as a level indicator, but the signal required to activate it was too large.

R22 and D4-5 clip any input to 1 volt peak to peak. If larger signals are allowed IC3 tends to detect harmonics of low frequency signals. C15 and R23-24 determine the centre detection frequency, allowing tuning from 700 Hz to 3000 Hz. C14 affects detection bandwidth and loop response time, while C16 is part of a filter on the output switch. These two capacitors were varied considerably from design values to optimize switching times and dropout rejection. R25-26 add some hysteresis to the output, and C17 increases this feedback during switching transitions.

IC1-2 are an AGC (automatic gain control) amplifier for microphone inputs. IC1a provides a gain of 100, with 3db rolloffs at 300 Hz (R1, C2) and 16 KHz (R2, C3). Q1 and

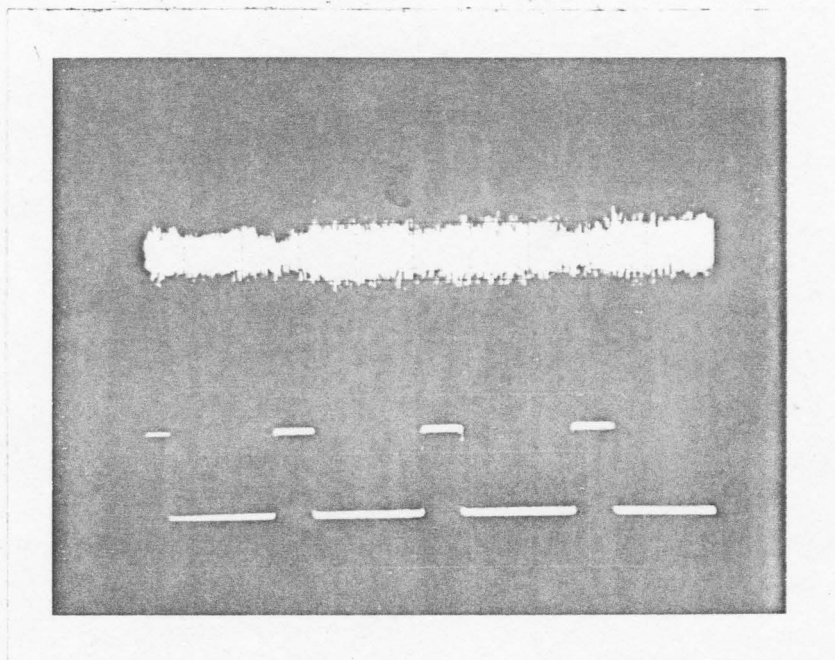
R3 are the variable gain attenuator necessary for AGC action. IC1b is the amplifier inside the AGC loop, and it also lowers the upper 3db point to 10 KHz. The external gain control varies the gain of this section from 27 to 150, exclusive of the attenuator. It may seem odd to have this control inside the loop, but it limits the maximum gain for low level signals when low gain is selected. Combined with the nonlinear network of R5b, D3, D6, at low gain levels low level signals are not well amplified by IC1b and are further rejected at the nonlinear network. By reducing these unwanted signals IC3 performs more reliably. High gain is available when required, but it should only be used in quiet surroundings.

D1-2 and C9-10 form a voltage doubler for AGC detection. Attack time is approximately 20ms (R11+R13, C10) and decay time 15s (R14, C10). IC1c adds a dc gain of 2 to the AGC control signal.

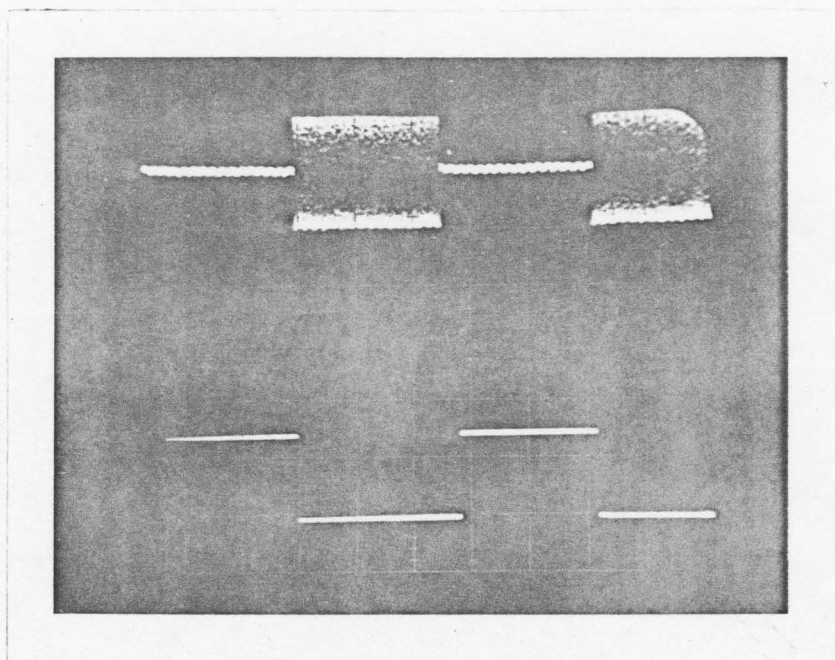
IC2 gates the AGC signal to C10 so that it is only increased to lower the gain when a tone has been detected by IC3. This stops high room noise from lowering the gain of the system and relates it more closely to the level of the desired signal.

Tests were done in a noisy room (fan noise) using the Sonalert Keyer and a microphone pickup several feet away. As Figure 12 shows, good digital signals were produced for

Figure 12. Phase Locked Loop Tests.



a) High ambient noise (short space, long mark)



b) Low ambient noise (long space, long mark)

both low and high noise inputs.

For tracking of noisier or faster code, one could add a bandpass filter in front of the microphone amplifier with its tuning coupled to the phase locked loop. This would greatly increase noise rejection and allow other time constants to be relaxed for better tracking of fast code.

4.6 Digital Input Board

This is a simple circuit (Figure 13) to condition the detected tone from the audio board and allow the direct use of single or double keys.

The three input circuits are identical and require a switch closure to ground for a mark. R25 is a pullup resistor and R26, C25 form a low pass filter with a cutoff near 70 Hz to reduce noise from IC3 or external keys. R27-28 and IC4 form a Schmitt trigger to square the signal for digital use. D10-11 give a visual indication of external switch activation.

IC5-6 combine the three possible inputs into two data lines. Pin 1 is high for a tone mark or the dot switch and Pin 2 is high for the dash switch. Both pins high indicates an error. SW2 informs the software whether a single or double switch is being used. Pins 1, 2 and 4 connect to data lines 0, 1 and 2 respectively on an input port on the microprocessor. The dual key and SW2 sense software was not

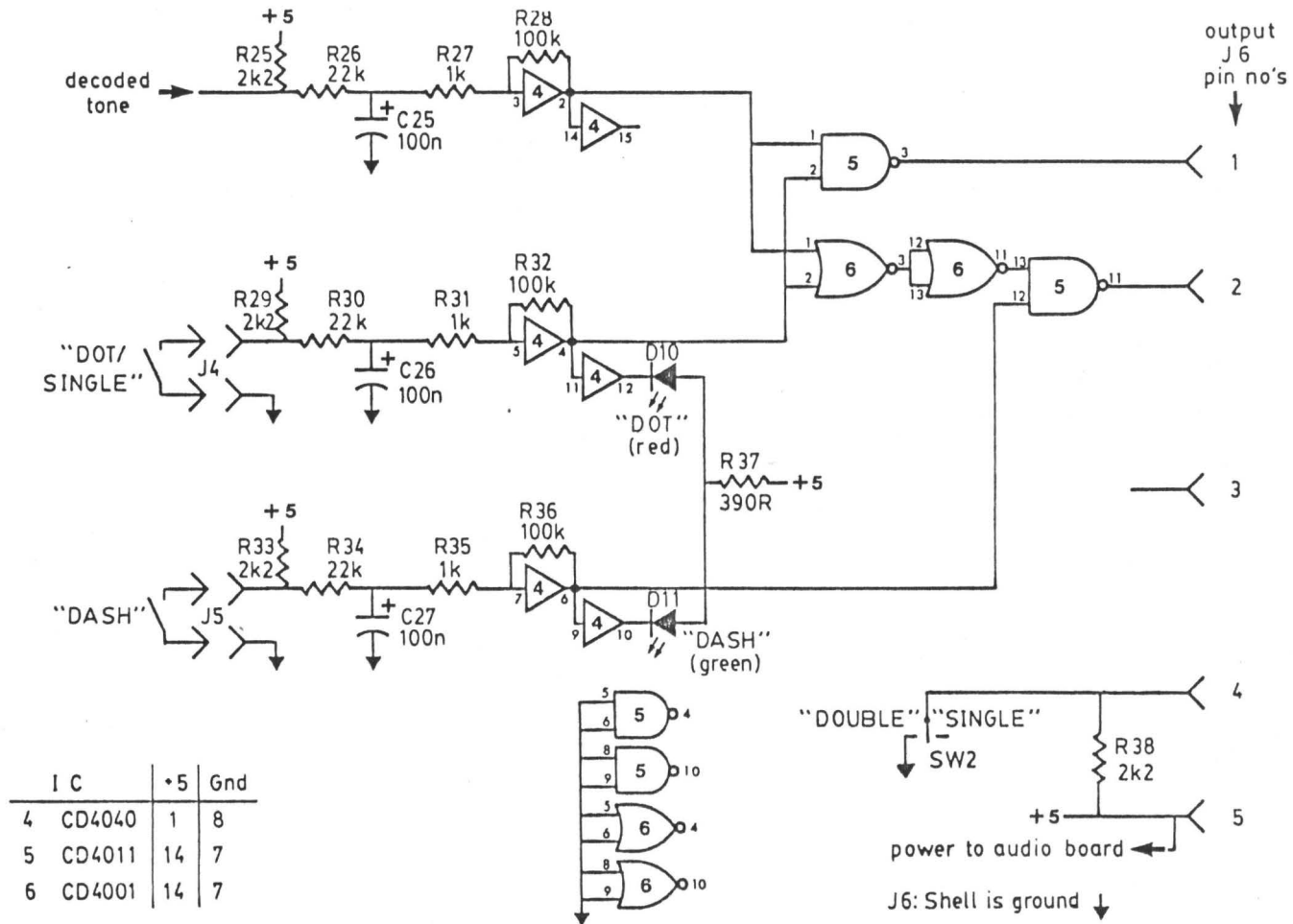


Figure 13. Digital Input Schematic

needed or written during development, but its addition would be trivial.

4.7 Microprocessor System

The microcomputer used for development is shown in Figure 14. It is based on the single board demonstration unit marketed by Tektron Inc. using the RCA 1802 CMOS microprocessor. It was chosen as its power requirements are a few milliamps at 5 volts, and it is thus suited for the ultimate goal of a portable device. It is an 8-bit machine with sixteen general purpose 8-bit registers and a suitable instruction set. Apart from RAM and ROM it is nearly a one chip processor, as the clock oscillator, DMA, interrupt and I/O ports are provided on the chip. Further details can be found in reference 16.

The system used for development was more complex than the capabilities of the chip might imply. The unit had to drive TTL circuits, which required buffering the data and address lines. The final system would be completely CMOS and the microprocessor could drive everything directly. Two serial I/O ports, including a cassette interface, were included, as well as a hexadecimal led display and an ASCII keyboard. On the demonstration board itself was an 8-bit latched output port driving eight leds and a hexadecimal keypad which was enlarged and brought out to the front

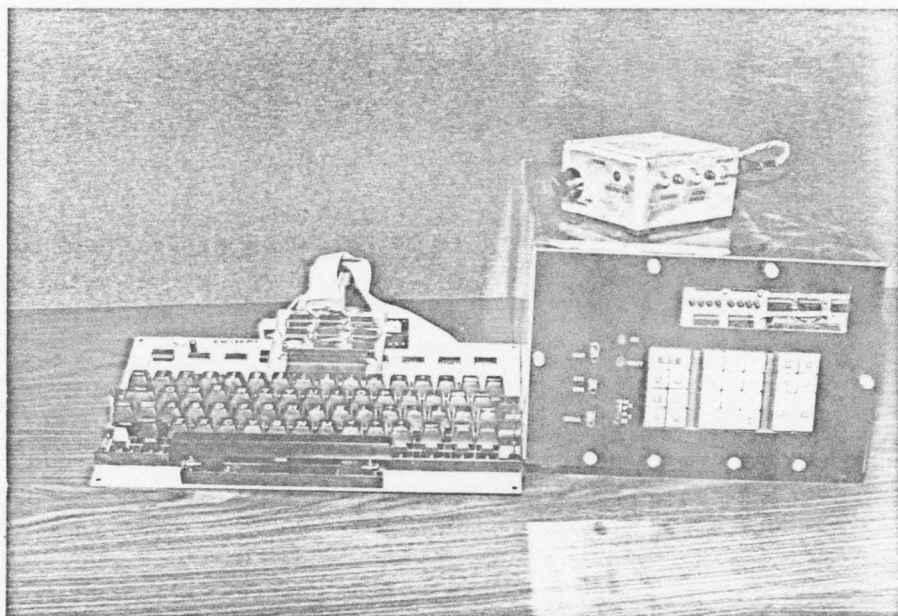


Figure 14. Development System.

panel. Memory consisted of 2K of RAM, but the first 512 bytes were replaced with EPROM containing the monitor.

The monitor program given in Appendix C allowed editing and execution of programs, and also contained utility routines for driving the various interfaces. Operation is based on the ASCII keyboard and the hexadecimal display, but it can be run from a terminal with less flexibility.

Circuitry for this system is not given as it does not relate directly to the final system. As pointed out above, that would consist of a one board computer without the extra buffering and peripheral capabilities.

CHAPTER V

RESULTS AND CONCLUSIONS

5.0 Testing and Results

Testing involved two main areas, testing with the patient and without. Both the hardware and the software were initially developed using the general rules for Morse code, and then modified as required.

The footswitch progressed from a single switch and a buzzer to a dual switch designed for the feet, plus an automatic keyer. The final combination of an adjustable wooden frame and microswitches covered in plexiglass proved to be functionally adequate and quite rugged. Required pressure for activation was easily varied by inserting foam rubber under the plexiglass. For a patient with weaker movements and more control it should be possible to place a single or dual switch entirely within a shoe using the Rancho type switch.

The automatic keyer was found to be necessary due to the subject's relatively poor control compared to a normal Morse operator. In this respect Morse code with its timing requirements might not be the ideal code for the cerebral palsied. Again, for a subject with better control the keyer

may not be needed, or could be incorporated into the software.

The program began as a short decoder of ideal code with fixed dot/dash ratios. This performed perfectly in tests using perfect code from another 1802 system, but failed miserably using hand sent code. The final program was much longer than anticipated (as usual!), but produced readable code. Some tests were conducted using code from a shortwave radio using the phase locked loop detector described in the hardware section. This code was being sent at rates up to 20 times those anticipated for this project, but readable code was still produced. One caution to others trying this test is that amateur radio operators often use a bewildering array of abbreviations which at first appear to be bad decoding. Due to international regulations, it is unfortunately illegal to reproduce here any of those tests.

Tests with the patient were quite brief due to logistic problems. These sessions were taped however and used repeatedly. It was found that for use as a conversational aid extra non-standard characters had to be added to do things such as blank the display and perform a carriage return. Due to the very long times between letters the automatic letter space should be deleted and replaced with a special space code, although this was not done.

5.1 Trial Unit

Near the end of the project a summer student built a portable decoder based on the work presented here. (15) The result, shown in Figure 15, consisted of the footswitch and keyer, the processor, the battery and charger, and a 20 character alphanumeric display. The size of the battery was due mainly to the fluorescent display.

Several trials were conducted with this apparatus on a wheelchair and it fulfilled the requirements. At the present time, intermittent electrical failures are preventing its use, but a more reliable unit will be produced.

5.2 Conclusions and Recommendations

The main objectives of the project were met, in that Morse code was shown to be a feasible means of portable communication for a cerebral palsied person.

Improvements can now be made all round, due to a better understanding of the problem, and more compact and lower power technology. Studies of the histograms generated by the user should be made so that the decoding algorithm can be tailored to the user. Convenience features can be included, such as the facility for the unit to operate as a computer terminal. A liquid crystal display would allow week long operation on a single charge, or a smaller battery. The actual switch required depends on the individual user, as with

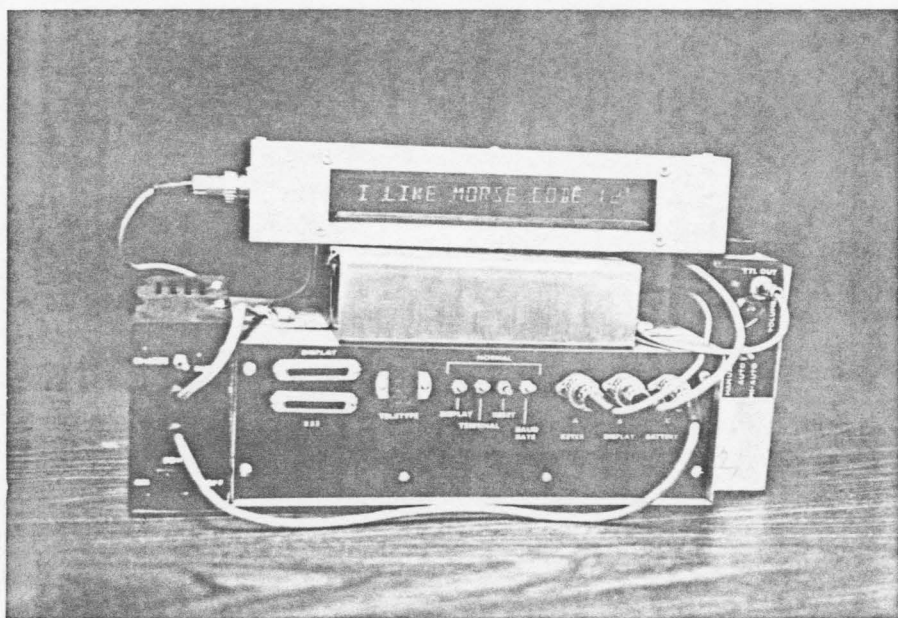


Figure 15. Trial Unit.

all aids, but a lighter and more visually pleasing switch could be designed.

It should be noted that most of the above improvements are aesthetic (editing, lower weight, etc.) rather than functionally essential. It is the author's experience however that unless an aid is almost invisible in its active and passive states it is apt to be rejected by the user.

In closing, an aid is only useful if it is used! It is the author's hope that further work will be done in this area, both in clinical trials and final placing of devices with the handicapped.

APPENDIX A

MORSE CODE CHARACTERS

AND

PROGRAM COMMANDS

MORSE CODE CHARACTERS

A . -	N - .	1 . - - - -	Period . - . - . -
B - . . .	O - - -	2 . . - - -	Comma - - . . - -
C - . - .	P . - - .	3 . . . - -	Question . . - - . .
D - . .	Q - - . -	4 -	Error
E .	R . - .	5	Colon - - - . . .
F . . - .	S . . .	6 -	Semicolon - . - . - .
G - - .	T -	7 - - . . .	Bracket - . - - . -
H	U . . -	8 - - - . .	Backslash - . . - .
I . .	V . . . -	9 - - - - .	
J . - - -	W . - -	0 - - - - -	
K - . -	X - . . -		
L . - . .	Y - . - -		
M - -	Z - - . .		

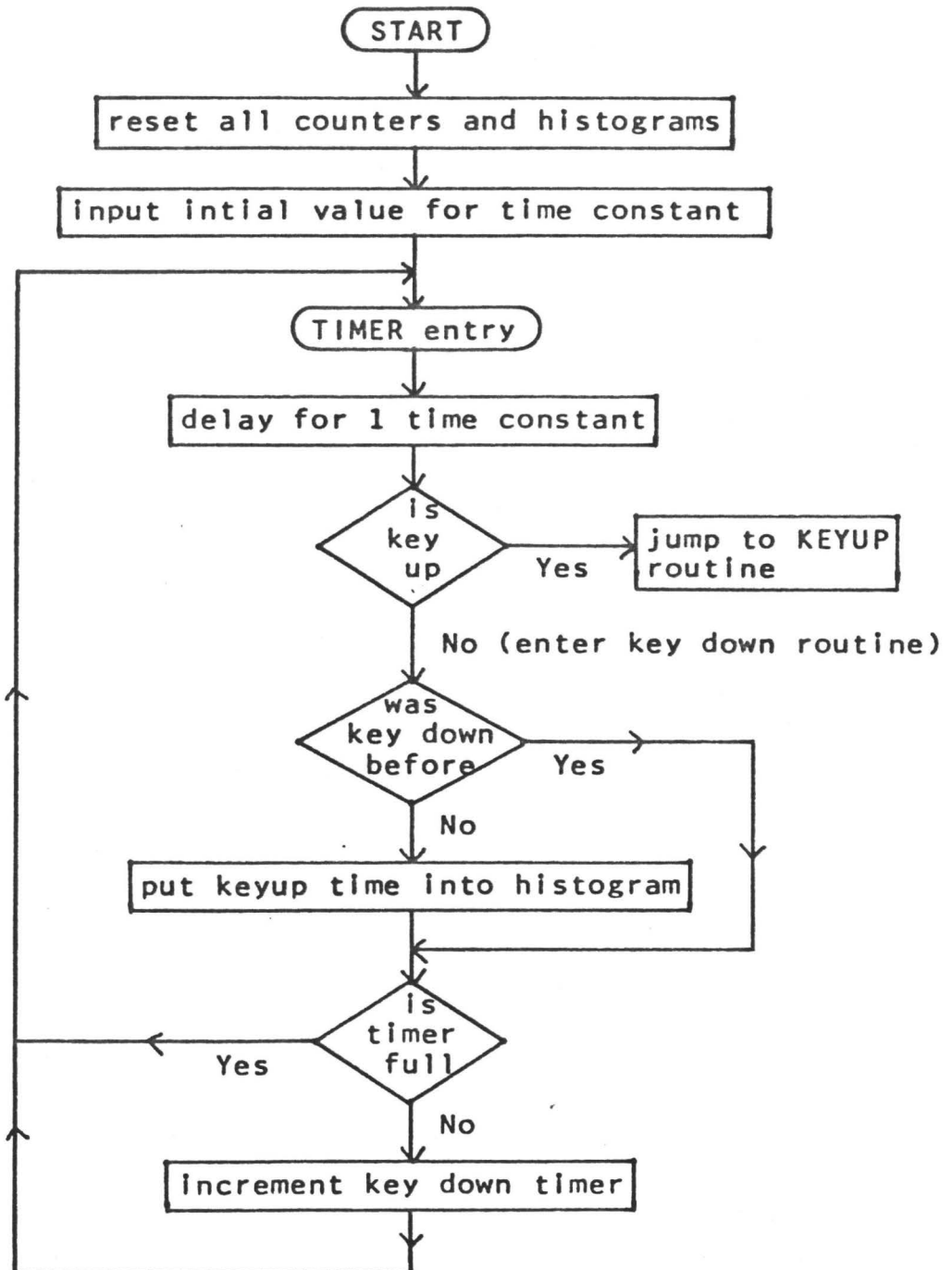
PROGRAM COMMANDS FROM ASCII KEYBOARD

/ Restart at 03E0, call monitor for new speed
H Output histogram immediately, wait for more commands
K Send "COPY" command to graphics terminal, wait
M Restart at 03F6, some histogram, speed and thresholds
N Set non-auto mode, restart at 03F6 (cancel a "Z")
X Clear histograms and continue
Y Divide histogram data by 2 and continue
Z Set automatic mode (output a histogram on any overflow)

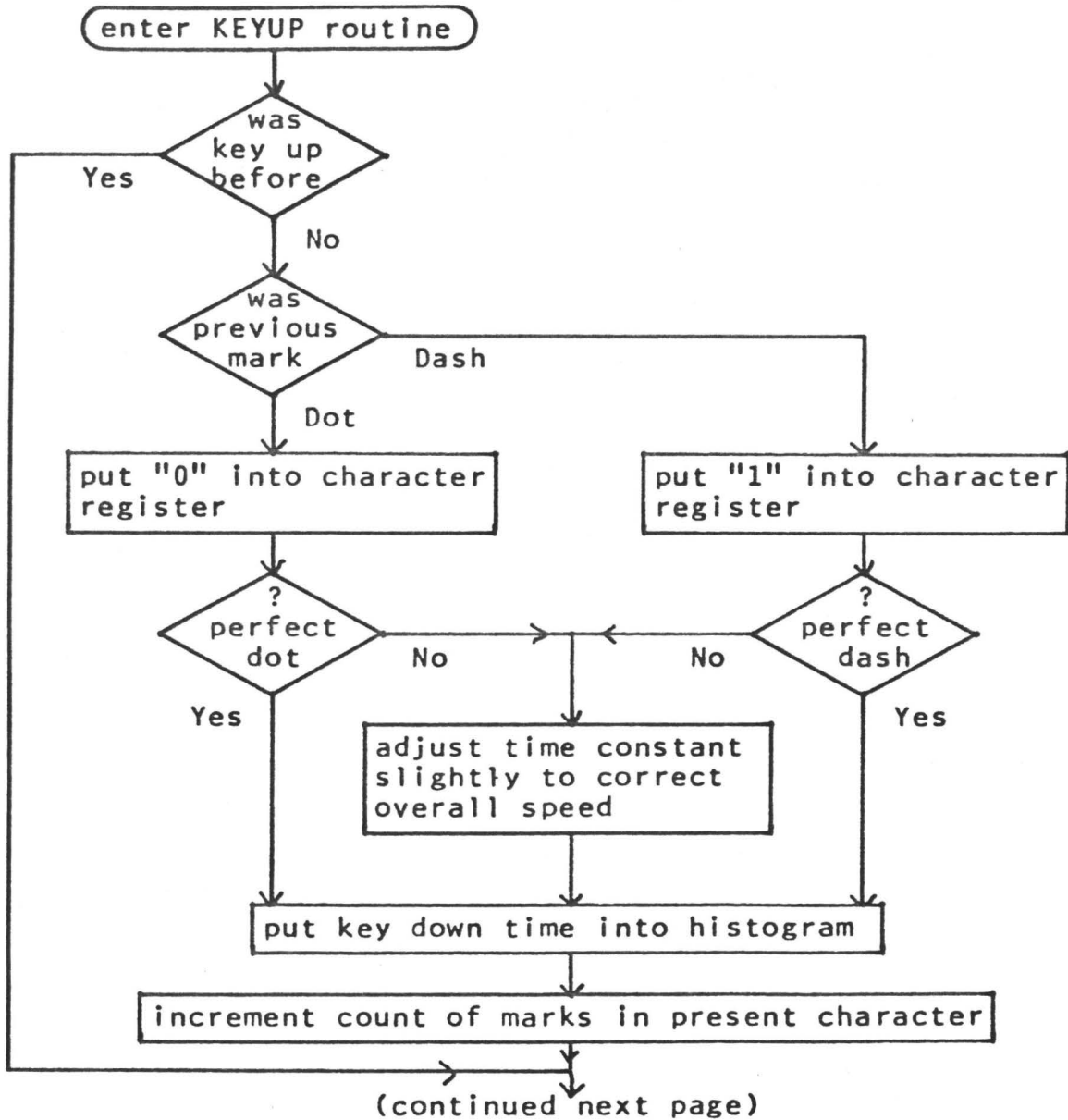
APPENDIX B

FLOWCHARTS

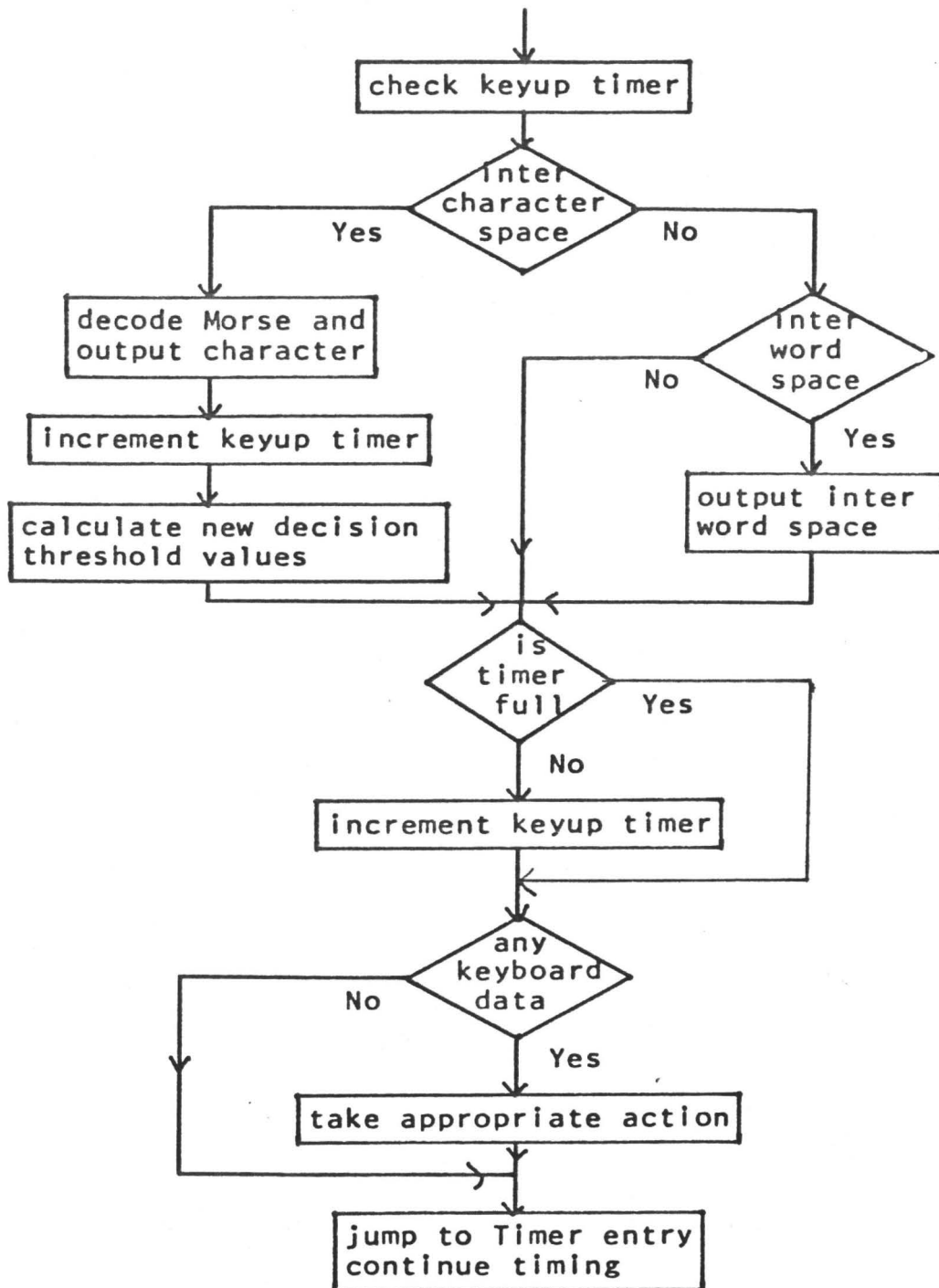
PROGRAM OVERVIEW



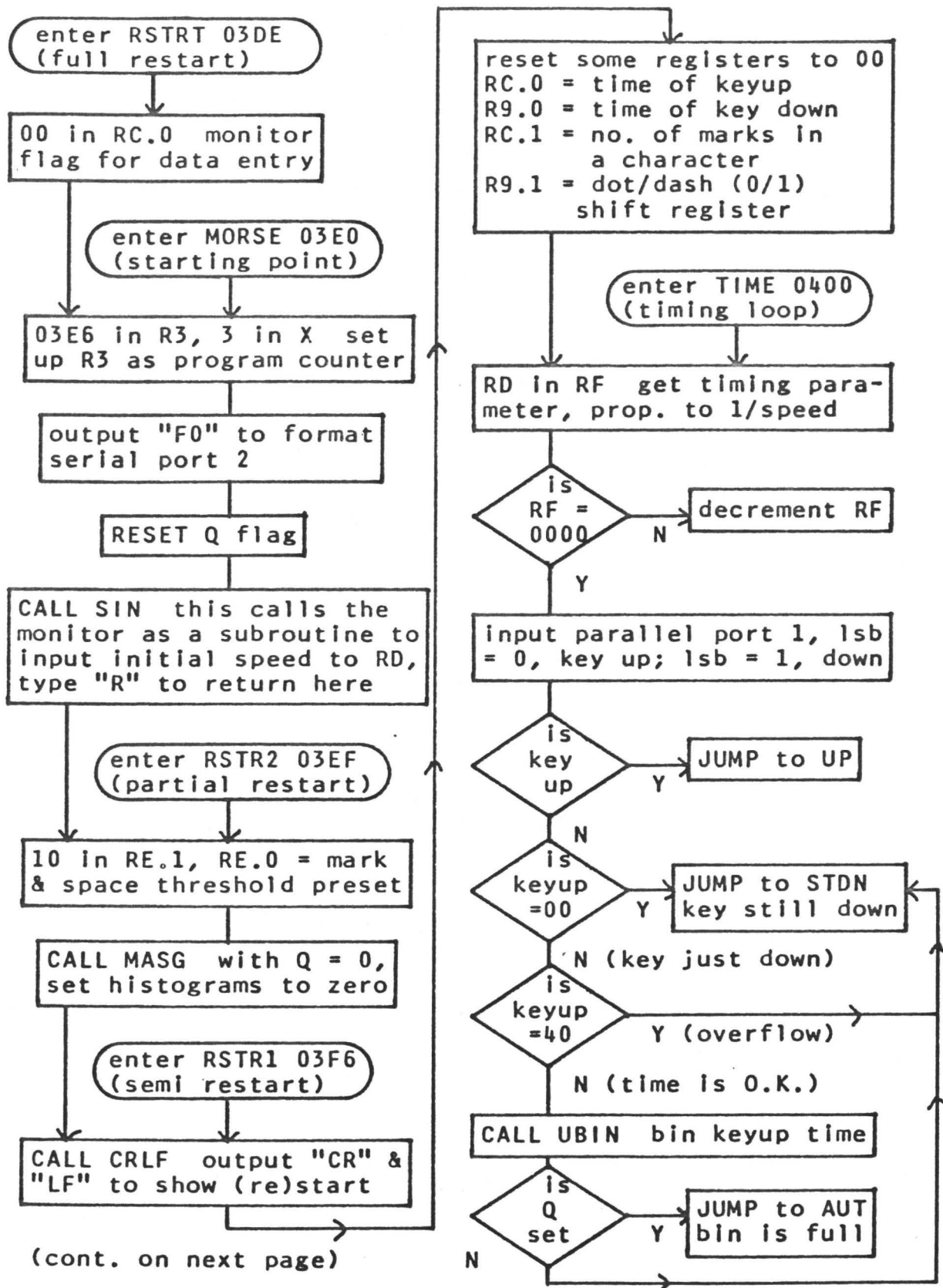
PROGRAM OVERVIEW (continued)



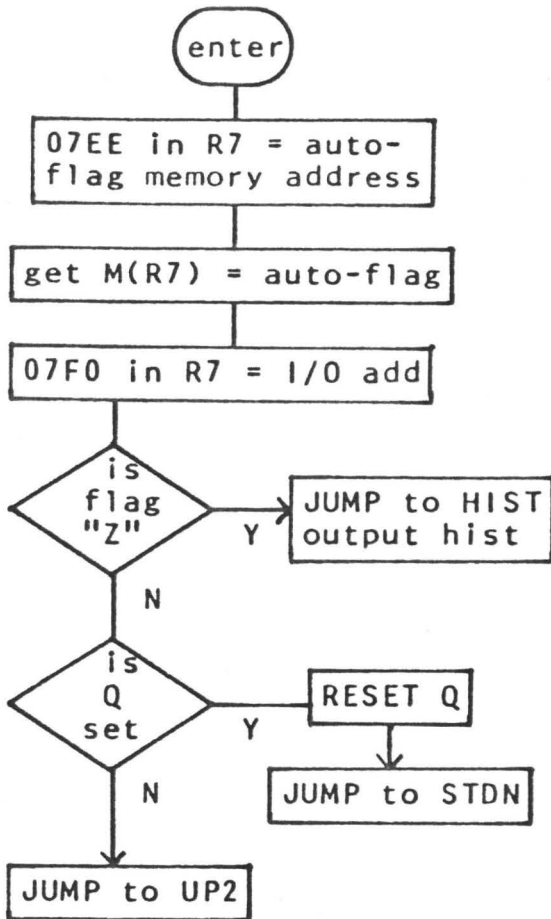
PROGRAM OVERVIEW (continued)



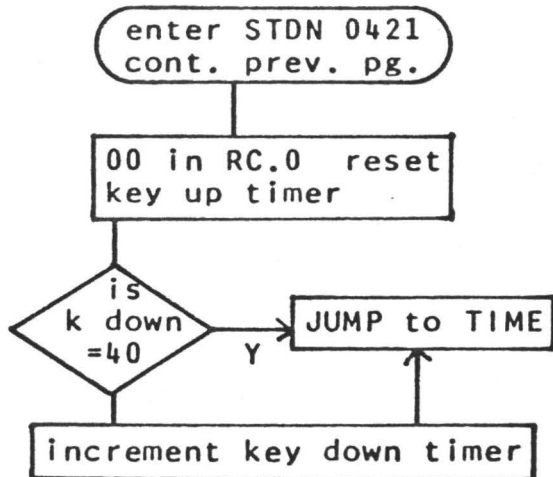
MORSE 03DE-03FF



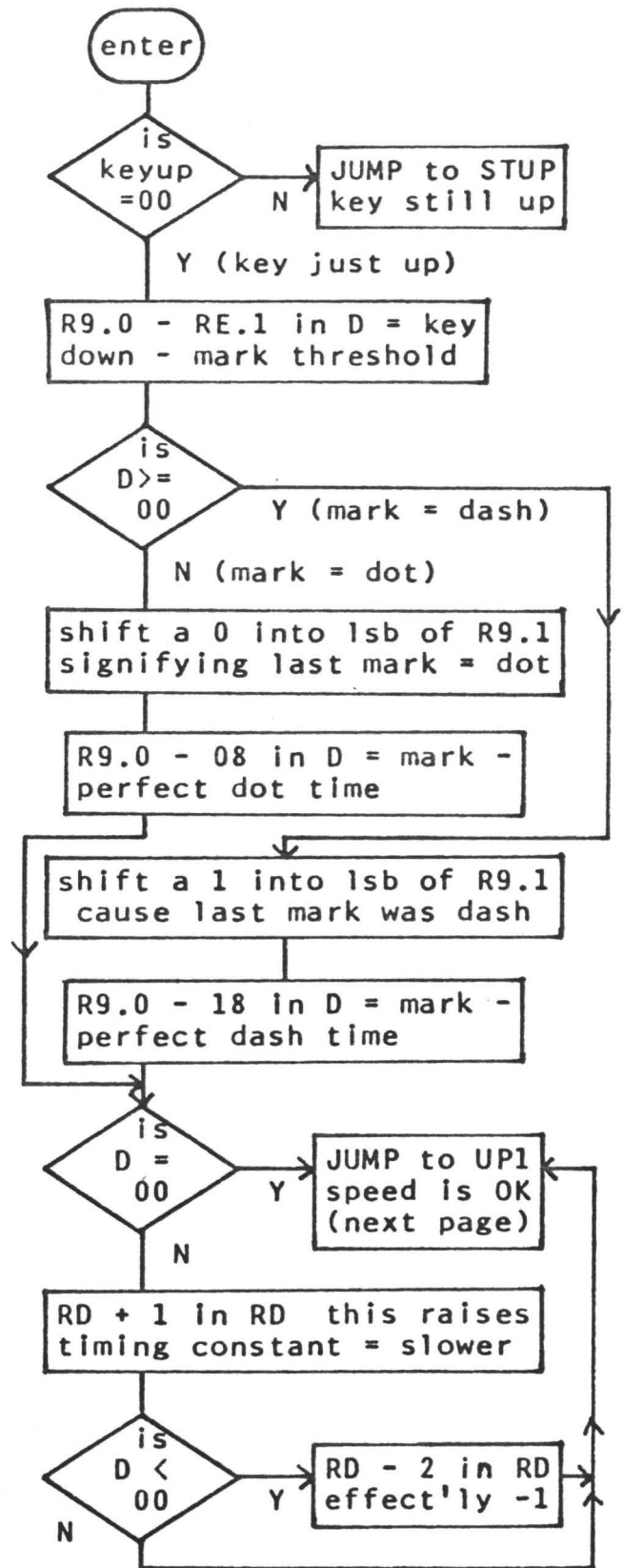
AUT? 0568-057F

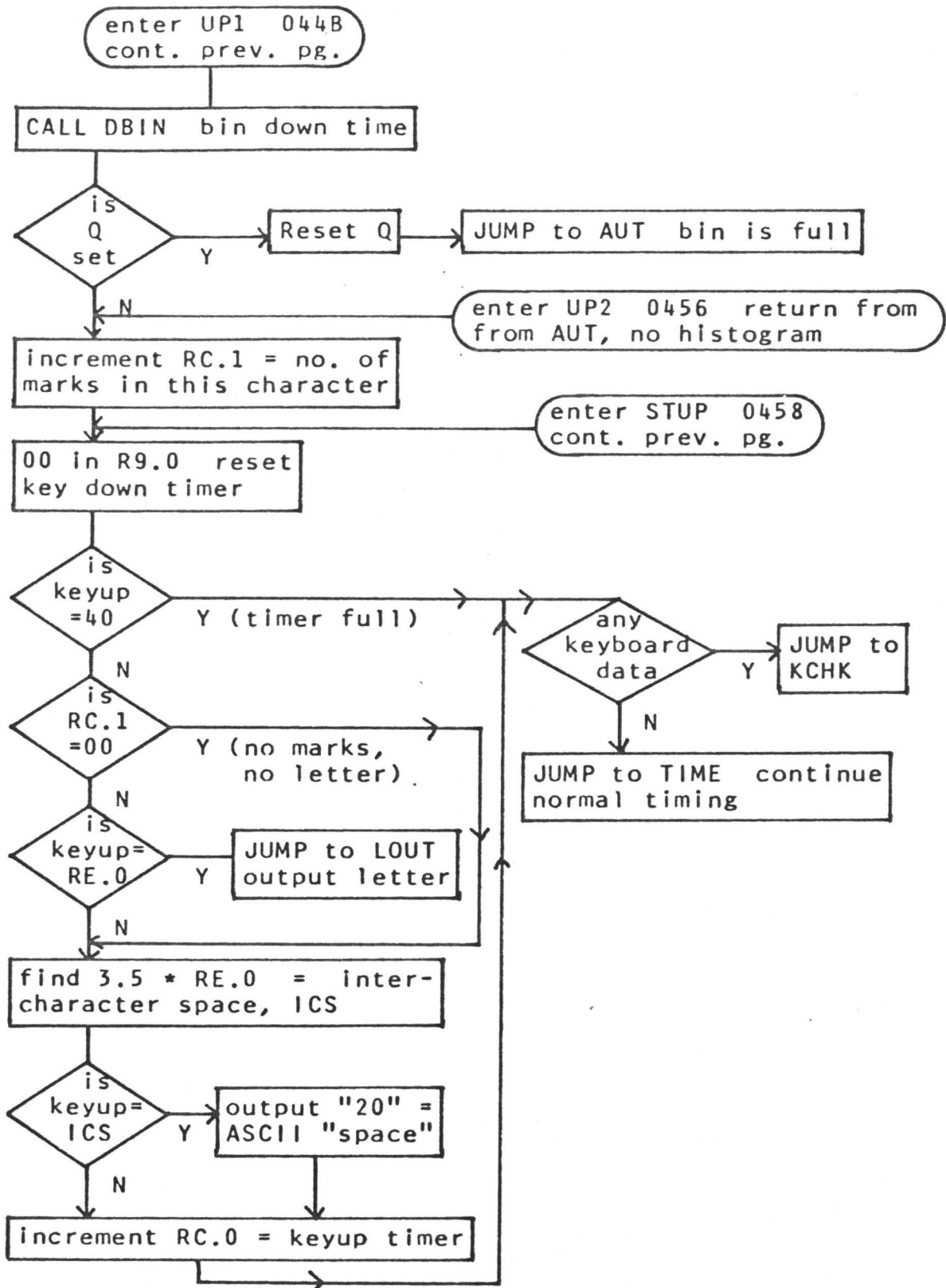


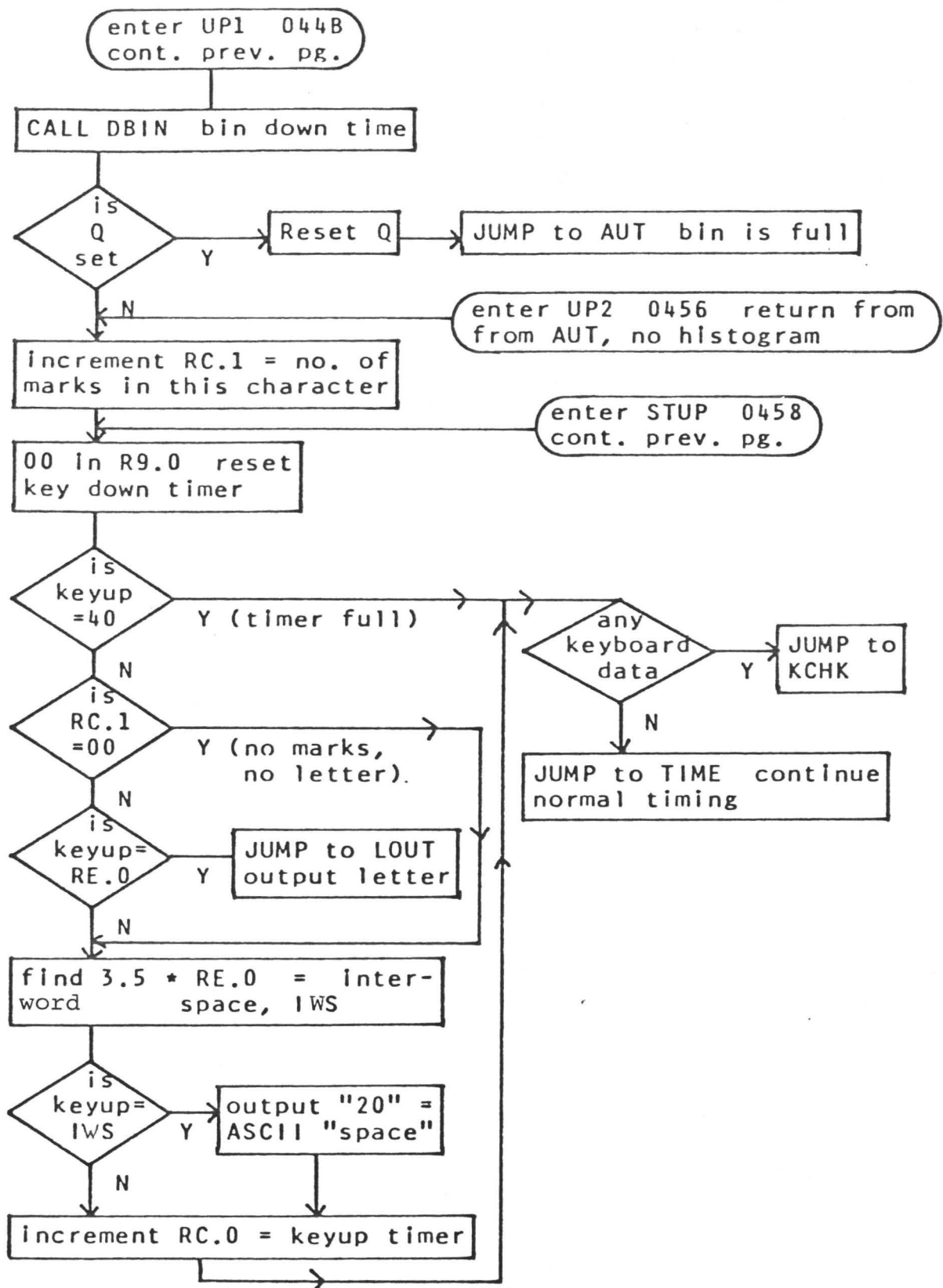
The "AUT" section of code is in a difficult position. It acts like a subroutine, but in fact "Q" is used to determine the "return" address.



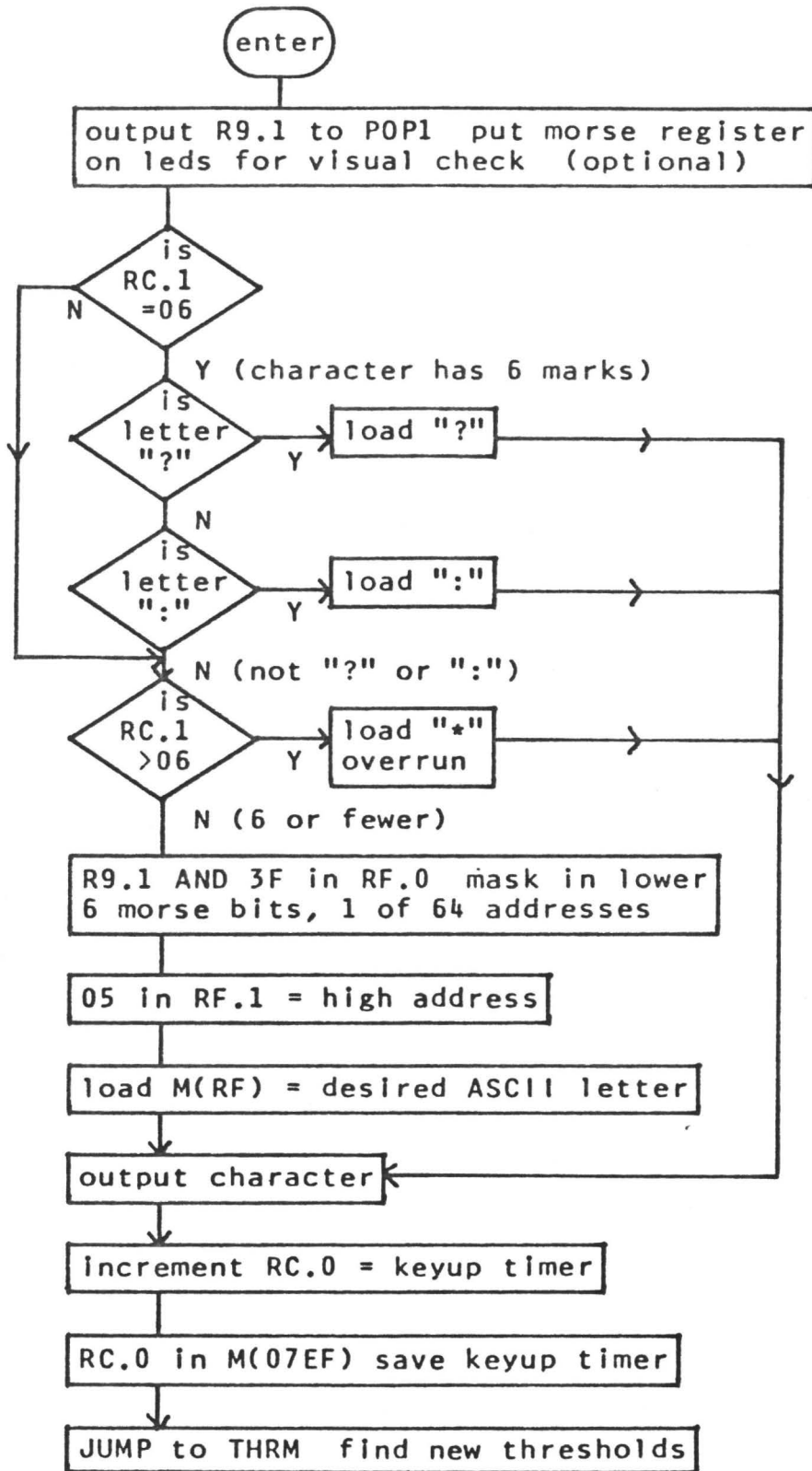
UP 042C-047F



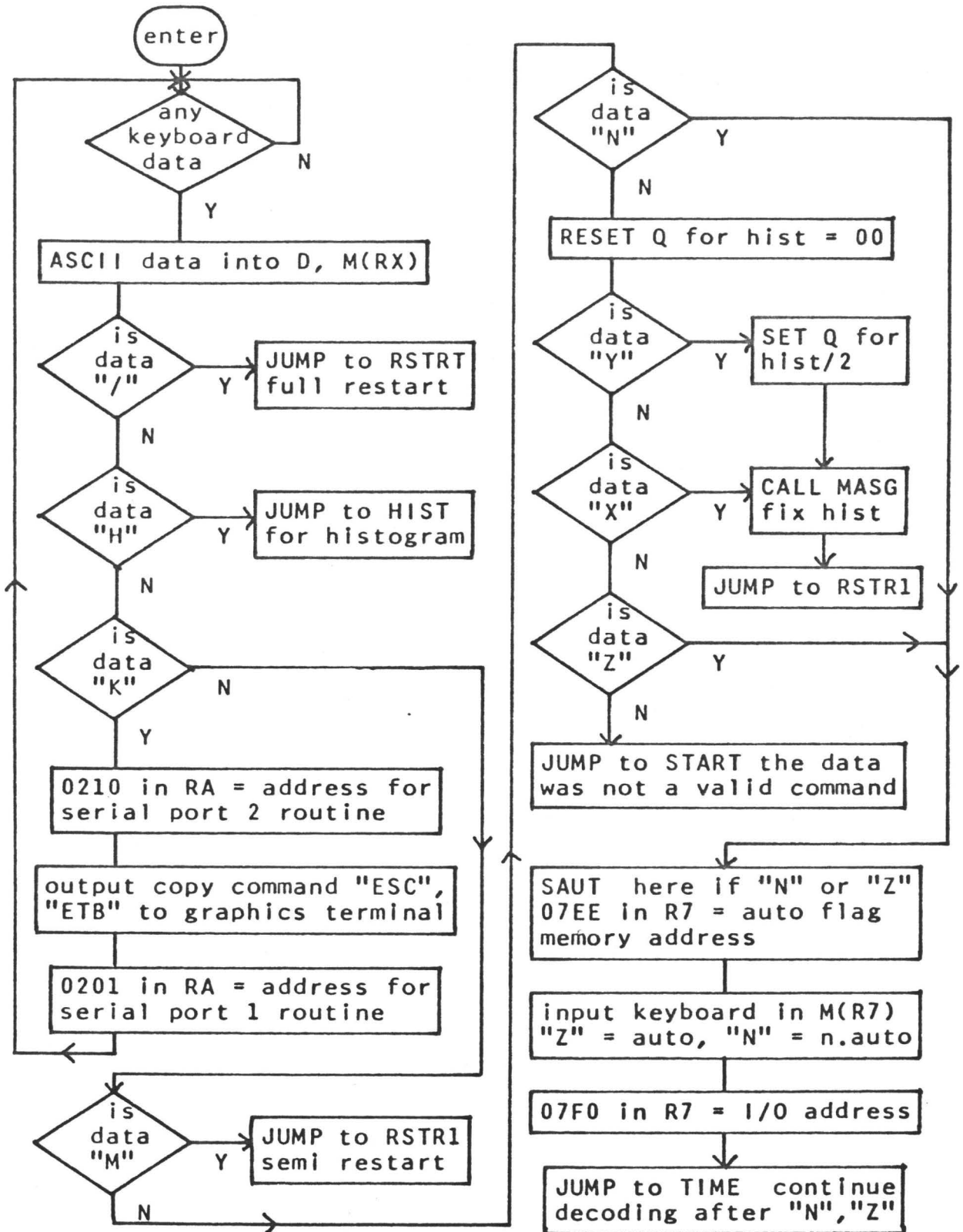




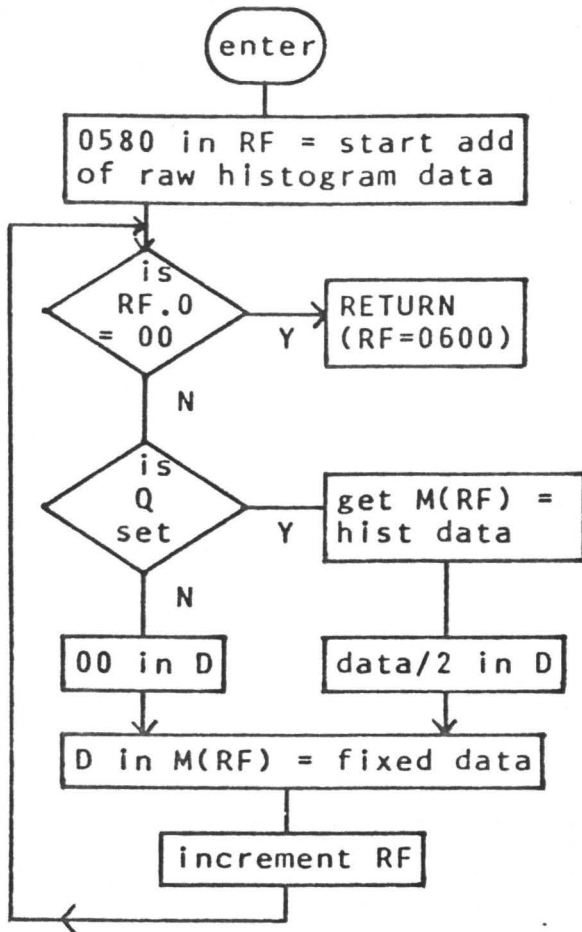
LOUT 0480-04B0



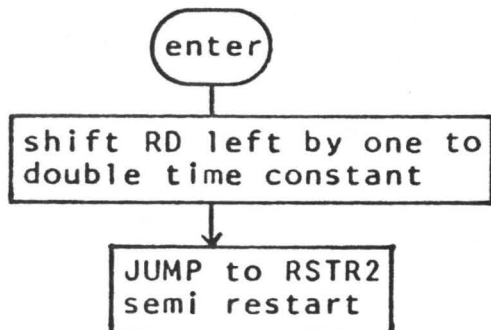
KCHK 04B8-04FC



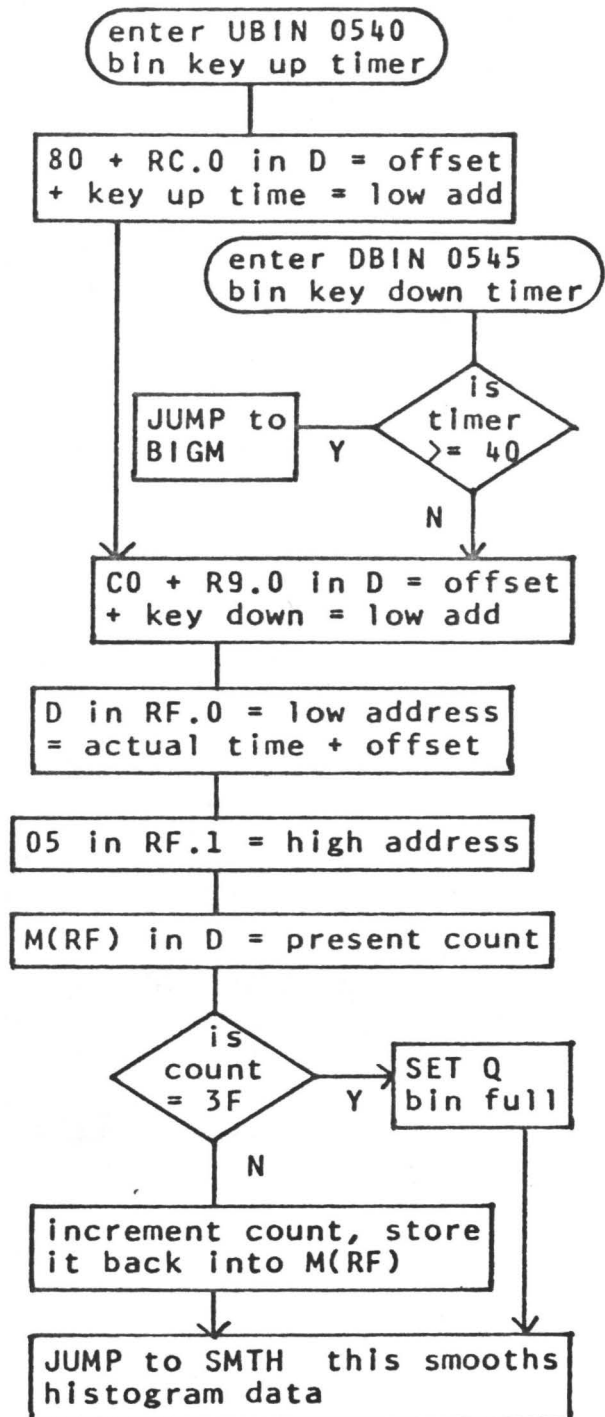
MASG 03C0-03D4



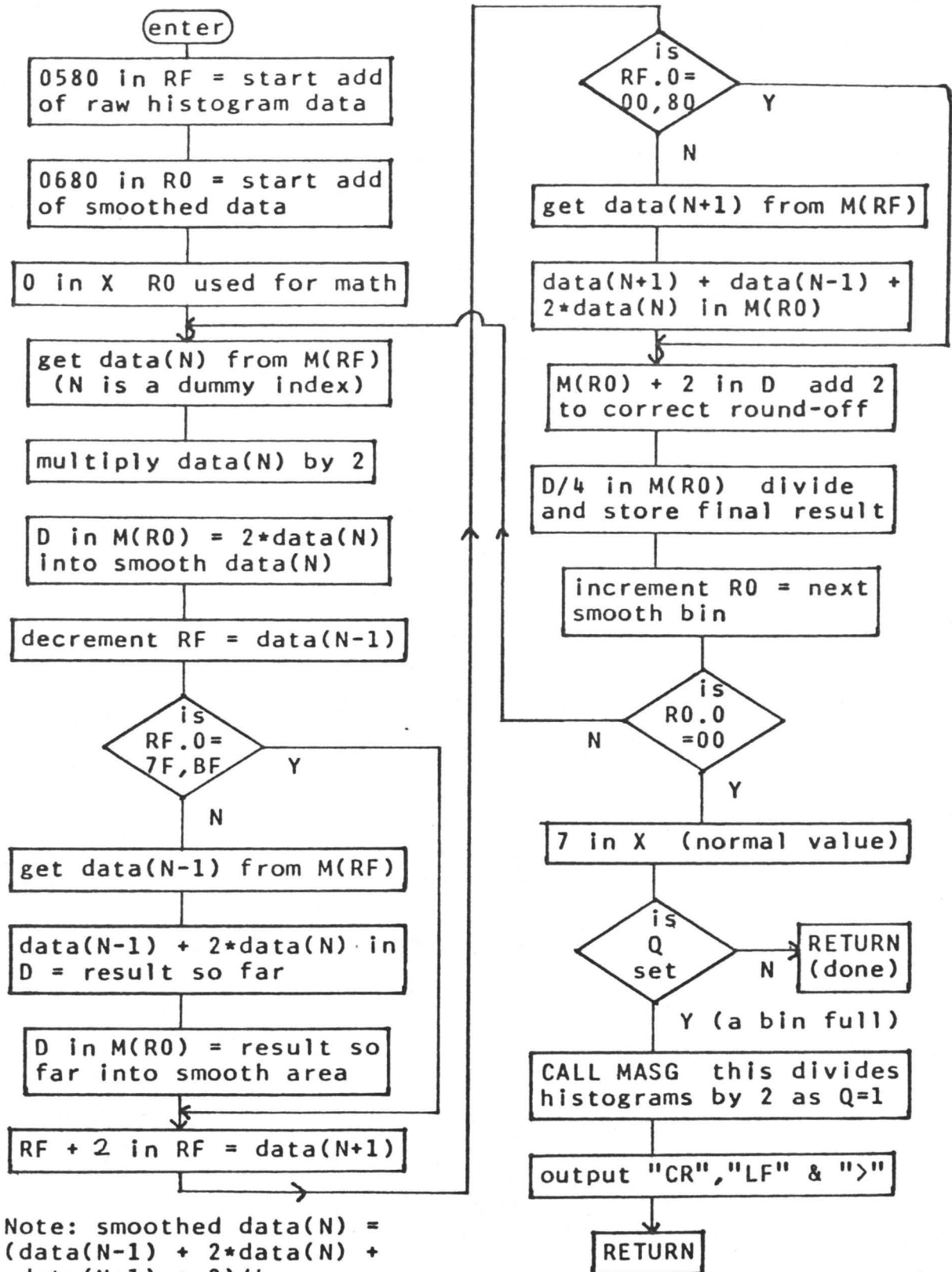
BIGM 055F-0565



UBIN, DBIN 0540-055E



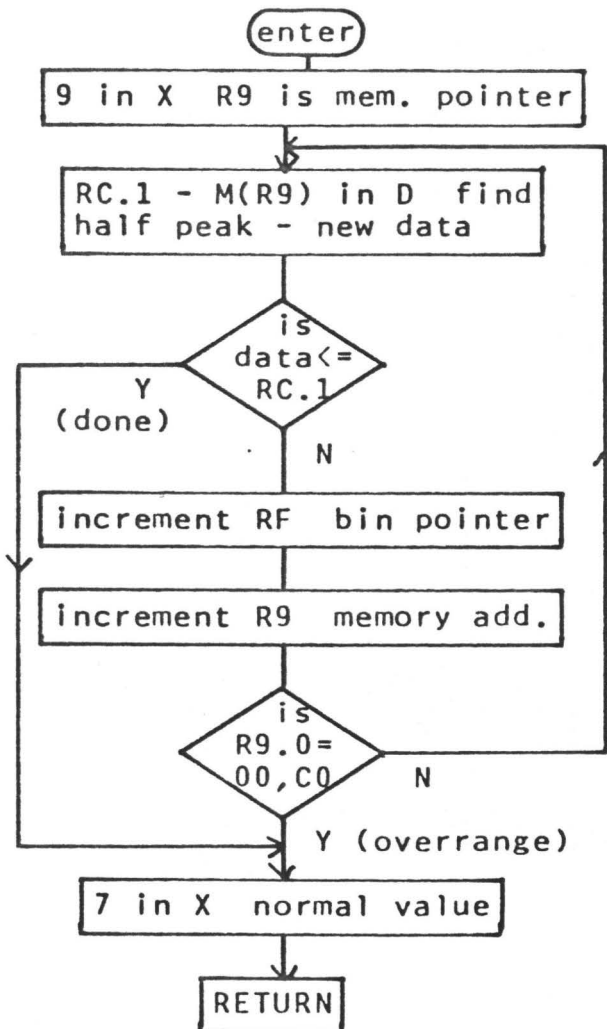
SMTH 0760-079F



LTOR 0710-0723

entry: half peak in RC.1
 start address of search
 in R9
 00 in RF.0

exit: RF.0 = # of bins
 from peak to first bin
 $\leq 1/2$ peak from left
 to right

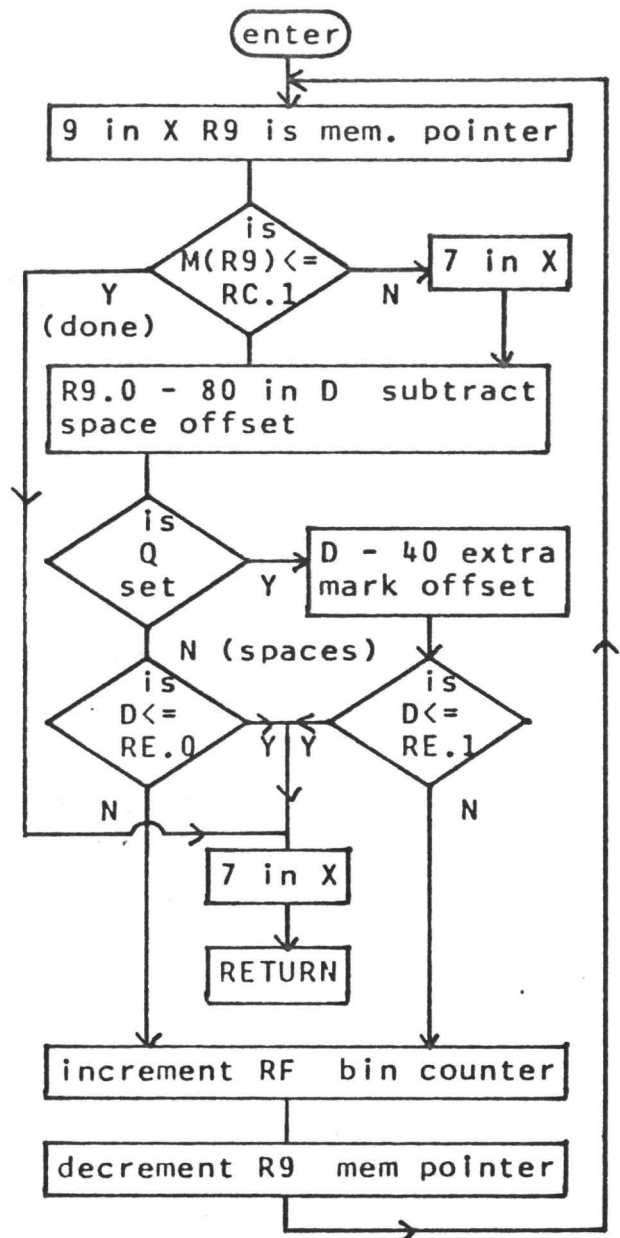


Note: LTOR and RTOL are very similar and differ in direction of scan and overrange criteria.

RTOL 0728-073F

entry: half peak in RC.1
 start address in R9.0
 00 in RF.0
 lowermost bin # in RE.0 for
 spaces, RE.1 for marks
 $Q = 0$, spaces; $Q = 1$, marks

exit: RF.0 = # of bins from
 peak to first bin \leq peak
 from right to left



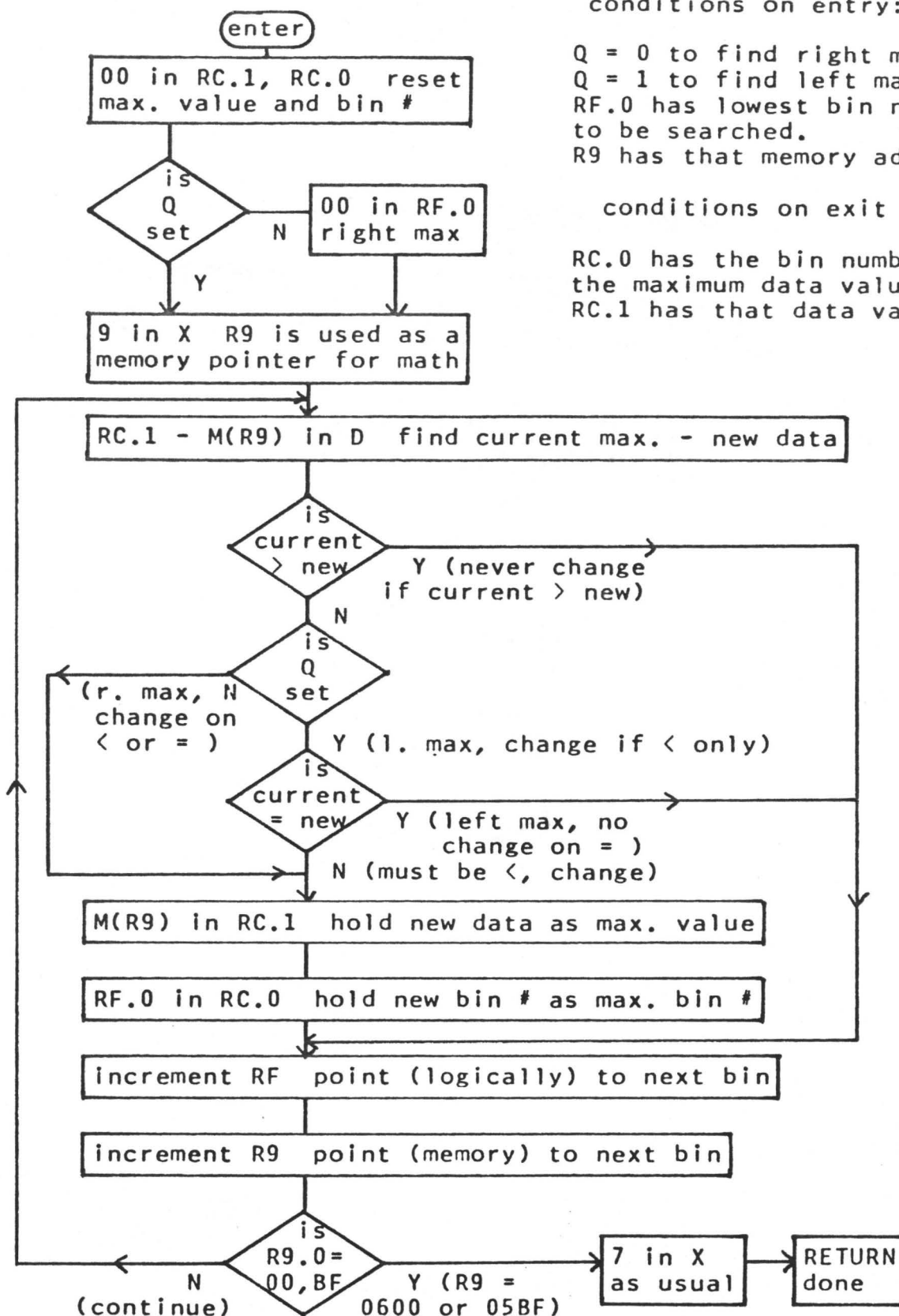
FMAX 0740-075D

conditions on entry:

Q = 0 to find right maximum
 Q = 1 to find left maximum
 RF.0 has lowest bin number
 to be searched.
 R9 has that memory address.

conditions on exit

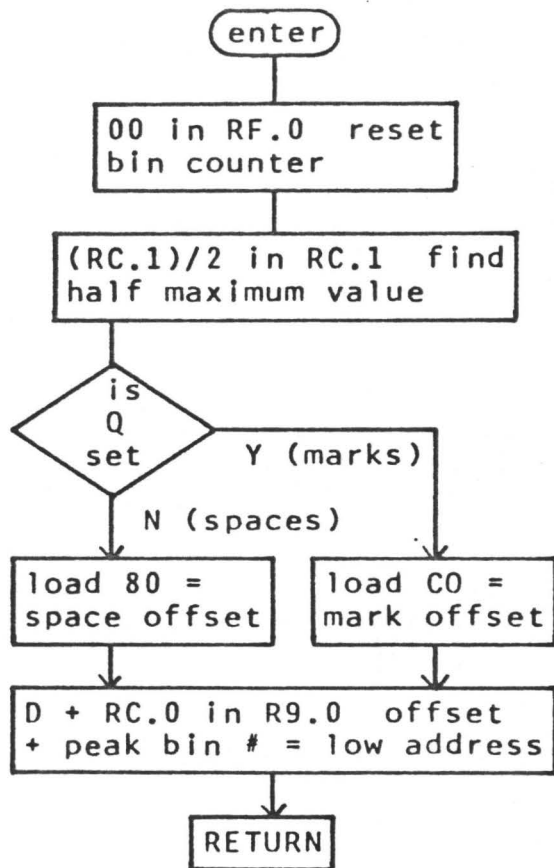
RC.0 has the bin number with
 the maximum data value.
 RC.1 has that data value.



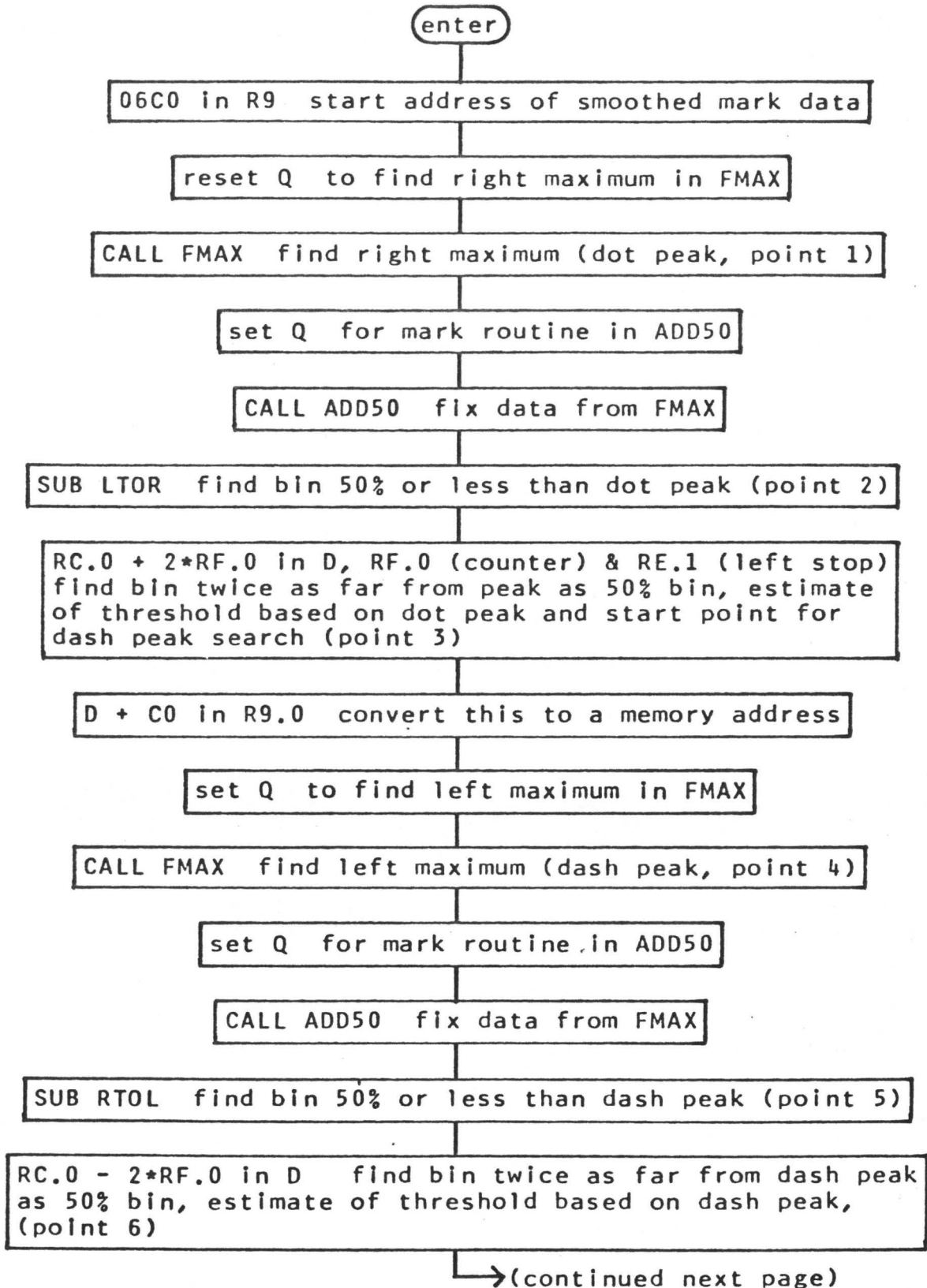
ADD50 0700-070F

entry: peak value in RC.1
peak bin # in RC.0
Q = 1 for marks
Q = 0 for spaces

exit: half peak in RC.1
actual memory add.
of peak in R9.0



THRM 0600-063B



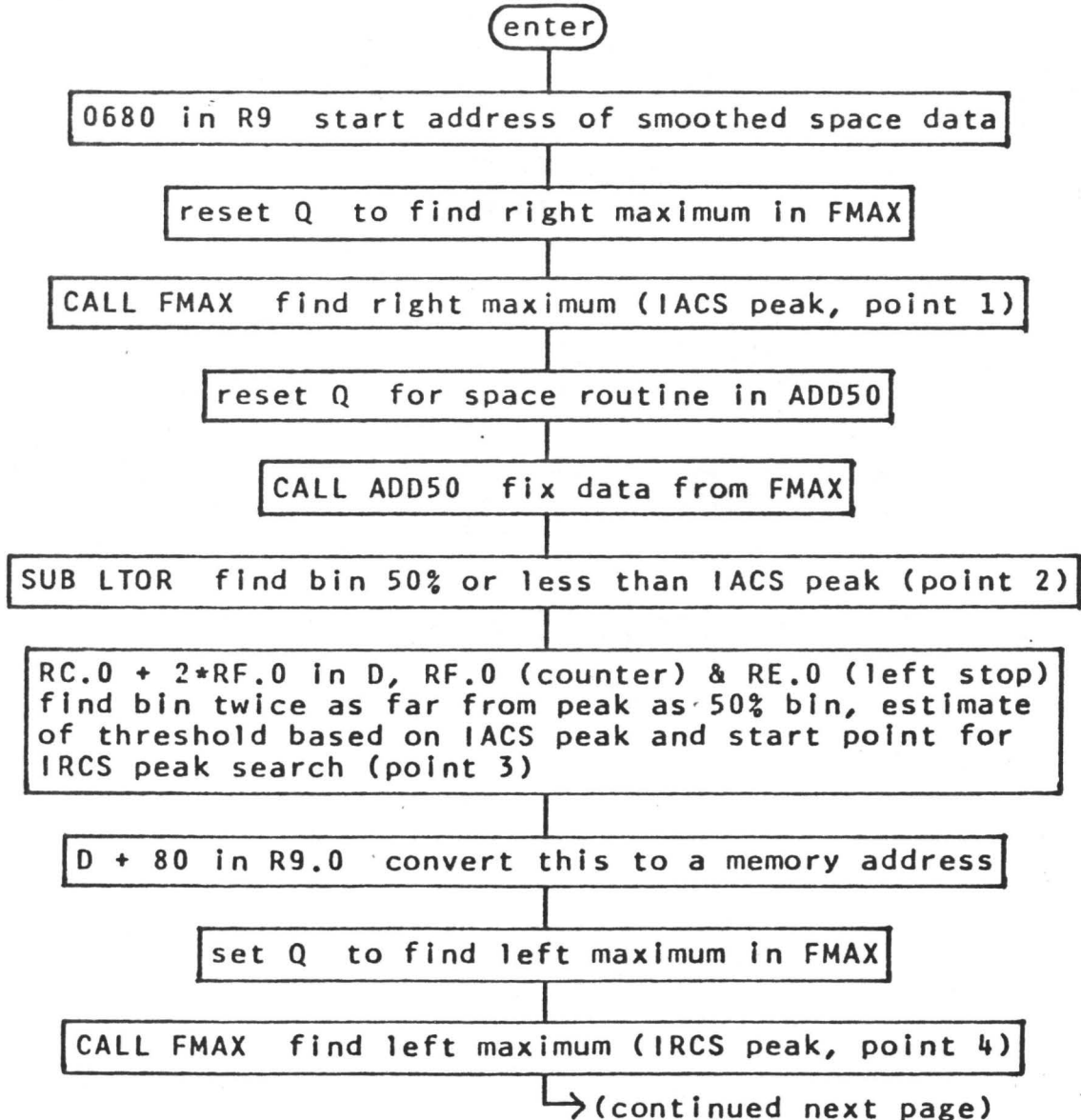
THRM (continued previous page)

(RE.1 + D)/2 in RE.1 average two estimates (dot, dash) for mark threshold, store it in RE.1 (point 7)

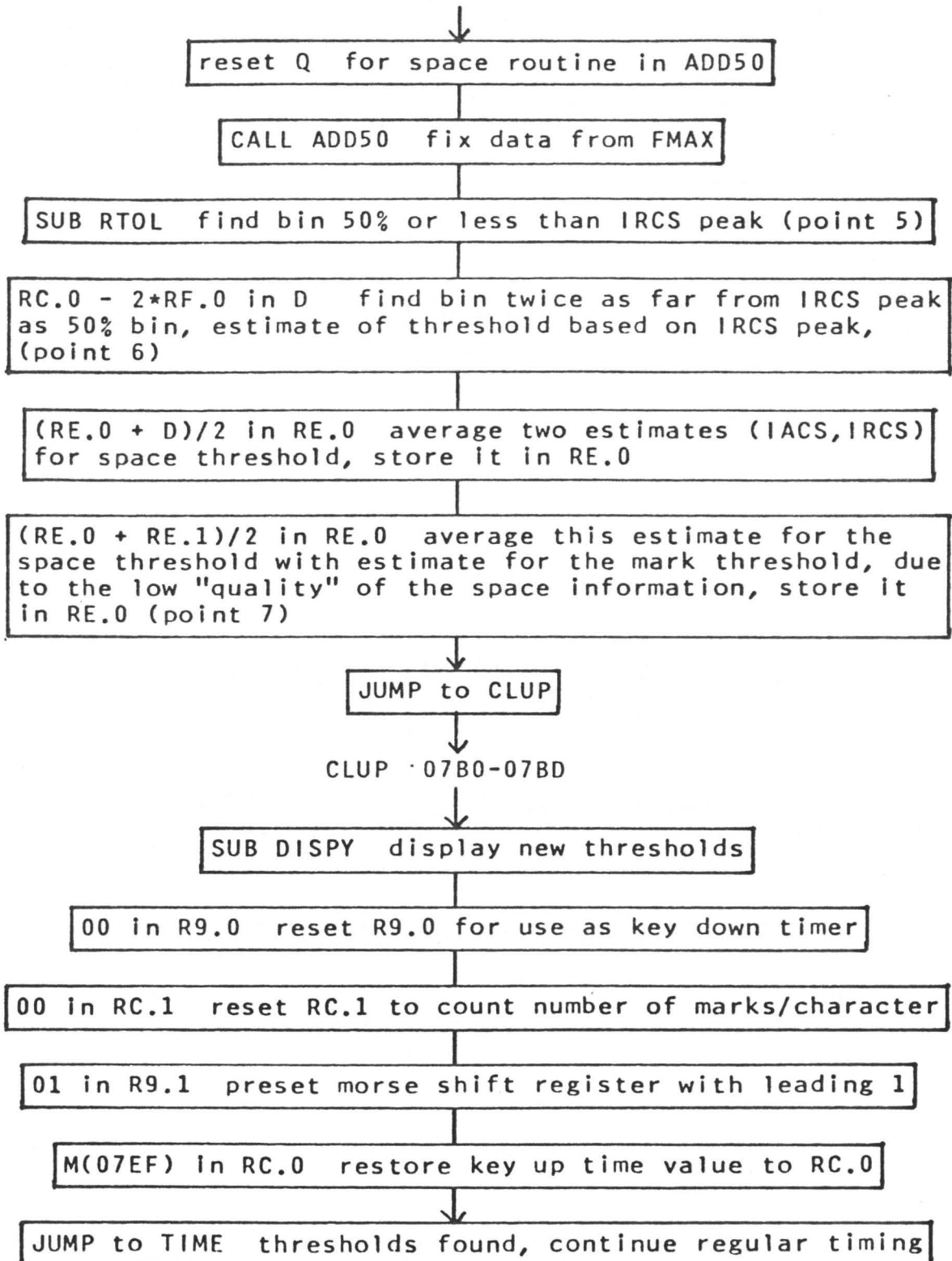
JUMP to THRS

THRS 0640-067F

This routine is very similar to THRM, and differs in detail only. IRCS refers to the intercharacter space, and IACS to the intracharacter space.

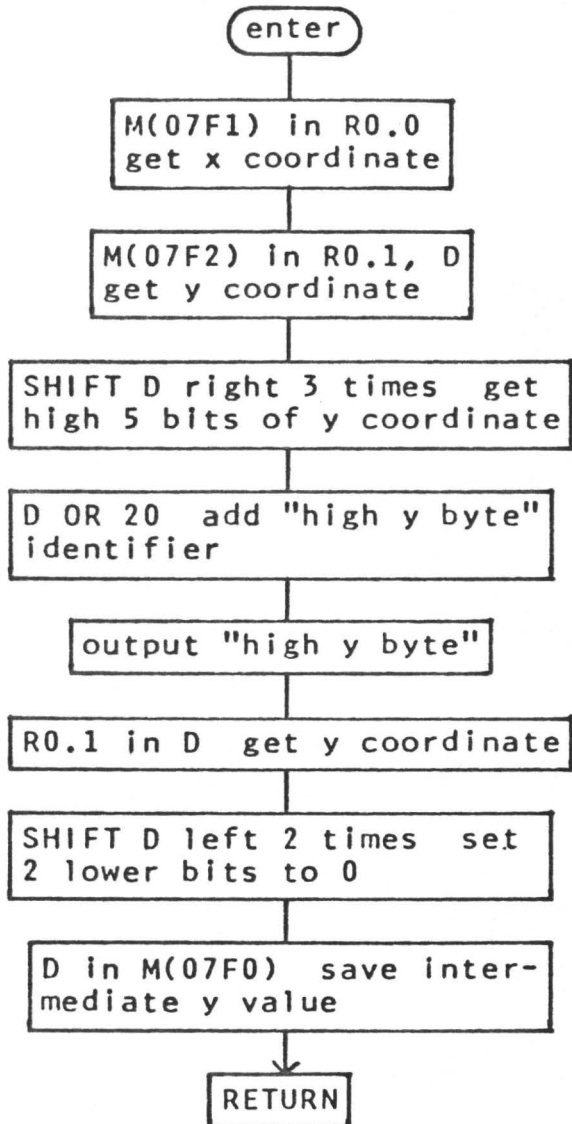


THRS (continued previous page)



XYOU1 02E0-02F1

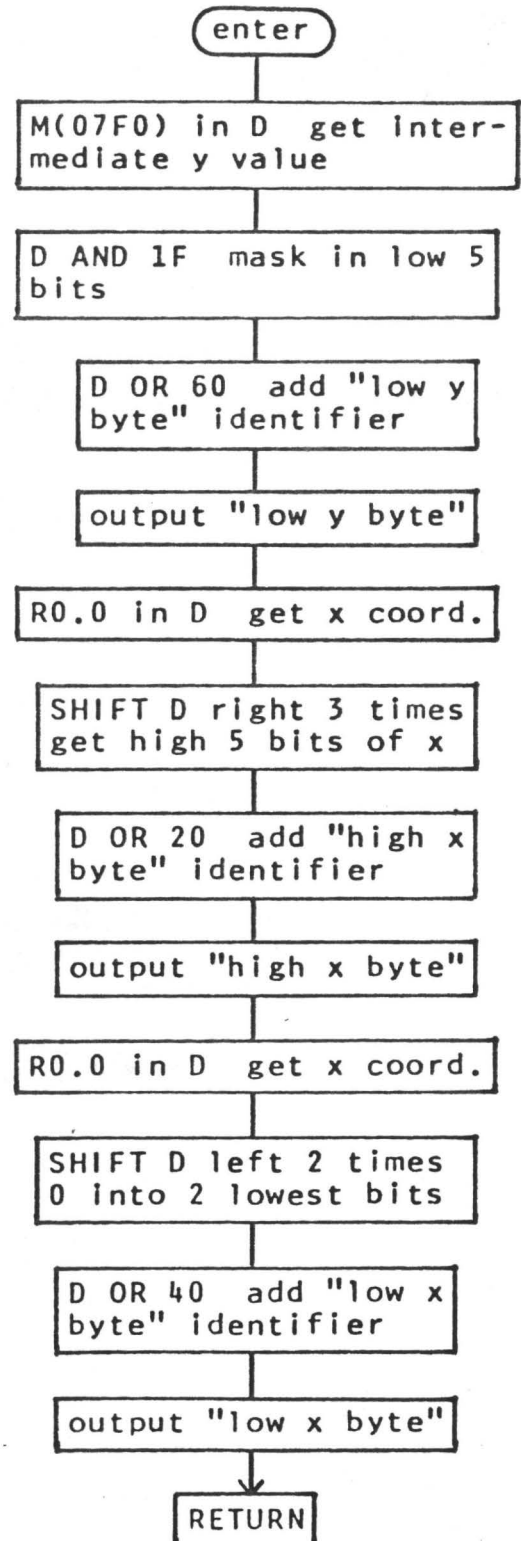
entry: 8 bit x coordinate
in M(07F1), y in M(07F2)



At this point the intermediate y value in M(07F0) can have the two lowest bits set by the main program. This would allow the full 10 bit accuracy of the graphics terminal to be used.

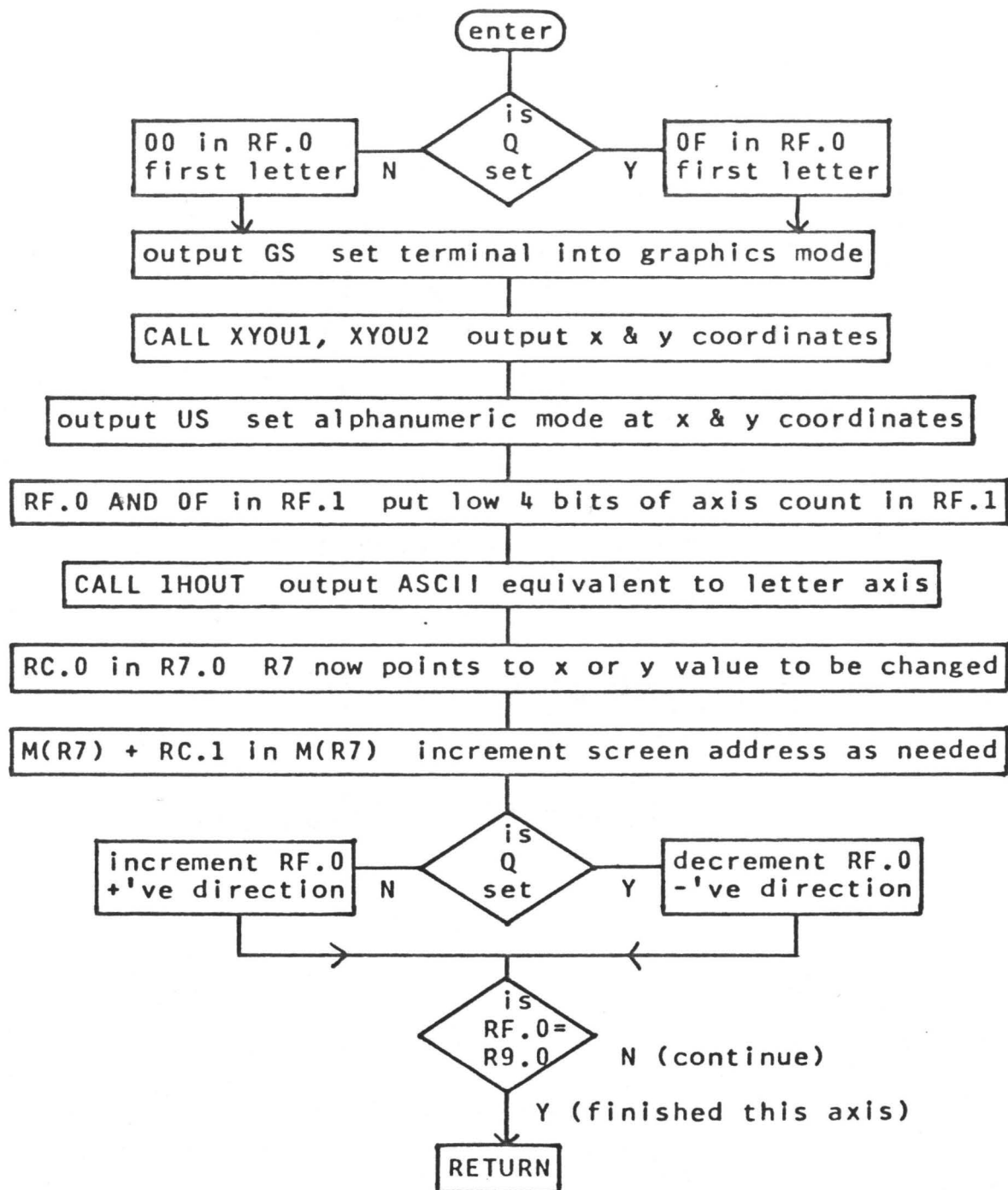
XYOU2 02F8-030D

entry: intermediate y
value in M(07F0)

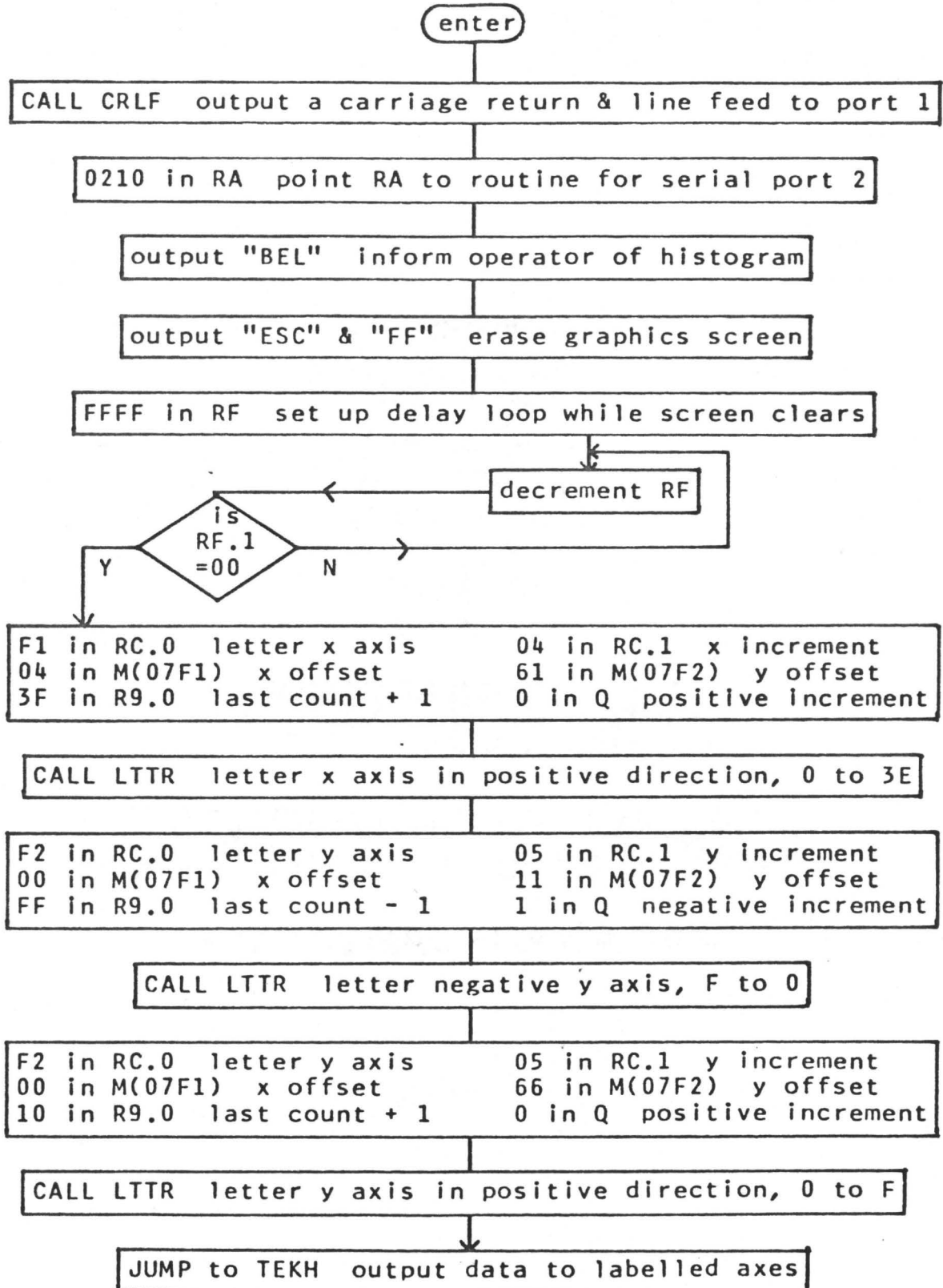


LTTR 0220-024B

entry: RC.0 = F1 for x axis, F2 for y axis; a low mem. add.
 Q = 0 to letter in +'ve direction (x or y) and v.v.
 RC.1 = screen address increment per letter
 M(07F1) = x coordinate, M(07F2) = y coordinate
 R9.0 = last position to be lettered, +/- 1

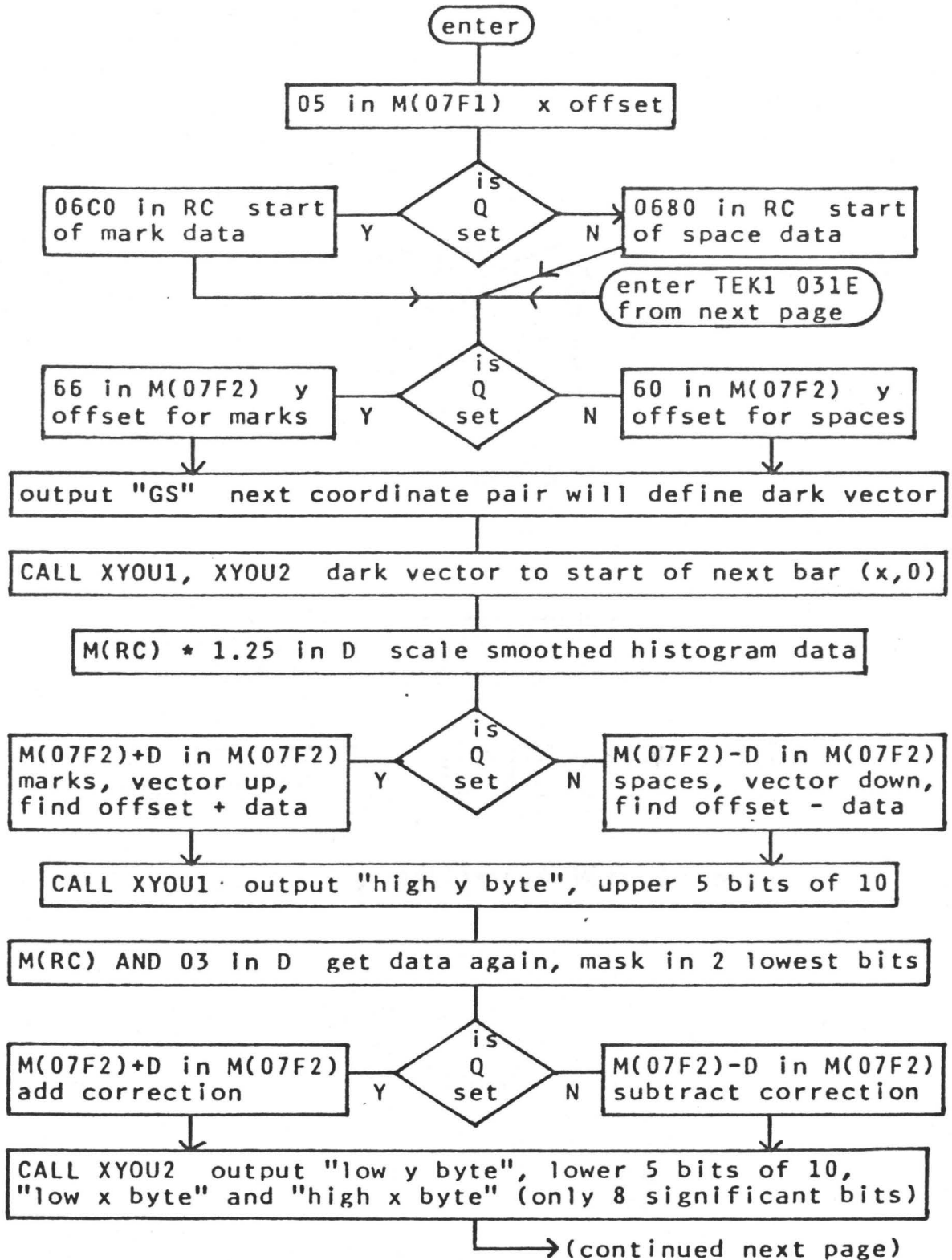


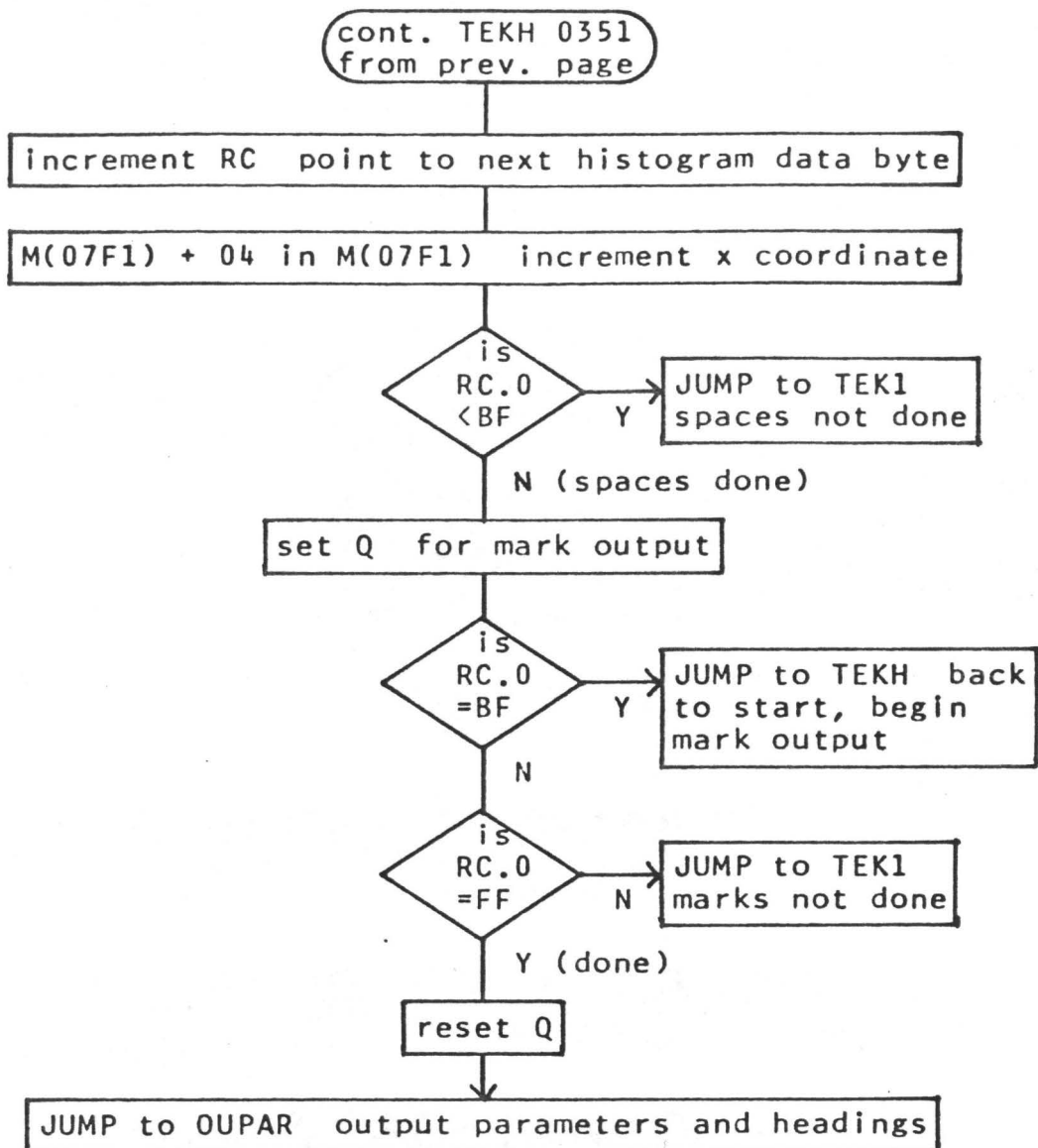
HIST 0250-036E



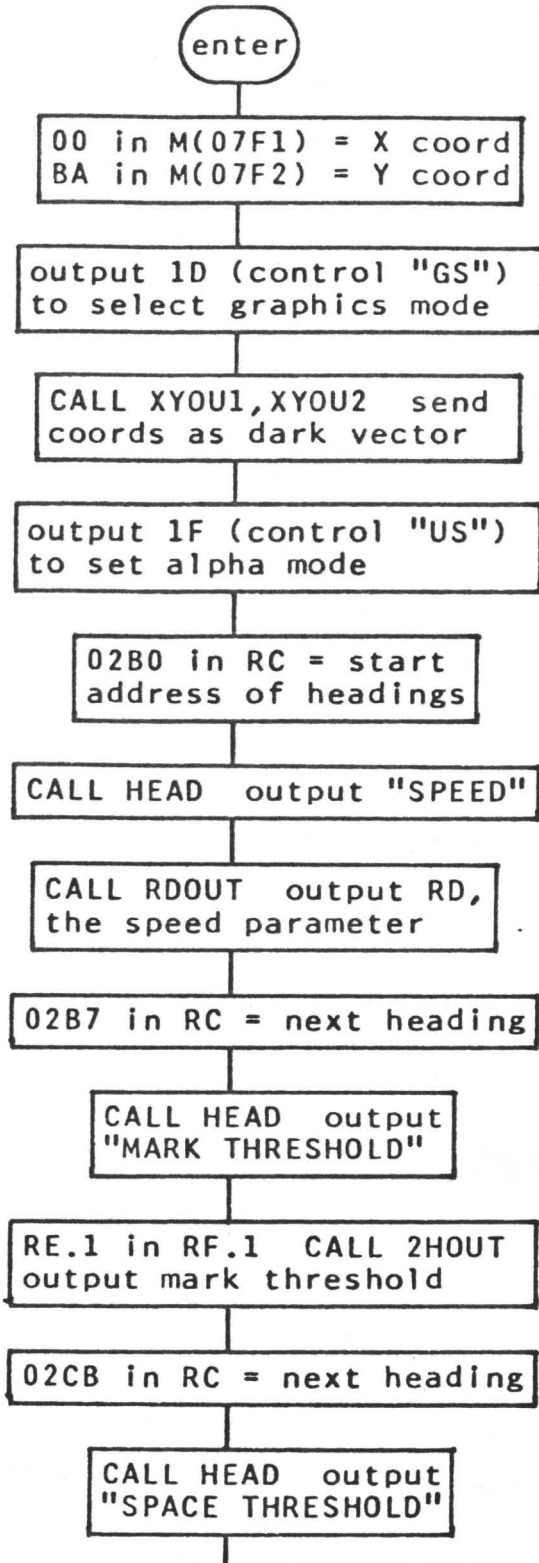
TEKH 0310-036E

entry: Q = 0, space data is output first

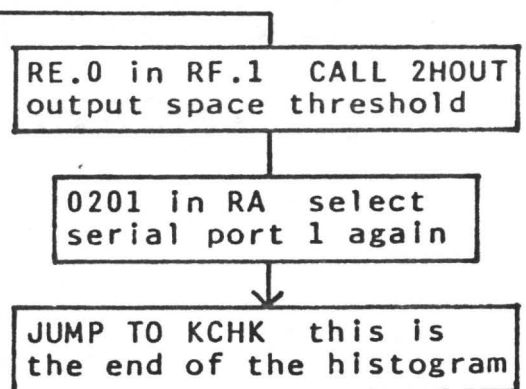
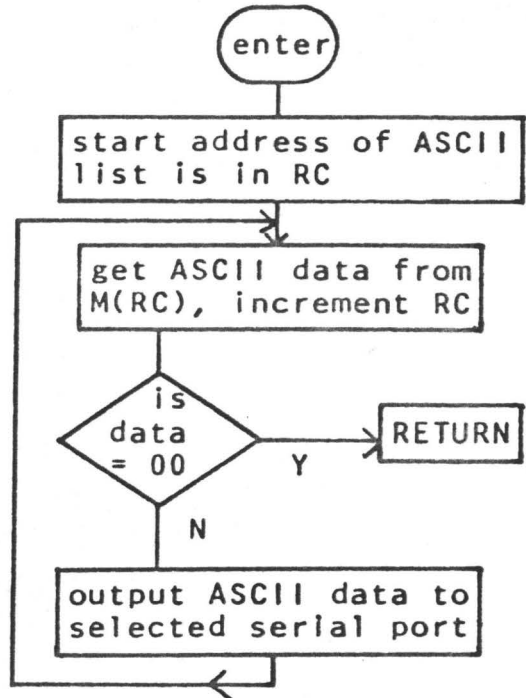




OUPAR 0370-03AF



HEAD 03B6-03BC



APPENDIX C

MONITOR PROGRAM

MONITOR OPERATION

Most of the monitor resides in EPROM from 0000 to 01FF. The two serial output drivers sit in RAM from 0200 to 021F and must be reloaded before use. To start the monitor hit "Reset" and "Run" on the front panel. The available commands are all single key commands which are immediately acted upon, no carriage return is required. All addresses or data are displayed on the 8 digit hexadecimal display above the ASCII keyboard. RD is displayed in the left four digits and RE in the right.

Key	Function
break	;Interrupts the processor and restarts the monitor at 0000. This may be used at any time to regain control from a bad program which has NOT used R1.
space	;Normally numeric data goes to RD, but after the space bar is pressed data goes to RE. After the next command data flow returns to RD.
0-9, A-F	;Enters hexadecimal numbers to RD or RE.
?	;The data in M(RD) is displayed in RE.0, RE.1=00
I	;Use after "?" to display successive locations.
!	;The data in RE.0 is entered into M(RD). If "!" is hit again, RD is advanced by 1 and RE.1=00. New data can now be entered in RE.
\$;Start execution of a program at M(RD) with R0 as RP. R3 should be made RP at the start of that program.
P	;All the data from M(RD) to M(RE) is moved down by one to M(RD+1) to M(RE+1). This is useful but any jump addresses must be corrected by hand.
G	;Any data from serial port 1 is dumped into M(RD) to M(RE). When (RE-RD) bytes have been collected or if there were any errors the program returns to the monitor. This is usually used for loading from a cassette.
W	;All the data from M(RD) to M(RE) is dumped out serial port 1. This is usually used to record a cassette.
R	;If the monitor was called as a subroutine by a running program this command is used to return to that program. This is useful for entering data during execution
TT	;This is a two key command. The data from M(RD) to M(RE) is output to serial port 1 as double hexadecimal ASCII characters, with each pair followed by "CR" & "LF". This is used to obtain listings on the Teletype.
TP	;This is the same as "TT" except that the address is output at the beginning and at every nnn0 address.

Many useful routines are hidden in the monitor and can be called by other programs. The whole monitor can be called, with return effected by "R". Some routines will return automatically if "Q" is set.

I/O PORTS & FLAG ASSIGNMENTS

The three "N" lines from the 1802 are decoded to seven lines (called "DEC N") to allow direct access to seven I/O ports.

Code	Mnemonic	Comments
61	OUT SER1	;This is serial port 1, a UART which drives an FSK cassette, 20 ma loop TTY, and EIA RS 232 interfaces. Baud rate is 110 or 600, format is 7 or 8 bits with optional parity, selectable by front panel switches. ;EF3 = 1 signifies transmitter busy, EF3 = 0, ready.
69	INP SER1	;This is serial input port 1, and it is configured the same way as the output port. The three input interfaces are OR'd, so only one should be used at time. Input and output may proceed simultaneously. UART error flags are OR'd onto a flag: EF1 = 1 means error, ;EF1 = 0, no error. EF2 = 0 means new data ready, EF2 = 1 ;no data. There is a front panel switch to connect EF2 to ;the DMA so tapes may be loaded with no bootstrap at 0000.
62	OUT POP1	;Parallel output port 1, a set of 8 leds on the CPU card and a dip connector.
6A	INP HEX	;This is the hex keyboard on the front panel. If the CPU is in the load mode, each pair of hex digits will be DMA'd into M(R0). It can be read with 6A.
63	OUT DISP	;Output to the hex led display. The data must be formatted by subroutine DISPY. This display is used extensively by all programs.
6B	INP ASKEY	;Input the ASCII encoded keyboard. ;EF4 = 1 means new data, EF4 = 0, no data.
64	OUT POP2	;This is a one bit (D0) output port used by the morse program to output the sampled key as a check on sampling rate. It is not essential.
6C	INP 4	;Not used.
65	OUT 5	;Not used.
6D	INP 5	;Not used.
66	OUT SER2	;Output data to serial port 2. This port has 20 ma TTY and EIA RS 232 interfaces. The baud can be set on the back panel from 75 to 4800.
6E	INP SER2	;Input data from serial port 2.
67	OUT CON2	;This sets the control register on serial port 2. D7=PI, D6=SBS, D5=WLS1, D4=WLS2, D3=PS. ;D2-D0 are not used. Output "F0" to set 8 data bits, no parity and 2 stop bits.
6F	INP PIP1	;D7=1 if serial port 2 is ready to send, D6=1 for new data ready, D5=1 for any errors. D4 and D3 are not used. D2 - D0 connect to the morse interface, D0=1 for key down.

LABELS

;Some labels are entry points of significance, but most
;are merely for internal branches and deserve no comment.

Add.	Label	Comments
0000	START	;Beginning of monitor, resets all required ;registers. The interrupt is not enabled ;until this is complete. Many routines and ;the interrupt come to this point.
0030	CHECK	;If Q=1, executes an SCRT RETURN from here.
0033	SIN	;Enter monitor here as a subroutine.
0036	DIN	;Enter here if input already in RF.1.
003E	S1	;
004B	S1B	;
0070	ATOF	;Jump here if monitor recognizes hex A to F
0074	NUM	;and here if 0 to 9.
007D	MOSH	;
008D	LRE	;
0098	SHOV?	;
00A4	LOGAN	;Here from look routine to examine next add.
00A5	LOOK	;Routine to examine memory. Not a subroutine.
00B8	RENM	;
00C0	MOD	;Routine to modify memory. Not a subroutine.
00DB	SOVER	;
00DC	SRX4	;SEP subroutine to shift D right four places.
00E2	HOVER	;
00E3	HXOUT	;SEP subroutine to output to hardware display
00F4	INPUT	;SCRT subroutine waits for input from serial ;port 1 or the ASCII keyboard. Data returned ;in M(RX) and RF.1.
00F9	KB?	;
00FC	GOTIT	;
0100	EXITC	;
0101	SUB	;SEP subroutine for subroutine calls by SCRT.
0112	EXITR	;
0113	RETRN	;SEP subroutine for SCRT subroutine exits. ;See RCA 1802 Users Manual for SEP and SCRT.
0120	WRITE	;Cassette or paper tape dump. Not subroutine.
0121	TFRE?	;
0134	PUSH	;Routine to make a hole. Not a subroutine.
0136	MPUSH	;
0147	PDONE	;
0150	GET	;Routine to input tapes. Not a subroutine.
0157	DELAY	;
015D	SER?	;
0161	GMORE	;
016C	MSER?	;
0170	GDONE	;

LABELS

Add.	Label	Comments
0178	TTY	;Routines for hex/ASCII dumps. Not subroutine
0188	MORTY	;
018B	NOTTP	;
01A1	NOADD	;
01CE	RDOUT	;SCRT subroutine outputs RD as 4 hex/ASCII.
01DA	CRLF	;SCRT subroutine outputs an ASCII "CR" & "LF"
01E2	2HOUT	;SCRT output subroutine, 2 hex/ASCII from RF.1
01F2	1HOUT	;SCRT output subroutine, 1 hex/ASCII from RF.1
0200	S1DON	;
0201	SOUT1	;SEP subroutine outputs D to serial port 1.
0202	S1OK?	;
020F	S2DON	;
0210	SOUT2	;SEP subroutine outputs D to serial port 2.
0214	T2OK?	;

MONITOR REGISTER ASSIGNMENTS

Register	Initial Value	Use
R0	0000	;Reset internally by 1802; initial ;program counter (RP)/ RP after "\$" command/ DMA pointer.
R1	0000	;RP after interrupt to restart mon- ;itor, reset before interrupt reenabled. To ensure correct ;interrupt action, R1 should NEVER have other uses.
R2	07E0	;Stack pointer for SCRT addresses.
R3	0007	;RP for most programs.
R4	0101	;RP for SEP routine CALL for SCRT.
R5	0113	;RP for SEP routine RETURN for SCRT.
R6	----	;Scratchpad for CALL and RETURN.
R7	07F0	;RX for most programs. 07F0 is a ;free location for I/O and R7 should point here when not ;in use. The area from 07E1 to 07FF is available as a stack ;for data, but only 07F0 is used by the monitor.
R8	00DC	;RP for SEP routine SRX4.
R9	----	;Scratchpad
RA	0201	;RP for SEP routine SOUT1.
RB	00E3	;RP for SEP routine HXOUT.
RC	--00	;Scratchpad, RC.0 is used as a flag.
RD	----	;Basically scratchpad, but since RD ;is output to the led display, it is used by many routines ;for addresses or data. The display should be updated each ;time RD is modified.
RE	----	;Same as RD.
RF	----	;Scratchpad.

MONITOR PROGRAM

Add.	Code	Label	Mnemonic	Comments
0000	F8 00	START	LDI 00	;Load 0007 in R3 for use as
0002	B3		PHI 3	;the program counter (RP).
0003	F8 07		LDI 07	;
0005	A3		PLO 3	;
0006	D3		SEP 3	;Make R3 = RP
0007	7A		REQ	;If monitor restarted from
0008	F8 00		LDI 00	;jump or interrupt, Q flag
000A	A1		PLO 1	;may need to be reset.
000B	B1		PHI 1	;Initialize the registers
000C	B8		PHI 8	;as required.
000D	AC		PLO C	;
000E	EB		PHI B	;0000 R1
000F	F8 01		LDI 01	;07E0 R2
0011	A4		PLO 4	;0101 R4
0012	B4		PHI 4	;0113 R5
0013	B5		PHI 5	;07F0 R7
0014	AA		PLO A	;00DC R8
0015	F8 02		LDI 02	;0201 RA
0017	BA		PHI A	;00E3 RB
0018	F8 E0		LDI E0	; 00 RC.0
001A	A2		PLO 2	;
001B	F8 DC		LDI DC	;
001D	A8		PLO 8	;
001E	F8 E3		LDI E3	;
0020	AB		PLO B	;
0021	F8 F0		LDI F0	;
0023	A7		PLO 7	;
0024	F8 07		LDI 07	;
0026	B2		PHI 2	;
0027	B7		PHI 7	;
0028	F8 13		LDI 13	;
002A	A5		PLO 5	;
002B	E3		SEX 3	;Make R3 = RX. This enables
002C	70 73		RET 73	;the RET to load the immed-
002E	C4		NOP	;iate byte and enable the
002F	C4		NOP	;interrupt. R3=RP, R7=RX

MONITOR PROGRAM

Add.	Code	Label	Mnemonic	Comments
0030	39 33	CHECK	BNQ SIN	;Jmp 1 if Q=0,else return
0032	D5		RETURN	;to calling program
0033	D4 00F4	SIN	SUB INPUT	;Get input from device
0036	9F	DIN	GHI F	;Input into D
0037	FF 20		SMI 20	;Subtract 20, 0 if "space"
0039	3A 3E		BNZ S1	;Jmp if not
003B	1C		INC C	;Make RC.0 > 0, data to RE
003C	30 30		BR CHECK	;Back for more input
003E	FF 01	S1	SMI 01	;Subtract 01, 0 if "!"
0040	32 C0		BZ MOD	;Jmp to modify memory if so
0042	FF 03		SMI 03	;Subtract 03, 0 if "\$"
0044	3A 4B		BNZ S1B	;Jmp if not
0046	8D		GLO D	;Put RD into R0, and begin
0047	A0		PLO 0	;execution of a new program
0048	9D		GHI D	;with R0 as the PC. The PC
0049	E0		PHI 0	;should be changed back to
004A	D0		SEP 0	;R3 by the new program.
004B	FF 1B	S1B	SMI 1B	;Subtract 1B, 0 if "?",
004D	3B 74		BM NUM	;If neg, was 0 to 9, input it
004F	32 A5		BZ LOOK	;Jmp to look at memory if "?"
0051	FF 08		SMI 08	;Subtract 08, 0 if "G"
0053	3B 70		BM ATOF	;If neg, was A to F, input it
0055	C2 0150		LBZ GET	;Jmp to GET if "G"
0058	FF 09		SMI 09	;Subtract 09, 0 if "P"
005A	C2 0134		LBZ PUSH	;Jmp to PUSH if "P"
005D	FF 02		SMI 02	;Subtract 02, 0 if "R"
005F	C6		LSNZ	;Skip 2 if not
0060	D5		RETURN	;Return to the program
0061	00		IDL	;which called the monitor
0062	FF 02		SMI 02	;Subtract 02, 0 if "T"
0064	C2 0178		LBZ TTY	;Jmp to TTY routines if "T"
0067	FF 03		SMI 03	;Subtract 03, 0 if "W"
0069	C2 0120		LBZ WRITE	;Jmp to write if "W"
006C	D1		SEP 1	;If no command recognized,
006D	00		IDL	;there was an error, so
006E	00		IDL	;restart monitor.
006F	00		IDL	;

MONITOR PROGRAM

Add.	Code	Label	Mnemonic	Comments
0070	F8 09	ATOF	LDI 09	;Add 09 to the ASCII from
0072	F4		ADD	;A to F to remove offset
0073	57		STR 7	;Restore it like 0 to 9
0074	F0	NUM	LDX	;Here if it was 0 to 9
0075	FE		SHL	;Shift left by 4 to remove
0076	FE		SHL	;last of offset. There is
0077	FE		SHL	;now a hex number in D4 to D7
0078	FE		SHL	;of the D register.
0079	57		STR 7	;Save it
007A	F8 04		LDI 04	;Load R9.0 with the number of
007C	A9		PLO 9	;shifts to be done.
007D	8C	MOSH	GLO C	;If RC.0=0, shift hex to RD
007E	3A 8D		BNZ LRE	;Else jmp to load RE
0080	8D		GLO D	;This section of code shifts
0081	FE		SHL	;RD left by one in a 16 bit
0082	AD		PLO D	;shift. This leaves a 0 in
0083	9D		GHI D	;the 1sb.
0084	7E		SHLC	;
0085	BD		PHI D	;Shift done here
0086	F0		LDX	;Get the hex character
0087	FE		SHL	;Shift the msb into DF
0088	3B 98		BNF SHOVS?	;If 0, check if shifting done
008A	1D		INC D	;Else inc RD, ie set 1sb to 1
008B	30 98		BR SHOVS?	;Check if shifting done
008D	8E	LRE	GLO E	;This is the same as 0080 to
008E	FE		SHL	;008A, but the hex is shifted
008F	AE		PLO E	;into RE because RC.0 > 0.
0090	9E		GHI E	;
0091	7E		SHLC	;
0092	BE		PHI E	;
0093	F0		LDX	;
0094	FE		SHL	;
0095	3B 98		BNF SHOVS?	;
0097	1E		INC RE	;
0098	57	SHOVS?	STR 7	;Save the shifted hex
0099	29		DEC R9	;Dec the shift counter
009A	89		GLO 9	;Get the counter
009B	3A 7D		BNZ MOSH	;Jmp for more if not 0
009D	D4 01 AB		SUB DISPY	;If done, display result,
00A0	30 30		BR CHECK	;then back for more.
00A1	00		IDL	;
00A2	00		IDL	;

MONITOR PROGRAM

Add.	Code	Label	Mnemonic	Comments
00A4	1D	LOGAN	INC D	;Point to next mem. location.
00A5	0D	LOOK	LDN D	;This is the beginning of the
00A6	AE		PLO E	;routine to examine memory.
00A7	F8 00		LDI 00	;Pu; the data in RE.0 and
00A9	BE		PHI E	;blank RE.1.
00AA	D4 01AB		SUB DISPY	;Display M(RD) in RE.0.
00AD	C5		LSNQ	;SKP 2 if Q=0, else since one
00AE	D5		RETURN	;location was displayed,
00AF	00		IDL	;return to calling program.
00B0	D4 00F4		SUB INPUT	;Go and get an input,
00B3	9F		GHI F	;put it into D.
00B4	FF 49		SMI 49	;Subtract 49, 0 if "I".
00B6	32 A4		BZ LOGAN	;If so, jump to look again.
00B8	F8 00	RENM	LDI 00	;If not, reset RC.0 (not
00BA	AC		PLO C	;always needed), and jump to
00BB	30 36		BR DIN	;monitor to check last input.
00ED	00		IDL	;
00BE	00		IDL	;MOD is the routine to modify
00BF	00		IDL	;a memory location. The piece
00C0	8E	MOD	GLO E	;of data in RE.0 is put into
00C1	5D		STR D	;M(RD).
00C2	F8 00		LDI 00	;Blank RE.1.
00C4	BE		PHI E	;
00C5	D4 01AB		SUB DISPY	;Display that.
00C8	C5		LSNQ	;SKP 2 if Q=0, else since one
00C9	D5		RETURN	;location was modified,
00CA	00		IDL	;return to calling program.
00CB	D4 00FA		SUB INPUT	;Go and get an input,
00CE	9F		GHI F	;put it into D.
00CF	FF 21		SMI 21	;Subtract 21, 0 if "!".
00D1	3A B8		BNZ RENM	;If not, back to monitor,
00D3	1D		INC D	;else point to next location
00D4	D4 01AB		SUB DISPY	;and display that address.
00D7	30 30		BR CHECK	;Go back to the monitor and
00D9	00		IDL	;enter data for that location
00DA	00		IDL	;Note that data goes to RE.
00DB	D3	SOVER	SEP 3	;Return to calling program.
00DC	F6	SRX4	SHR	;This is an SEP subroutine
00DD	F6		SHR	;which shifts the D register
00DE	F6		SHR	;right by four places. It
00DF	F6		SHR	;puts 0's into D4 to D7.
00E0	30 30		BR SOVER	;Over and out

MONITOR PROGRAM

Add.	Code	Label	Mnemonic	Comments
00E2	D3	HOVER	SEP 3	;Return to calling program
00E3	FA 0F	HXOUT	ANI 0F	;This is an SEP subroutine
00E5	FB 0F		XRI 0F	;which outputs a single hex
00E7	57		STR 7	;digit to the correct place
00E8	89		GLO 9	;in the led display. Enter
00E9	FC 10		ADI 10	;with the digit in D0 to D3
00EB	A9		PLO 9	;of D, and the digit # minus
00EC	F1		OR	;one in D4 to D7 of R9.0, so
00ED	57		STR 7	;F0 is leftmost digit, and
00EE	63		OUT DISP	;60 is the rightmost digit.
00EF	27		DEC 7	;This digit # is upped by 10.
00F0	30 E2		BR HOVER	;Over and out. Note that this
00F2	00		IDL	;routine depends heavily on
00F3	00		IDL	;the display hardware.
;This is the input subroutine. It loops until it finds an				
;input from the ASCII keyboard or serial interface 1.				
00F4	35 F9	INPUT	B2 KB?	;Jump if no serial input
00F6	69		INP SER1	;Else input it
00F7	30 FC		BR GOTIT	;Jump a bit
00F9	3F F4	KB?	BN4 INPUT	;Loop if no keyboard either.
00FB	6B		INP ASKEY	;Else input it
00FC	BF	GOTIT	PHI F	;Salt it away
00FD	D5		RETURN	;Over and out
00FE	00		IDL	;
00FF	00		IDL	;
;This is the SUB subroutine. It is used to call subroutines				
;via the RCA SCRT (Standard Call and Return) conventions.				
0100	D3	EXITC	SEP 3	;Go to the called subroutine.
0101	E2	SUB	SEX 2	;R2 points to the return
0102	96		GHI 6	;address stack.
0103	73		STXD	;Push R6 onto stack and
0104	86		GLO 6	;leave stack at a free
0105	73		STXD	;location.
0106	93		GHI 3	;Copy old RP (R3) into R6
0107	B6		PHI 6	;to save it.
0108	83		GLO 3	;R6 now points to two byte
0109	A6		PLO 6	;inline address in calling
010A	46		LDA 6	;program. Get it and put it
010B	B3		GHI 6	;into R3 as it is the start
010C	46		LDA 6	;of the called routine.
010D	A3		PLO 3	;
010E	E7		SEX 7	;Put R7 back as RX as R2 only
010F	30 00		BR EXITC	;points to the subroutine
0111	C0		IDL	;address stack.

MONITOR PROGRAM

This is the RETURN subroutine. It is used to return from a subroutine called by standard RCA SCRT techniques. This and the above routines are not used for SEP subroutines.

Add.	Code	Label	Mnemonic	Comments
0112	D3	EXITR	D3	;Return to calling program.
0113	96	RETRN	GHI 6	;Copy R6 into R3 so it will
0114	B3		PHI 3	;contain the return program
0115	86		GLO 6	;counter.
0116	A3		PLO 3	;
0117	E2		SEX 2	;R2 points to the return
0118	60		IRX	;address stack. Decrement
0119	72		LDXA	;it from the free location so
011A	A6		PLO 6	;it points to the previously
011B	F0		LDX	;stored value of R6 and put
011C	B6		PHI 6	;it back into R6.
011D	E7		SEX 7	;Put R7 back as RX.
011E	30 12		BR EXITR	;Over and out

;This is the "W" (Write) routine. It is used to dump ;memory locations via serial output port 1. The data is ;sent out as a stream of 8 bit bytes, but final format ;depends on the UART control switches. It is not converted ;to hex or ASCII, so this is the routine for writing to ;cassette. The data from M(RD) to M(RE) is sent out.

0120	ED	WRITE	SEX D	;RD is used for output (RX).
0121	36 21	TFRE?	B3 TFRÉ?	;Loop 'till transmitter free.
0123	61		OUT SER1	;Output byte and advance RD.
0124	E7		SEX 7	;R7 = RX for display.
0125	D4 01AB		SUB DISPY	;Display advanced address.
0128	D4 01C4		SUB SRERD	;Sub for 16 bit RE-RD and
012B	33 20		BGE WRITE	;back for more if RE>=RD.
012D	C5		LSNQ	;SKp 2 if Q=0, else this was
012E	D5		RETURN	;called as a subroutine, so
012F	00		IDL	;return to calling program.
0130	D1		SEP 1	;Done, so back to monitor.
0131	00		IDL	;
0132	00		IDL	;
0133	00		IDL	;

MONITOR PROGRAM

;This is the "P" (Push) routine. All the data from M(RD) to
;M(RE) is pushed down by one location. This is useful for
;inserting an instruction, but all jumps must be corrected.

Addr.	Code	Label	Mnemonic	Comments
0134	4D	PUSH	LDA D	;Get the first byte & advance
0135	B9		PHI 9	;Save it.
0136	0D	MPUSH	LDN D	;Get the second byte.
0137	A9		PLO 9	;Save that too.
0138	99		CHI 9	;Get the first one again.
0139	5D		STR D	;Put it in new location.
013A	89		GLO 9	;Get the second one again.
013B	B9		PHI 9	;Save it as new first byte.
013C	D4 01AB		SUB DISPY	;Display advanced address.
013F	D4 01C4		SUB SRERD	;Sub for 16 bit RE-RD.
0142	3B 47		BL PDONE	;Jump ahead if done
0144	1D		INC D	;Else increment RD
0145	30 36		BR MPUSH	;and go back for more.
0147	C5	PDONE	LSNQ	;Skp 2 if Q=0, else this was
0148	D5		RETURN	;called as a subroutine, so
0149	00		IDL	;return to calling program.
014A	D1		SEP 1	;Done, so back to monitor.
014B	00		IDL	;IDL from 014B to 014F inc.

;This is the "G" (Get) routine. It is used to input a
;string of data bytes into memory from M(RD) to M(RE) from
;serial port 1. Parity depends on the front panel switches.

;This is the routine for loading a cassette or paper tape.

0150	69	GET	INP SER1	;Input to clear UART flags.
0151	F8 FF		LDI FF	;Load FFFF into R9 for use as
0153	B9		PHI 9	;a timing constant. A delay
0154	F8 FF		LDI FF	;is used so the cassette
0156	A9		PLO 9	;interface can stabilize.
0157	29	DELAY	DEC 9	;Decrement timing constant.
0158	B9		CHI 9	;Get the high 8 bits.
0159	C4		NOP	;Extra delay.
015A	C4		NOP	;
015B	3A 57		BNZ DELAY	;Loop until R9.1 = 00.
015D	35 5D	SER?	B2 SER?	;Tiny loop until data appears
015F	34 50		B1 GET	;Start again if error.
0161	69	GMORE	INP SER1	;Get some data (finally!!)
0162	5D		STR D	;and store it.
0163	1D		INC D	;Increment memory pointer.
0164	D4 01AB		SUB DISPY	;Display pointer.
0167	D4 01C4		SUB SRERD	;Sub for 16 bit RE - RD
016A	3B 70		BL GDONE	;and jump ahead if done.
016C	35 6C	MSER?	B2 MSER?	;Wait for more serial.
016E	3C 61		BN1 GMORE	;Back for more if no error.
0170	C5	GDONE	LSNQ	;Skp 2 if Q=0, else this was
0171	D5		RETURN	;called as a subroutine, so
0172	00		IDL	;return to calling program.
0173	D1		SEP 1	;All done, back to monitor

MONITOR PROGRAM

;These are the two teletype routines, "TT" and "TP".
 ;TT takes sequential 8 bit memory locations and outputs
 ;them as two hex ASCII characters - so 11000101 would
 ;be sent out as C5. Each pair of characters is followed by
 ;a line feed/carriage return. TP is similar except that the
 ;starting address and every nn00 address is also output.
 ;These routines are used for teletype listings. The memory
 ;from M(RD) to M(RE) is output.

Addr.	Code	Label	Mnemonic	Comments
0174	00		IDL	;
0175	00		IDL	;
0176	00		IDL	;
0177	00		IDL	;
0178	D4 00F4	TTY	SUB INPUT	;Get second letter of inst-
017B	9F		GHI F	;ruktion and put it into D.
017C	FF 50		SMI 50	;Subtract 50, 0 if "P", put
017E	BC		PHI C	;result in RC.1 for flag use.
017F	D4 01DA		SUB CRLF	;Output a CR and LF.
0182	9C		GHI C	;Get flag and
0183	3A 8B		BNZ NOTTP	;jump a bit if not TP.
0185	D4 01CE		SUB RDOUT	;Output RD to serial port 1.
0188	D4 01DA	MORTY	SUB CRLF	;Output a CR and LF.
018B	4D	NOTTP	LDA D	;Get the data, put it in RF.1
018C	BF		PHI F	;and advance the pointer.
018D	D4 01E2		SUB 2HOUT	;Sub to output two characters
0190	D4 01AB		SUB DISPY	;Display advanced pointer.
0193	9C		GHI C	;Get flag and
0194	3A A1		BNZ NOADD	;don't output add. if not TP.
0196	8D		GLO D	;Get low 8 bits of pointer,
0197	FA 0F		ANI 0F	;mask in lower 4 bits only.
0199	3A A1		BNZ NOADD	;Jump if address not = nn00.
019B	D4 01DA		SUB CRLF	;Output a CR and LF.
019E	D4 01CE		SUB RDOUT	;Output RD.
01A1	D4 01C4	NOADD	SUB SRERD	;Sub for 16 bit RE - RD.
01A4	33 8B		BGE MORTY	;Go back, for more if not done
01A6	C5		LSNQ	;SKp 2 if Q=0, else this was
01A7	D5		RETURN	;called as a subroutine, so
01A8	00		IDL	;retrun to calling program.
01A9	D1		SEP 1	;All done, back to monitor
01AA	00		IDL	;

MONITOR PROGRAM

;This is the display subroutine. When it is called the
 ;current contents of RD and RE are output to an eight digit
 ;hexadecimal led display. The display is self scanning and
 ;refresh by the program is not required. The display should
 ;be updated whenever RD or RE is changed.

Add.	Code	Label	Mnemonic	Comments
01AB	F8 F0	DISPY	LDI F0	;Preset R9.0 for use by HXOUT
01AD	A9		PLO 9	;
01AE	9D		GHI D	;Get the high 8 bits of RD.
01AF	DB		SEP 8	;Sub to shift right by four.
01B0	DB		SEP B	;Sub to output first (left-
01B1	9D		GHI D	;most) digit, get RD.1 again.
01B2	DB		SEP B	;Output second digit.
01B3	8D		GLO D	;Get the low 8 bits of RD.
01B4	DB		SEP 8	;Shift right by four.
01B5	DB		SEP B	;Output third digit.
01B6	8D		GLO 8	;Get RD.0 again.
01B7	DB		SEP B	;Output fourth digit.
01B8	9E		GHI E	;Get the high 8 bits of RE.
01B9	DB		SEP 8	;Shift right by four.
01BA	DB		SEP B	;Output fifth digit.
01BB	9E		GHI E	;Get RE.1 again.
01BC	DB		SEP B	;Output sixth digit.
01BD	8E		GLO E	;Get the low 8 bits of RE.
01BE	DB		SEP 8	;Shift right by four.
01BF	DB		SEP B	;Output seventh digit.
01C0	8E		GLO E	;Get RE.0 again.
01C1	DB		SEP B	;Output eighth digit.
01C2	D5		RETURN	;Return to calling program.
01C3	00		IDL	;

;This routine performs a 16 bit subtraction, RE - RD.
 ;Only the sign, in DF, is returned.

01C4	8D	SRERD	GLO D	;Get low 8 bits of RD and put
01C5	57		STR 7	;it in a free location.
01C6	8E		GLO E	;Get low 8 bits of RE.
01C7	F7		SM	;D-M(RX) = RE.0-RD.0.
01C8	9D		GHI D	;Get high 8 bits of RD and
01C9	57		STR 7	;put it in a free location.
01CA	9E		GHI E	;Get high 8 bits of RE.
01CB	77		SMB	;RE.1 - RD.1 - Borrow.
01CC	D5		RETURN	;Return to calling program.
01CD	00		IDL	;

MONITOR PROGRAM

;This is a routine to output RD as 4 ASCII characters.

Add.	Code	Label	Mnemonic	Comments
01CE	9D	RDOUT	GHI D	;Get the high 8 bits of RD
01CF	BF		PHI F	;and put them in RF.1
01D0	D4 01E2		SUB 2HOUT	;Sub to output two characters
01D3	8D		GLO D	;Get the low 8 bits of RD
01D4	BF		PHI F	;and put them in RF.1.
01D5	D4 01E2		SUB HOUT	;Sub to output two characters
01D8	D5		RETURN	;Return to calling program.
01D9	00		IDL	;

;This routine outputs a "CARRIAGE RETURN" & "LINEFEED".

01DA	F8 0D	CRLF	LDI 0D	;Load ASCII for "CR".
01DC	DA		SEP A	;Sub to output D.
01DD	F8 0A		LDI 0A	;Load ASCII for "LF".
01DF	DA		SEP A	;Sub to output D.
01E0	D5		RETURN	;Return to calling program.
01E1	00		IDL	;

;This routine outputs two ASCII hexadecimal characters from
;an 8 bit number in RF.1.

01E2	9F	2HOUT	GHI F	;Get data passed from program
01E3	B9		PHI 9	;Hide it somewhere.
01E4	D8		SEP 8	;Sub to shift right by four.
01E5	BF		PHI F	;Put it in RF.1.
01E6	D4 01F2		SUB 1HOUT	;Sub to output one character.
01E9	99		GHI 9	;Get original data again and
01EA	FA 0F		ANI 0F	;mask in lower four bits.
01EC	BF		PHI F	;Put it in RF.1.
01ED	D4 01F2		SUB 1HOUT	;Sub to output one character.
01F0	D5		RETURN	;Return to calling program.
01F1	00		IDL	;

;This routine outputs one ASCII hex character from the
;lower four bits of RF.1.

01F2	9F	1HOUT	GHI F	;Get data passed from program
01F3	FF 0A		SMI 0A	;Subtract 0A, <0 if 0 to 9.
01F5	9F		GHI F	;Get data again, but leave DF
01F6	C7		LSNF	;Skp 2 if neg, was 0 to 9.
01F7	FC 07		ADI 07	;Add partial offset if A to F
01F9	FC 30		ADI	;Add offset for 0 to F.
01FB	DA		SEP A	;Sub to output D.
01FC	D5		RETURN	;Return to calling program.

MONITOR PROGRAM

;This is the subroutine to output the contents of D to
;serial port 1. It does not modify any registers, but
;does require RX to point to a free memory location.

Add.	Code	Label	Mnemonic	Comments
01FD	00		IDL	;
01FE	00		IDL	;Note that RX must be R7
01FF	00		IDL	;on entry.
0200	D3	S1DON	SEP 3	;Return to calling program.
0201	57	SOUT1	STR 7	;Put D into memory.
0202	36 02	S10K?	B3 S10K?	;Loop 'till transmitter free.
0204	61		OUT SER1	;Output data.
0205	27		DEC 7	;Restore RX value.
0206	30 00		BR S1DON	;Done, so get out.
0208	00		IDL	;IDL's (00) until 020F.

;This is the routine to output the contents of D to serial
;port 2. If RA is set to 0210 this routine will be used
;instead of SOUT1 whenever an SEP A (DA) is executed. This
;port is harder to use as the status must be read in with
;an input instruction and cannot be checked with the flags.
;It is more flexible as the control register is set with
;software rather than switches, but because of this the
;control register must be set before calling this program.
;This routine is not used by the monitor. RX must be R7 and
;point to 07F0 on entry. Location 07FF must be free.

020F	D3	S2DON	SEP 3	;Return to calling program.
0210	57	SOUT2	STR 7	;Save D in 07F0.
0211	F8 FF		LDI FF	;Point to a free location.
0213	A7		PLO 7	;
0214	6F	T20K?	INP PIP1	;Input parallel port 1.
0215	FE		SHL	;Put transmitter status bit
0216	3B 14		BNF T20K?	;into DF, loop until ready.
0218	F8 F0		LDI F0	;Point to stored D in normal
021A	A7		PLO 7	;I/O location and
021B	66		OUT SER2	;output data.
021C	27		DEC 7	;Restore RX value and
021D	30 0F		BR S2DON	;get out.
021F	00		IDL	;

APPENDIX D

MORSE DECODING PROGRAM

MORSE PROGRAM

;This is the Morse code decoding program. Program segments
;are more in logical than numerical order.

Addr.	Code	Label	Mnemonic	Comments
03DE	F8 00	RSTRT	LDI 00	;Program restarts here, RC.0
03DF	AC		PLO C	;is a monitor flag, reset it.
03E0	F8 03	MORSE	LDI 03	;Program starts here. R3 is
03E2	B3		PHI 3	;set for use as RP.
03E3	F8 E6		LDI E6	;
03E5	A3		PLO 3	;
03E6	D3		SEP 3	;Make R3 be RP.
03E7	E3		SEX 3	;Make R3 be RX as well. This
03E8	67		OUT CON2	;will output "F0" to serial
03E9	F0		(data)	;port 2 to set the UART.
03EA	E7		SEX 7	;Make R7 be RX.
03EB	7A		REQ	;Reset the Q flag.
03EC	D4 0033		SUB SIN	;Call the monitor as a sub-
				routine. Enter a 4 hex digit
				initial speed in the lefthand
				display (typ. 0000 to 0300),
				then type "R" to return.
03EF	F8 10	RSTR2	LDI 10	;Preset RE.1 and RE.0, the
03F1	AE		PLO E	;mark and space thresholds
03F2	BE		PHI E	;respectively, to "10".
03F3	D4 03C0		SUB MASC	;Sub to zero histograms, Q=0.
03F6	D4 01DA	RSTR1	SUB CRLF	;Output a "CR" & "LF".
03F9	F8 00		LDI 00	;Reset various registers:
03FB	AC		PLO C	;RC.0 time of keyup
03FC	BC		PHI C	;RC.1 # of keydowns in letter
03FD	A9		PLO 9	;R9.0 time of keydown
03FE	B9		PHI 9	;R9.1 Morse shift register
03FF	7A		REQ	;Reset the Q flag.
0400	9D	TIME	GHI D	;This is the start of the
0401	BF		PHI F	;basic timing loop. The
0402	8D		GLO D	;timing constant is moved
0403	AF		PLO F	;from RD to RF.
0404	8F	TIM1	GLO F	;Check RF.0,
0405	32 0A		BZ TIM2	;jump ahead if zero,
0407	2F	TIM3	DEC F	;else decrement RF,
0408	30 04		BR TIM1	;and jump back for more.
040A	9F	TIM2	GHI 9	;Check RF.1, and jump back
040B	3A 07		BNZ TIM3	;if not also zero,
040D	6F		INP PIP1	;else input Morse port, and
040E	64		OUT POP2	;output it to a one bit (D0)
040F	27		DEC 7	;port to check sampling.

MORSE PROGRAM

Add.	Code	Label	Mnemonic	Comments
0410	FA 01		ANI 01	;Mask in lsb, 0=up, 1=down.
0412	32 2C		BZ UP	;Jump if key up.
0414	8C		GLO 8	;If keyup timer zero, then
0415	32 21		BZ STDN	;key was down before, so jmp.
0417	FF 40		SMI 40	;If timer = 40, timer has
0419	32 21		BZ STDN	;reached limit, so jump.
041B	D4 0540		SUB UBIN	;Else sub to bin key up time
041E	C1 0568		L3Q AUT?	;If Q was set by BIN, then
;a histogram bin was full, so jump to AUT? for action.				
0421	F8 00	STDN	LDI 00	;Reset key up timer.
0423	AC		PLO C	;
0424	89		GLO 9	;Get key down timer,
0425	FF 40		SMI 40	;subtract 40, zero if full,
0428	32 00		BZ TIME	;so jump back to timer.
042A	19		INC 9	;Otherwise increment key
042A	30 00		BR TIME	;down counter and go to timer
042C	8C	UP	GLO C	;Here if key up, get timer.
042D	3A 58		BNZ STUP	;Jump if 0, was up before.
042F	9E		GHI E	;Else get mark threshold.
0430	57		STR 7	;Put it in free location.
0431	89		GLO 9	;Get key down timer, find
0432	F7		SM	;key down - mark threshold.
0433	33 3D		BGE DASH	;Jump if >= 0, it was a dash.
0435	99		GHI 9	;Get Morse shift register.
0436	FE		SHL	;Shift left, 0=dot into lsb.
0437	E9		PHI 9	;Put it back.
0438	89		GLO 9	;Get timer AGAIN!!!
0439	FF 08		SMI 08	;Subt. 08, 0 if perfect dot.
043B	30 45		BR SPD?	;Jump ahead to check speed.
043D	99	DASH	GHI 9	;Get shift register.
043E	FE		SHL	;Shift left, 0 into lsb.
043F	FC 01		ADI 01	;Add 01, 1=dash into lsb.
0441	E9		PHI 9	;Put it back.
0442	89		GLO 9	;Get timer again.
0443	FF 18		SMI 18	;Subt. 18, 0 if perfect dash.
;At this point, the D register is zero if the dot or dash				
;was of a perfect length. DF=1 if the speed is too fast				
;(dot > 08 or dash > 18), and DF=0 if too slow.				
0445	32 4B	SPD?	BZ UP1	;Jump if D=0, no change.
0447	1D		INC D	;Inc.(slower) timing constant
0448	CF		LSDF	;Skip 2 if DF=1, speed was
0449	2D		DEC D	;fast, else dec. by 2, effec-
044A	2D		DEC D	;tively by 1, raise speed.
044B	D4 0545	UP1	SUB DBIN	;Sub to bin key down.
044E	39 54		BNQ UP2	;Jmp if Q=0, no bins full,
0450	7A		REQ	;else reset Q,
0451	C0 0568		LBR AUT?	;and jump for action.

MORSE PROGRAM

Add.	Code	Label	Mnemonic	Comments
0454	9C	UP2	GHI C	;Get counter for number of
0455	FC 01		ADI 01	;Key downs in a letter incre-
0457	BC		PHI C	;ment and restore it.
0458	F8 00	STUP	LDI 00	;Reset key down timer.
045A	A9		PLO 9	;
045B	8C		GLO C	;Get key up timer,
045C	FF 40		SMI 40	;subt. 40, 0 if timer full,
045E	32 79		BZ UP3	;jump if so.
0460	9C		GHI 9	;If RC.0=0, there are no
0461	32 69		BZ UP4	;elements in letter, so jump.
0463	8E		GLO E	;Get space threshold, and put
0464	57		STR 7	;it in a free location.
0465	8C		GLO C	;Get key up timer and find
0466	F7		SM	;key up - space threshold.
0467	32 80		BZ LOUT	;If 0, now output letter.
0469	8E	UP4	GLO E	;Find inter-character space:
046A	F6		SHR	;
046B	57		STR 7	;ICS = 4*RE.0 - RE.0/2
046C	8E		GLO E	; = 3.5*RE.0
046D	FE		SHL	;For perfect morse it should
046E	FE		SHL	;be 3.0, but hand sent is
046F	F7		SM	;better with this ratio.
0470	57		STR 7	;Put ICS in free location.
0471	8C		GLO C	;Get key up timer and find
0472	F7		SM	;key up - ICS
0473	3A 78		BNZ UP5	;Jump if not zero, no space.
0475	F8 20		LDI 20	;Load ASCII "space".
0477	DA		SEP A	;Output to serial port 1
0478	1C	UP5	INC C	;Increment key up timer.
0479	3F 00	UP3	BN4 TIME	;Jump to timer if no key-
047B	30 B8		BR KCHK	;board input, else go and
047D	00		IDL	;see what it is.
047E	00		IDL	;Following section outputs
047F	00		IDL	;a letter.
0480	99	LOUT	GHI 9	;Get Morse shift register,
0481	57		STR 7	;put it in a free location,
0482	62		OUT POP1	;output it to 8 leds for a
0483	27		DEC 7	;visual check of dots/dashes.
0484	9C		GHI C	;Check if 6 key downs in
0485	FF 06		SMI 06	;letter, jump if not for
0487	3A 9A		BNZ N06	;standard decoding.

MORSE PROGRAM

;The next segment of code is for Morse characters with
;exactly six elements as some require different decoding.

Add.	Code	Label	Mnemonic	Comments
0489	99		GHI 9	;Get Morse shift register and
048A	FF 4C		SMI 4C	;subtract Morse "?".
048C	3A 92		BNZ NO?	;Jump if not that,
048E	F8 3F		LDI 3F	;else load ASCII "?" and
0490	30 AB		BR MOUT	;jump to output it.
0492	FF 2C	NO?	SMI 2C	;Subtract further offset,
0494	3A A0		BNZ OKLET	;jump if not 0,
0496	F8 3A		LDI 3A	;else load ASCII ":" and
0498	30 AB		BR MOUT	;jump to output it.

;The remaining 6 element characters are handled normally.
;At this point the D register holds # of key downs - six.

049A	3B A0	NO6	BL OKLET	;If less than six, ok, but if
049C	F8 2A		LDI 2A	;not load ASCII "*" for over-
049E	30 AB		BR MOUT	;run and jump to output it.
04A0	99	OKLET	GHI 9	;Get Morse shift register,
04A1	FA 3F		ANI 3F	;mask in lower 6 bits, put
04A3	AF		PLO F	;it in RF.0 as low address.
04A4	F8 05		LDI 05	;Put high address in RF.1,
04A6	BF		PHI F	;point to Morse/ASCII table.
04A7	0F		LDN F	;Get ASCII
04A8	DA	MOUT	SEP A	;and output it.
04A9	1C		INC C	;Increment key up timer.
04AA	27		DEC 7	;Point to free memory (07EF).
04AB	8C		GLO C	;Get key up timer,
04AC	57		STR 7	;save it in memory, and point
04AD	60		IRX	;back to I/O location (07F0).
04AE	C0 0600		LBR THRM	;Jump to find new thresholds.
04B1	00		IDL	;
04B2	00		IDL	;
04B3	00		IDL	;
04B4	00		IDL	;
04B5	00		IDL	;
04B6	00		IDL	;
04B7	00		IDL	;

MORSE PROGRAM

;The next segment of code checks for input from the ASCII
;keyboard and acts accordingly.

Add.	Code	Label	Mnemonic	Comments
04B8	3F B8	KCHK	BN4 KCHK	;Loop until keyboard input.
04BA	6B		INP ASKEY	;Input keyboard into D, M(RX)
04BB	FF 2F		SMI 2F	;Subtract ASCII "/".
04BD	C2 03DD		LBZ RSTRT	;Complete restart if "/"
04C0	FF 19		SMI 19	;Subtract again, 0 if "H".
04C2	C2 0250		LBZ HIST	;Output histogram if "H".
04C5	FF 03		SMI 03	;Subtract again, 0 if "K".
04C7	3A D7		BNZ NOKO	;Jump if not, else point
04C9	F8 10		LDI 10	;register A to output
04CB	AA		PLO A	;routine for serial port 2.
04CC	F8 1B		LDI 1B	;Output "ESC" & "ETB" to
04CE	DA		SEP A	;serial port 2. This sends
04CF	F8 17		LDI 17	;a copy command to the
04D1	DA		SEP A	;Tektronix Graphics Terminal.
04D2	F8 01		LDI 01	;Point register A back to the
04D4	AA		PLO A	;routine for serial port 1.
04D5	30 B8		BR KCHK	;Back for more input.
04D7	FF 02	NOKO	SMI 02	;Subtract again, 0 if "M".
04D9	C2 03F6		LBZ RSTR1	;If so partial restart.
04DC	FF 01		SMI 01	;Subtract again, 0 if "N".
04DE	32 F6		BZ SAUT	;If so, jump to cancel auto.
04E0	FF 0A		SMI 0A	;Subtract again, 0 if "X".
04E2	7A		REQ	;Reset Q (will clear hist.)
04E3	32 EA		BZ GOMAS	;If "X", go clear histograms.
04E5	FF 01		SMI 01	;Subtract again, 0 if "Y".
04E7	3A F1		BNZ Z?	;Jump if not, else
04E9	7B		SEQ	;set Q (will histogram/2)
04EA	D4 03C0	GOMAS	SUB MASC	;and divide hist. data by 2.
04ED	7A		REQ	;Reset Q (X & Y return here)
04EE	C0 03F6		BR RSTR1	;and do partial restart.
04F1	FF 01	Z?	SMI 01	;Subtract AGAIN!, 0 if "Z".
04F3	CA 0000		LBNZ START	;If not, invalid, to monitor.
04F6	27	SAUT	DEC 7	;Here if "Z" or "N", point to
04F7	27		DEC 7	;location for auto command,
04F8	6B		INP ASKEY	; (07EE) and put input there.
04F9	60		IRX	; "Z" for auto, "N" nonauto.
04FA	60		IRX	;Point back to I/O location.
04FB	30 00		BR TIME	;Go back to timing loop.
04FD	00		IDL	;
04FE	00		IDL	;
04FF	00		IDL	;

MORSE PROGRAM

```

;This is the lookup table to convert Morse to ASCII.
;A register is preset to 0000 0001 and the elements are
;shifted into the lsb, 0 for dot and 1 for dash. The
;leading 1 shows the beginning of the character. The letter
;"L" is ".-..."and would appear as 0001 0100. The upper two
;bits are set to 00. This forms the lower eight bits of the
;address, and 05 is the offset for the high eight bits.
;This scheme is memory efficient, but causes characters
;with six elements to show up anywhere in the table.
;This is only a problem for "?" and ":", so these are dealt
;with in software. "@" is used for an invalid character,
;and "<" for the Morse error ".....".

```

Add.	Data	Character	Add.	Data	Character
0500	3C	<	0520	35	5
0501	40	@	0521	34	4
0502	45	E .	0522	40	@
0503	54	T -	0523	33	3
0504	49	I ..	0524	40	@
0505	41	A .-	0525	40	@
0506	4E	N -. ,	0526	40	@
0507	4D	M --	0527	32	2
0508	53	S ...	0528	40	@
0509	55	U ..-	0529	40	@
050A	52	R .-. ,	052A	3B	;
050B	57	W .--	052B	40	@
050C	44	D -..	052C	40	@
050D	4B	K -.-	052D	5B	[
050E	47	G --. ,	052E	40	@
050F	4F	O ---	052F	31	1
0510	48	H	0530	36	6
0511	56	V ...-	0531	5F	-
0512	46	F ... ,	0532	2F	/
0513	40	@	0533	2C	,
0514	4C	L .-..	0534	40	@
0515	2E	. -.-.-	0535	40	@
0516	50	P .-- ,	0536	40	@
0517	4A	J .---	0537	40	@
0518	42	B -... ,	0538	37	7
0519	58	X -.- ,	0539	40	@
052A	43	C -. ,	053A	40	@
051B	59	Y -. --	053B	40	@
051C	5A	Z --.. ,	053C	38	8
051D	51	Q --.- ,	053D	40	@
051E	40	@	053E	39	9
051F	40	@	053F	30	0

MORSE PROGRAM

;This is the subroutine to enter the mark and space times
 ;into their respective histograms. The raw space data is
 ;stored from 0580 to 05BF, the raw mark data from 05C0
 ;to 05FF.

Add.	Code	Label	Mnemonic	Comments
0540	8C	UBIN	GLO C	;Enter here to bin space, get
0541	FC 80		ADI 80	;Key up timer, add offset,
0543	30 4D		BR SP	;and jump ahead.
0545	89	DBIN	GLO 9	;Enter here for mark, get
0546	FF 40		SMI 40	;timer, subtract 40, jump if
0548	33 5F		BGE BIGM	; >=0, therefore overflow.
054A	89		GLO 9	;If ok, get timer again, add
054B	FC C0		ADI C0	;offset.
054D	AF	SP	PLO F	;For mark or space, put data
054E	F8 05		LDI 05	;plus offset into RF.0, and
0550	BF		PHI F	;high address into RF.1.

;The value from the timer, e.g. 2C, is used to point to the
 ;2Cth bin of the appropriate histogram. The number in the
 ;2Cth bin indicates the number of times the particular
 ;timer reached exactly 2C before the key changed.

0551	0F		LDN F	;Get the data from that bin.
0552	FF 3F		SMI 3F	;Subt. 3F (63), 0 if full.
0554	7B		SEQ	;Set Q in case it's full.
0555	32 5C		BZ BIND	;If full, jump out with Q=1.
0557	7A		REQ	;Not full, so reset Q,
0558	0F		LDN F	;and get data again from bin.
0559	FC 01		ADI 01	;Increment the data
055B	5F		STR F	;and restore it.
055C	C0 0760	BIND	BR SMTH	;Done, jmp to smooth routine.

;Program ends up here if key down timer overflowed. This
 ;usually means the overall speed is far too high.

055F	8D	BIGM	GLO D	;The timing constant is
0560	FE		SHL	;stored in RD, and this is
0561	AD		PLO D	;shifted left by one in a 16
0562	9D		GHI D	;bit shift. This multiplies
0563	7E		SHLC	;the constant by 2 and halves
0564	BD		PHI D	;the speed of the timer.
0565	C0 03EF		LBR RSTR2	;Do a partial restart with

;the new slower speed and zeroed histograms.

MORSE PROGRAM

```

;The binning subroutine sets the Q flag if a bin is full.
;On return to the Key up or Key down routines the flag
;is sampled. If set, control passes to this routine which
;checks the auto histogram flag stored in 07EE. This
;is not a subroutine as it can return to different places.
;If entered from the key down section, Q is still set, but
;the Key up section resets Q first. Q tells this section
;where to return if no histogram was required.

```

Add.	Code	Label	Mnemonic	Comments
0568	27	AUT?	DEC 7	;Point to location which
0569	27		DEC 7	;holds auto flag.
056A	72		LDXA	;Get it and point back to
056B	60		IRX	;free location (07F0).
056C	FF 5A		SMI 5A	;Subtract 5A, 0 if "Z".
056E	C2 0250		LBZ HIST	;If so, jump for histogram.
0571	C9 0454		LBNQ UP2	;If Q=0, back to Key up,
0574	7A		REQ	;else reset Q
0575	C0 0421		LBR STDN	;and back to key down.
0578	00		IDL	;IDL's to 057F inc.

```

;0580 to 05BF contains raw histogram data for spaces.
;05C0 to 05FF contains raw histogram data for marks.

```

```

;This is the subroutine to clear (Q=0) or divide by two
;(Q=1) the raw histogram data.

```

Add.	Code	Label	Mnemonic	Comments
03C0	F8 05	MASG	LDI 05	;Put high address pointer
03C2	BF		PHI F	;into RF.1.
03C3	F8 80		LDI 80	;Low address is start of raw
03C5	AF		PLO F	;space data, marks follow.
03C6	8F	MAS?	GLO F	;If RF.0=00, then done.
03C7	C6		LSNZ	;Skip two if not done,
03C8	D5		RETURN	;else return.
03C9	00		IDL	;
03CA	C5		LSNQ	;Skip two if Q=0 (clear),
03CB	0F		LDN F	;else get data from bin and
03CC	F6		SHR	;shift right to halve it.
03CD	CD		LSQ	;Skip two if Q=1 (halve),
03CE	F8 00		LDI 00	;else load 00 to clear.
03D0	5F		STR F	;Put new data into bin
03D1	1F		INC F	;and increment pointer.
03D2	30 C6		BR MAS?	;Loop back for more.
03D4	00		IDL	;IDL's to 03DC inc.

MORSE PROGRAM

;This is the smoothing subroutine. It is a continuation of
 ;the binning subroutine which returns from here. The raw
 ;histogram data from 0580-05BF, 05C0-05FF is smoothed by
 ;a simple algorithm and stored from 0680-06BF, 06C0-06FF.
 ;The Q flag is still set if a bin was full).

Add.	Code	Label	Mnemonic	Comments
0760	F8 05	SMTH	LDI 05	;Load 0580 into RF, point to
0762	BF		PHI F	;raw data, 0680 into R0,
0763	F8 06		LDI 06	;space for smoothed data.
0765	B0		PHI 0	;
0766	F8 80		LDI 80	;
0768	A0		PLO 0	;
0769	AF		PLO F	;
076A	E0		SEX 0	;R0 used for math, make it RX
076B	0F	SMTH1	LDN F	;Get raw data(N)
076C	FE		SHL	;multiply by 2,
076D	50		STR 0	;store it in smooth space.
076E	2F		DEC F	;Point to raw data(N-1).
076F	8F		GLO F	;If RF=057F, out of data
0770	FF 7F		SMI 7F	;storage area,
0772	32 7B		BZ BDAT	;so jump ahead.
0774	FF 40		SMI 40	;If RF=05BF, marks intruding
0776	32 7B		BZ BDAT	;into spaces, so jump ahead.
0778	0F		LDN F	;Get raw data(N-1) and add it
0779	F4		ADD	;to 2*raw data(N).
077A	50		STR 0	;Put result in smoothed area.
077B	1F	BDAT	INC F	;Point to raw data(N+1).
077C	1F		INC F	;
077D	8F		GLO F	;If RF=0600, out of data,
077E	32 87		BZ BDAT2	;so jump ahead.
0780	FF 80		LDI 80	;If RF=0580, spaces intruding
0782	32 87		BZ BDAT2	;into marks, so jump ahead.
0784	0F		LDN F	;Get raw data(N+1) and add it
0785	F4		ADD	;to value from line 0779
0786	50		STR 0	;Put result in smoothed area.
0787	F0	BDAT2	LDX	;Get result so far,
0788	FC 02		ADI 02	;add two for rounding.
078A	F6		SHR	;Divide by four
078B	F6		SHR	;
078C	50		STR 0	;and store result.

;The data was smoothed with the following formula:
 ; smooth(n) = (2*raw(n) + raw(n-1) + raw(n+1) + 2)/4

MORSE PROGRAM

;This is the continuation of the smoothing routine.

Add.	Code	Label	Mnemonic	Comments
078D	60		IRX	;Point to next smooth space.
078E	80		GLO 0	;If R0 not 0700, not done,
078F	3A 6B		BNZ SMTH1	;so jump back for more,
0791	E7		SEX 7	;else restore R7 as RX.
0792	CD		LSQ	;If Q=0, no bins full,
0793	D5		RETURN	;so return, otherwise skip 2.
0794	00		IDL	;
0795	D4 03C0		SUB MASC	;Sub to MASC, Q=1, so data/2.
0798	D4 01DA		SUB CRLF	;Output "CR" & "LF"
079B	F8 3E		LDI 3E	;Load ASCII ">" and output it
079D	DA		SEP A	;to show histogram change.
079E	D5		RETURN	;Return with Q still set for
079F	00		IDL	;auto histogram check.

;The next set of routines computes the mark and space
 ;thresholds from the smoothed data. The main routine begins
 ;at 0600, but the subroutines are presented first.
 ;They all rely on the position of the data within memory.
 ;This subroutine returns the maximum value of a mark or
 ;space histogram in RC.1 and the bin number (00 to 3F),
 ;not the memory address, in RC.0. If Q is set the leftmost
 ;maximum is returned and if Q is reset, the rightmost.
 ;These values are different if two bins equal the maximum.
 ;R9 contains the starting address of data and RF.0 contains
 ;the bin number of that start data as not all scans start
 ;from bin 00 when looking for the leftmost maximum.

Add.	Code	Label	Mnemonic	Comments
0740	F8 00	FMAX	LDI 00	;Reset bin number
0742	AC		PLO C	;and maximum value.
0743	BC		PHI C	;
0744	CD		LSQ	;If right maximum set RF.0
0745	AF		PLO F	;to 00, else leave it alone.
0746	C4		NOP	;
0747	E9	MAXM	SEX 9	;Make R9 be RX for math use.
0748	9C		GHI C	;Get current maximum and sub-
0749	F5		SD	;tract data pointed to by R9.
074A	3B 53		BL NCHG	;If less, no change, so jump.
074C	C5		LSNQ	;Skip two if rightmost max.,
;that means change on >=. If leftmost max, change on > only				
074D	32 53		BZ NCHG	;Jump if leftmost and equal.
074F	09		LDN 9	;If a change, get new maximum
0750	BC		PHI C	;and put it in RC.1.
0751	8F		GLO F	;Get new bin number and put
0752	AC		PLO C	;it in RC.0.

MORSE PROGRAM

;This is a continuation of the "FMAX" subroutine.

Add.	Code	Label	Mnemonic	Comments
0753	1F	NCHG	INC F	;Increment bin number
0754	19		INC 9	;and memory pointer.
0755	89		GLO 9	;If R9=0600, mark scan done,
0756	32 5C		BZ MAXD	;jump ahead.
0758	FF C0		SMI C0	;If R9=05BF, space scan done,
075A	3A 47		BNZ MAXM	;else jump back for more.
075C	E7	MAXD	SEX 7	;Restore R7 as RX
075D	D5		RETURN	;and get out.
075E	00		IDL	;This next subroutine takes
075F	00		IDL	;the data passed from FMAX
;and sets it up for later use. Enter with Q=1 for marks.				
0700	F8 00	ADD50	LDI 00	;Reset RF.0
0702	AF		PLO F	;
0703	9C		GHI C	;Get maximum peak height.
0704	F6		SHR	;divide by two
0705	BC		PHI C	;and put it back.
0706	F8 80		LDI 80	;Get space offset
0708	C5		LSNQ	;skip 2 if spaces
0709	F8 C0		LDI C0	;else get mark offset.
070B	57		STR 7	;Put offset in 07F0.
070C	8C		GLO C	;Get maximum bin number
070D	F4		ADD	;add offset for low address
070E	A9		PLO 9	;and put it in R9.0.
070F	D5		RETURN	;

;This subroutine finds the first occurrence of a bin
 ; <= (maximum height)/2 while searching left to right.
 ;Enter with start address of search in R9, (maximum)/2
 ;in RC.1, and 00 in RF.0. On exit RF.0 contains the number
 ;of bins from the peak to <= (peak)/2.

Add.	Code	Label	Mnemonic	Comments
0710	00		IDL	;IDL's to 0713 inc.
0714	E9	LTOR	SEX 9	;Make R9 be RX
0715	9C	MOLR	GHI C	;Get 1/2 peak and
0716	F7		SM	;subtract data.
0717	33 22		BGE LROK	;Out if bin found,
0719	19		INC 9	;else increment pointer
071A	1F		INC F	;and bin offset.
071B	89		GLO 9	;If R9=0700
071C	32 22		BZ	;get out, end of marks.
071E	FF C0		SMI C0	;If R9=06C0, end of spaces,
0720	3A 15		BNZ MOLR	;else jump back for more.
0722	E7	LROK	SEX 7	;Restore R7 as RX.
0723	D5		RETURN	;
0724	00		IDL	;IDL's to 0727 inc.

MORSE PROGRAM

;This subroutine is similar to the one above except that
 ;it searches from right to left. Q=0 for spaces and Q=1
 ;for marks. RE.0 and RE.1 contain leftmost stop addresses.

Add.	Code	Label	Mnemonic	Comments
0728	E9	RTOL	SEX 9	;Make R9 be RX
0729	9C		GHI C	;Get 1/2 peak and subtract
072A	F7		SM	;data.
072B	33 3E		BGE RLOK	;Out if bin found,
072D	E7		SEX 7	;else restore R7 as RX.
072E	89		GLO 9	;Get memory pointer and
072F	FF 80		SMI 80	;subtract spaces offset.
0731	C5		LSNQ	;Skip 2 if spaces, else
0732	FF 40		SMI 40	;subtract extra mark offset.
0734	57		STR 7	;Stuff it in 07F0
0735	8E		GLO E	;Get overange for spaces,
0736	C5		LSNQ	;skip 2 if spaces,
0737	9E		GHI E	;else get it for marks.
0738	C4		NOP	;
0739	F7		SM	;Compare pointer and overange
073A	29		DEC 9	;Decrement memory pointer,
073B	1F		INC F	;increment bin coutner,
073C	3B 28		BM RTOL	;If minus, back for more,
073E	E7	RLOK	SEX 7	;else restore RX
073F	D5		RETURN	;and get out.

;This is the routine to calculate the mark threshold.
 ;It calls the above three routines, FMAX, LTOR, & RTOL.
 ;In the first part "peak" or "maximum" refers to the
 ;histogram peak generated by dots.

Add.	Code	Label	Mnemonic	Comments
0600	F8 06	THRM	LDI 06	;06C0 into R9, start address
0602	B9		PHI 9	;of smoothed mark data.
0603	F8 C0		LDI C0	;
0605	A9		PLO 9	;
0606	7A		REQ	;Reset Q, shows right maximum
0607	D4 0740		SUB FMAX	;Sub to find right maximum.
060A	7B		SEQ	;Set Q, shows marks.
060B	F8 06		LDI 06	;Put partial address in R9.1.
060D	B9		PHI 9	;
060E	D4 0700		SUB ADD50	;Sub to fix data.
0611	D4 0714		SUB LTOR	;Sub to find 50% peak bin.

MORSE PROGRAM

```

;This is the continuation of the mark threshold routine.
Add. Code      Label Mnemonic      Comments
0614 8F        GLO F          ;Get distance from peak
0615 57        STR 7          ;and stuff into 07F0.
0616 8C        GLO C          ;Get peak bin # and
0617 F4        ADD            ;find bin twice as far from
0618 F4        ADD            ;peak as 50% bin.
0619 BE        PHI E          ;Put it in RE.1 as left stop
061A AF        PLO F          ;and also into RF.0.
061B FC C0     ADI C0         ;Add offsets to make this
061D A9        PLO 9          ;a memory address in R9 in
061E F8 06     LDI 06         ;the smoothed mark range.
0620 B9        PHI 9          ;
;From this point "peak" refers to the dash cluster.
0621 7B        SEQ            ;Set Q to find left maximum.
0622 D4 0740   SUB FMAX        ;Sub to find left maximum.
0625 F8 06     LDI 06         ;Put partial address in R9.1.
0627 B9        PHI 9          ;
0628 7B        SEQ            ;Set Q for marks.
0629 D4 0700   SUB ADD50       ;Sub to fix data.
062C D4 0728   SUB RTOL        ;Sub to find 50% peak bin.
062F 8F        GLO F          ;Get distance from peak,
0630 FE        SHL            ;double it
0631 57        STR 7          ;and stuff it in 07F0.
0632 8C        GLO C          ;Get peak bin # and find
0633 F7        SM            ;bin twice as far from peak
0634 57        STR 7          ;as 50% and stuff it in 07F0.
0635 9E        CHI E          ;Add same result from dot
0636 F4        ADD            ;peak and divide by two to
0637 F6        SHR            ;find estimate for mark
0638 BE        PHI E          ;threshold, put it in RE.1.
0639 C0 0640   LBR THRS        ;Jump for space threshold.
063C 00        IDL            ;IDL's to 063F inc.

;This is the routine to calculate the space threshold.
;In the first part "peak" or "maximum", refers to the
;histogram peak generated by the intra-character spaces.
Add. Code      Label Mnemonic      Comments
0640 F8 06     THRS LDI 06         ;0680 into R9, start address
0642 B9        PHI 9          ;of smoothed space data.
0643 F8 80     LDI 80         ;
0645 A9        PLO 9          ;
0646 7A        REQ            ;Reset Q, shows right maximum
0647 D4 0740   SUB FMAX        ;Sub to find right maximum.
064A 7A        REQ            ;Reset Q, shows spaces.
064B F8 06     LDI 06         ;Put partial address in R9.1.
064D B9        PHI 9          ;

```

MORSE PROGRAM

;This is the continuation of the space threshold routine.

Add.	Code	Label	Mnemonic	Comments
064E	D4 0700		SUB ADD50	;Sub to fix data.
0651	D4 0714		SUB LTOR	;Sub to find 50% peak bin.
0654	8F		GLO F	;Get distance from peak
0655	57		STR 7	;and stuff into 07F0.
0656	8C		GLO C	;Get peak bin # and
0657	F4		ADD	;find bin twice as far from
0658	F4		ADD	;peak as 50% bin.
0659	AE		PLO E	;Put it in RE.0 as left stop
065A	AF		PLO F	;and also into RF.0.
065B	FC 80		ADI 80	;Add offsets to make this
065D	A9		PLO 9	;a memory address in R9 in
065E	F8 06		LDI 06	;the smoothed space range.
0660	B9		PHI 9	;

;Here "peak" refers to the inter-character space cluster.

0661	7B		SEQ	;Set Q to find left maximum.
0662	D4 0740		SUB FMAX	;Sub to find left maximum.
0665	F8 06		LDI 06	;Put partial address in R9.1.
0667	B9		PHI 9	;
0668	7A		REQ	;Reset Q for spaces.
0669	D4 0700		SUB ADD50	;Sub to fix data.
066C	D4 0728		SUB RTOL	;Sub to find 50% peak bin.
066F	8F		GLO F	;Get distance from peak,
0670	FE		SHL	;double it
0671	57		STR 7	;and stuff it in 07F0.
0672	8C		GLO C	;Get peak bin # and find
0673	F7		SM	;bin twice as far from peak
0674	57		STR 7	;as 50% and stuff it in 07F0.
0675	8E		GLO E	;Add result from intra-space
0676	F4		ADD	;peak and divide by two for
0677	F6		SHR	;first estimate of space
0678	57		STR 7	;threshold, put it in 07F0.
;Due the poor "quality" of the inter-character space peak,				
;it is averaged with the mark threshold which would be the				
;same for perfect code.				
0679	9E		CHI E	;Get mark threshold and
067A	F4		ADD	;simply average with estimate
067B	F6		SHR	;just calculated. Store this
067C	AE		PLO E	;final result in RE.0
067D	C0 07B0		LBR CLUP	;Jump to cleanup loose ends.

MORSE PROGRAM

;This short routine follows the space threshold calculation
;to clear up various odds and ends.

Add.	Code	Label	Mnemonic	Comments
07A0	00		IDL	;IDL's to 07AF inc.
07B0	D4 01AB	CLUP	SUB DISPY	;Display new thresholds.
07B3	F8 00		LDI 00	;
07B5	A9		PLO 9	;Reset key down timer and
07B6	BC		PHI C	;counter for key downs in
07B7	F8 01		LDI 01	;character. Set Morse shift
07B9	B9		PHI 9	;register to 01.
07BA	27		DEC 7	;Point to stored RC.0,
07BB	72		LDXA	;get it, point to 07F0,
07BC	AC		PLO C	;restore RC.0.
07BD	C0 0400		LBR TIME	;Jump back to timer.

;These are the commands which the program recognizes from
;an ASCII keyboard. They are immediately recognized, so
;no carriage return is required.

/ Restart at 03DD, sub to monitor to enter new speed,
press "R" to return.

H Output histogram immediately, wait for more commands.

K Send hardcopy command to graphics terminal, wait for
more commands.

M Restart at 03F6, same histogram, speed and thresholds.

N Cancel autohistograms, continue at 0400.

X Zero raw histogram data, restart at 03F6.

Y Divide raw histogram data by two, restart at 03F6.

Z Set automatic mode for histograms. If any raw bin is
full, the program stops and outputs a histogram. The
raw data is divided by two, but until the program is
restarted the smoothed data is untouched, so other
copies can be made. This command continues at 0400.

"N" or "Z" can be pressed at any time during normal
program execution. The other commands all disrupt
decoding. Any command not in the above table will
cause a restart of the entire system to the monitor.

MORSE PROGRAM

;The above routines were all essential to decoding Morse
 ;code. The remaining sections are only needed to output
 ;histograms to the Tektronix graphics terminal. Register A
 ;points to the routine for serial port 2, not 1 as usual.
 ;A lot of manipulation is necessary because the plotter
 ;works on a 1024X by 779Y absolute matrix, with each letter
 ;14X by 22Y, and the histogram requires 64X by 31Y
 ;numbered positions, each divisible by four!!! Calculations
 ;can only be done in integer math, so there are problems.
 ;This subroutine numbers an axis 0 to F repeatedly.

Add.	Code	Label	Mnemonic	Comments
0220	F8 00	LTTR	LDI 00	;00 into RF.0 if not Q,
0222	C5		LSNQ	;0F if Q, start number of
0223	F8 0F		LDI 0F	;hex-numbered axis.
0225	AF		PLO F	;
0226	F8 1D	MLTR?	LDI 1D	;Output "GS" to set
0228	DA		SEP A	;graphics mode.
0229	D4 02E0		SUB XY0U1	;Output XY coordinates.
022C	D4 02F8		SUB XY0U2	;
022F	F8 1F		LDI 1F	;Output "US" to set
0231	DA		SEP A	;alphameric mode.
0232	8F		GLO F	;Get axis counter, mask in
0233	FA 0F		ANI 0F	;lower 4 bits, put them
0235	BF		PHI F	;in RF.1 and output them as
0236	D4 01F2		SUB 1HOUT	;a single hex digit.
;RC.0 contains the low address of the current coordinate				
;being lettered: F1 for X or F2 for Y; high address 07.				
;RC.1 contains the axis coordinate increment: 04 for				
;horizontal, 05 for vertical. R9.0 holds stop count				
;for axis, 3F horizontal, FF or 10 vertical.				
0239	8C		GLO C	;Transfer low address to RX.
023A	A7		PLO 7	;RX points to coordinate.
023B	9C		GHI C	;Get coordinate increment,
023C	F4		ADD	;add it on,
023D	73		STXD	;restore coordinate.
023E	F8 F0		LDI F0	;Point RX back to 07F0.
0240	A7		PLO 7	;
0241	1F		INC F	;Increment axis counter, but
0242	C5		LSNQ	;if Q was set, decrement
0243	2F		DEC F	;counter by 2, effectively
0244	2F		DEC F	;by 1.
0245	8F		GLO F	;Get axis counter and
0246	57		STR 7	;stuff it in 07F0.
0247	89		GLO 9	;Get stop count and
0248	F7		SM	;subtract,
0249	3A 26		BNZ MLTR?	;jump back for more if not
024B	D5		RETURN	;equal, else return.
024C	00		IDL	;IDL's to 024F inc.

MORSE PROGRAM

;These two subroutines output an XY coordinate to the
 ;graphics terminal which requires 10 bit X and Y values,
 ;broken into four 5 bit ASCII characters. Enter with an
 ;8 bit X value in 07F1 and Y in 07F2. Normally call XYOU2
 ;immediately after XYOU1, but to obtain the full 10 bit
 ;precision available, the two lsb's in the Y value can be
 ;added at that point.

Add.	Code	Label	Mnemonic	Comments
02E0	60	XYOU1	IRX	;Point to X value,
02E1	72		LDXA	;get it and point to Y.
02E2	A0		PLO 0	;X into R0.0,
02E3	F0		LDX	;get Y,
02E4	B0		PHI 0	;Y into R0.1.
02E5	27		DEC 7	;Point to 07F0.
02E6	27		DEC 7	;
02E7	F6		SHR	;Get high 5 Y bits in lower
02E8	F6		SHR	;3 positions.
02E9	F6		SHR	;
02EA	F9 20		ORI 20	;OR in identifier and
02EC	DA		SEP A	;output high Y byte.
02ED	90		GHI 0	;Get Y again,
02EE	FE		SHL	;put 0's in two lsb's
02EF	FE		SHL	;and stuff it in 07F0.
02F0	57		STR 7	;
02F1	D5		RETURN	;Return for possible mods to
				;two lsb's to get 10 bit Y coordinate if required.
02F2	00		IDL	;IDL'S to 02F7 inc.
02F8	F0	XYOU2	LDX	;Get Y value back,
02F9	FA 1F		ANI 1F	;mask in lower 5 bits only
02FB	F9 60		ORI 60	;OR in identifier and
02FD	DA		SEP A	;output low Y byte.
02FE	80		GLO 0	;Get X value and
02FF	F6		SHR	;put high 5 bits in lower
0300	F6		SHR	;5 positions.
0301	F6		SHR	;
0302	F9 20		ORI 20	;OR in identifier and
0304	DA		SEP A	;output high X byte.
0305	80		GLO 0	;Get X again and
0306	FE		SHL	;put 0's in two lsb's.
0307	FE		SHL	;
0308	FA 1C		ANI 1C	;Mask in required bits,
030A	F9 40		ORI 40	;OR in identifier and
030C	DA		SEP A	;output low X byte.
030D	D5		RETURN	;All done, get out.
030E	00		IDL	;
030F	00		IDL	;

MORSE PROGRAM

;The above three subroutines are called by this section
 ;which handles the overhead for lettering the axes.
 ;This is the actual start of the histogram routine.

Add.	Code	Label	Mnemonic	Comments
0250	D4 01DA	HIST	SUB CRLF	;Output a "CR" & "LF"
0253	F8 10		LDI 10	;Point RA to output routine
0255	AA		PLO A	;for serial port 2.
0256	F8 07		LDI 07	;Output "BEL" to inform
0258	DA		SEP A	;operator of histogram.
0259	F8 1B		LDI 1B	;Output "ESC" & "FF" to
025B	DA		SEP A	;erase screen.
025C	F8 0C		LDI 0C	;
025E	DA		SEP A	;
025F	F8 FF		LDI FF	;This is a delay while
0261	BF		PHI F	;the screen clears.
0262	AF		PLO F	;FFFF into RF
0263	2F	SCR?	DEC F	;
0264	9F		GHI F	;Get RF.1, jump back if not
0265	3A 63		BNZ SCR?	;zero, else delay over.
0267	F8 F1		LDI F1	;Set up to label horizontal
0269	AC		PLO C	;axis (see LTTR for expl.)
026A	F8 04		LDI 04	;F1 in RC.0 = X axis
026C	BC		PHI C	;04 in RC.1 = increment
026D	60		IRX	;04 in 07F1 = X offset
026E	57		STR 7	;61 in 07F2 = Y offset
026F	60		IRX	;3F in R9.0 = last count+1
0270	F8 61		LDI 61	; 0 in Q = pos. incr.
0272	73		STXD	;
0273	27		DEC 7	;Point to 07F0 again.
0274	F8 3F		LDI 3F	;
0276	A9		PLO 9	;
0277	7A		REQ	;
0278	D4 0220		SUB LTTR	;Go and letter X axis.
027B	F8 FF		LDI FF	;F2 in RC.0 = Y axis
027D	A9		PLO 9	;05 in RC.1 = increment
027E	F8 F2		LDI F2	;00 in 07F1 = X offset
0280	AC		PLO C	;11 in 07F2 = Y offset
0281	F8 05		LDI 05	;FF in R9.0 = last count-1
0283	BC		PHI C	; 1 in Q = neg. incr.
0284	F8 11		LDI 11	;
0286	60		IRX	;
0287	60		IRX	;
0288	73		STXD	;
0289	F8 00		LDI 00	;
028B	73		STXD	;Point to 07F0 again.
028C	7B		SEQ	;
028D	D4 0220		SUB LTTR	;Letter Y axis from F to 0.

MORSE PROGRAM

Add.	Code	Label	Mnemonic	Comments
0290	F8 10		LDI 10	;Continue lettering axes.
0292	A9		PLO 9	;F2 in RC.0 = Y axis
0293	60		IRX	;05 in RC.1 = increment
0294	60		IRX	;00 in 07F1 = X offset
0295	F8 66		LDI 66	;66 in 07F2 = Y offset
0297	73		STXD	;10 in R9.0 = last count+1
0298	F8 00		LDI 00	; 0 in Q = pos. incr.
029A	73		STXD	;Point to 07F0 again.
029B	C4		NOP	;
029C	7A		REQ	;
029D	D4 0220		SUB LTTR	;Letter Y axis from 0 to F.
02A0	C0 0310		LBR TEKH	;Continue with histogram.
02A3	00		IDL	;IDL's to 02AF inc.
;This section outputs the data to the histogram.				
;The smoothed data is used, but this is easily changed.				
0310	F8 06	TEKH	LDI 06	;Enter with Q=0, spaces.
0312	BC		PHI C	;0680 in RC = spaces, if Q=0
0313	60		IRX	;06C0 in RC = marks, if Q=1
0314	F8 05		LDI 05	;05 in 07F1 = X offset
0316	57		STR 7	;60 in 07F2 = Y offset, Q=0
0317	60		IRX	;66 in 07F2 = Y offset, Q=1
0318	F8 80		LDI 80	;
031A	C5		LSNQ	;
031B	F8 C0		LDI C0	;
031D	AC		PLO C	;
031E	F8 60	TEK1	LDI 60	;
0320	C5		LSNQ	;
0321	F8 66		LDI 66	;
0323	73		STXD	;
0324	27		DEC 7	;Point to 07F0 again.
0325	C4		NOP	;
0326	F8 1D		LDI 1D	;Output "GS", make this dark
0328	DA		SEP A	;vector to new position.
0329	D4 02E0		SUB XYOU1	;Output coordinates.
032C	D4 02F8		SUB XYOU2	;
032F	60		IRX	;
0330	60		IRX	;Point R7 to Y coordinate.
0331	C4		NOP	;
0332	0C		LDN C	;Get smoothed data pointed
0333	F6		SHR	;to by RC,
0334	F6		SHR	;divide by 4,
0335	EC		SEX C	;
0336	F4		ADD	;add that to original,
0337	E7		SEX 7	;in effect multiply by 1.25.

MORSE PROGRAM

```

;.....continuation of histogram.....
;In order to fit histogram onto screen with even spacing,
;the very last (3Fth) bin is not output. This is not a
;problem as there should be no useful information there.

```

Add.	Code	Label	Mnemonic	Comments
0338	CD		LSQ	;Skip 2 if Q=1, marks.
0339	F5		SD	;Spaces, vector down, so
033A	C4		NOP	;subtract data from offset.
033B	C5		LSNQ	;Skip 2 if Q=0, spaces.
033C	F4		ADD	;Marks, vector up, so add
033D	C4		NOP	;data to offset.
033E	73		STXD	;Store Y coordinate in 07F2,
033F	27		DEC 7	;point back to 07F0.
0340	C4		NOP	;
0341	D4 02E0		SUB XYOU1	;Sub to output high Y byte.
0344	0C		LDN C	;Get data again,
0345	FA 03		ANI 03	;mask in lower 3 bits only.
0347	CD		LSQ	;Skip 2 if Q = 1 = marks,
0348	F5		SD	;else spaces, subtract.
0349	C4		NOP	;
034A	C5		LSNQ	;Skip 2 if Q = 0 = spaces,
034B	F4		ADD	;else marks, add. Store
034C	C4		NOP	;this result back in 07F2.
034D	57		STR 7	;This adds the two lsb's to
034E	D4 02F8		SUB XYOU2	;Y for 10 bits, output it.
0351	1C		INC C	;Increment RC, point to next
0352	60		IRX	;data, increment RX,
0353	F0		LDX	;get X coordinate.
0354	FC 04		ADI 04	;Add 04, move along axis,
0356	57		STR 7	;restore it.
0357	60		IRX	;Point RX to Y coordinate.
0358	8C		GLO C	;Get low data pointer,
0359	FF BF		SMI BF	;subtract BF, neg if still
035B	CF		LSDF	;spaces, skip 2 if positive.
035C	30 1E		BR TEK1	;Neg, jump back, more spaces
035E	7B		SEQ	;Set Q for marks,
035F	27		DEC 7	;point to 07F0 again.
0360	27		DEC 7	;If above subtraction was 0,
0361	32 10		BZ TEKH	;back to START mark output.
0363	60		IRX	;Else point to 07F2 = Y.
0364	60		IRX	;
0365	FF 40		SMI 40	;Subtract 40 from previous
0367	3A 1E		BNZ TEK1	;result, if nonzero, back
0369	7A		REQ	;for more spaces, else reset
036A	27		DEC 7	;Q, point back to 07F0,
036B	27		DEC 7	;
036C	C0 0370		LBR OUPAR	;and jump ahead to output
036F	00		IDL	;parameters.

MORSE PROGRAM

;This section outputs the speed and thresholds to the
;graphics terminal at the top of the histogram.

Add.	Code	Label	Mnemonic	Comments
0370	60	OUPAR	IRX	;00 in 07F1 = X offset
0371	60		IRX	;BA in 07F2 = Y offset
0372	F8 BA		LDI BA	;
0374	73		STXD	;
0375	F8 00		LDI 00	;
0377	73		STXD	;Point back to 07F0.
0378	C4		NOP	;NOP's to 037B inc.
037C	F8 1D		LDI 1D	;Output "GS" for dark vector
037E	DA		SEP A	;
037F	D4 02E0		SUB XYOU1	;Output coordinates for
0382	D4 02F8		SUB XYOU2	;first lettering position.
0385	F8 1F		LDI 1F	;Output "US" for alpha-
0387	DA		SEP A	;numeric mode.
0388	F8 02		LDI 02	;Point RC to the start of
038A	BC		PHI C	;the first heading stack
03BB	F8 B0		LDI B0	;at 02B0.
03BD	AC		PLO C	;
03BE	D4 03B6		SUB HEAD	;Output first heading.
0391	D4 01CE		SUB RDOUT	;Output speed from RD.
0394	F8 B7		LDI B7	;Point RC to second heading.
0396	AC		PLO C	;
0397	D4 03B6		SUB HEAD	;Output second heading.
039A	9E		CHI E	;Move mark threshold to RF.1
039B	BF		PHI F	;
039C	D4 01E2		SUB ZHOUT	;Output it as 2 hex digits.
039F	F8 CB		LDI CB	;Point RC to third heading.
03A1	AC		PLO C	;
03A2	D4 03B6		SUB HEAD	;Output third heading.
03A5	8E		GLO E	;Put space threshold in RF.1
03A6	BF		PHI F	;
03A7	D4 01E2		SUB ZHOUT	;Output it as 2 hex digits.
03AA	F8 01		LDI 01	;Histogram finished, point
03AC	AA		PLO A	;RA to routine for port 1,
03AD	C0 04B8		LBR KCHK	;wait for keyboard input.
03B0	00		IDL	;IDL's to 03B5 inc.
;This short subroutine outputs ASCII headings. Enter with				
;start of data stack in RC, returns when first 00 found.				
03B6	4C	HEAD	LDA C	;Get heading element,
03B7	C6		LSNZ	;skip 2 if nonzero,
03B8	D5		RETURN	;else return.
03B9	00		IDL	;
03BA	DA		SEP A	;Output heading and
03BB	30 B6		BR HEAD	;jump back for more.
03BC	00		IDL	;IDL's to 03BF inc.

MORSE PROGRAM

;This is the heading data. Each stack ends with 00.

Add.	Code	Character	Add.	Code	Character
02B0	53	S	02C9	20	space
02B1	50	P	02CA	00	end
02B2	45	E	02CB	20	space
02B3	45	E	02CC	20	space
02B4	44	D	02CD	20	space
02B5	20	space	02CE	20	space
02B6	00	end	02CF	53	S
02B7	20	space	02D0	50	P
02B8	20	space	02D1	41	A
02B9	20	space	02D2	43	C
02BA	20	space	02D3	45	E
02BB	4D	M	02D4	20	space
02BC	41	A	02D5	54	T
02BD	52	R	02D6	48	H
02BE	4B	K	02D7	52	R
02BF	20	space	02D8	45	E
02C0	54	T	02D9	53	S
02C1	48	H	02DA	48	H
02C2	52	R	02DB	4F	O
02C3	45	E	02DC	4C	L
02C4	53	S	02DD	44	D
02C5	48	H	02DE	20	space
02C6	4F	O	02DF	00	end
02C7	4C	L			
02C8	44	D			

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